



UNIVERSITY OF
TORONTO

Computer Science

Quantifier Scope Storage

CSC485/2501

December 2, 2021

Outline

- The Language of the Lambda Calculus
- Alpha Conversion
- Beta Reduction
- Logical Representations of Meaning
- Quantifier Storage

The Language of the Lambda Calculus

- Constants, e.g., a, f
- Variables, e.g., x, y (they often look like constants)
- Applications, e.g., $f(a)$
- Abstractions, e.g. $\lambda x.f(x)$

Alpha Conversion

The Basic Idea

$\lambda x.f(x) \Rightarrow \lambda y.f(y)$ [y fresh]

Beta Reduction

The Basic Idea

$$(\lambda x.f(x))(a) \Rightarrow f(a)$$

Logical Representations of Meaning

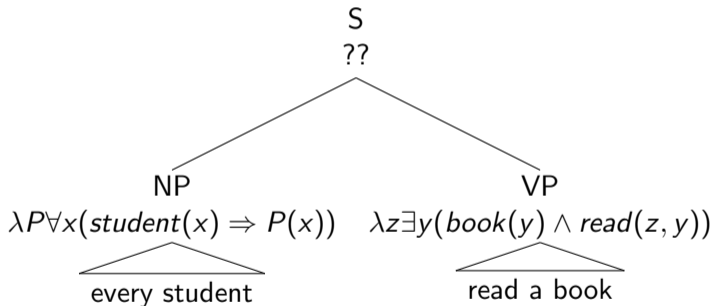
What's one of the logical forms?

Every student read a book

Logical Representations of Meaning

- Every student read a book

$$\forall x(\text{student}(x) \Rightarrow \exists y(\text{book}(y) \wedge \text{read}(x, y)))$$

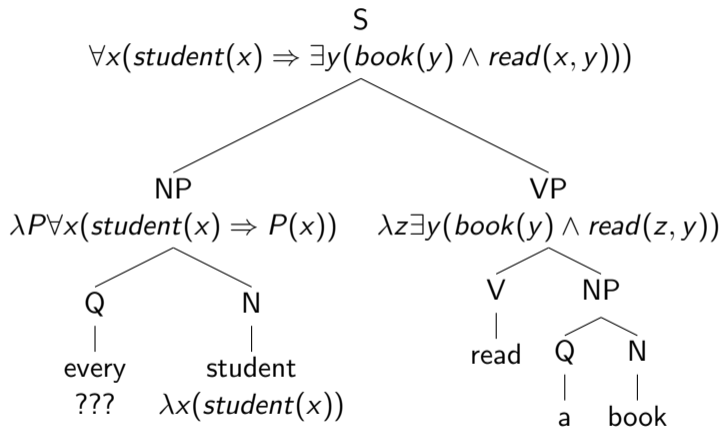


Beta Normalization

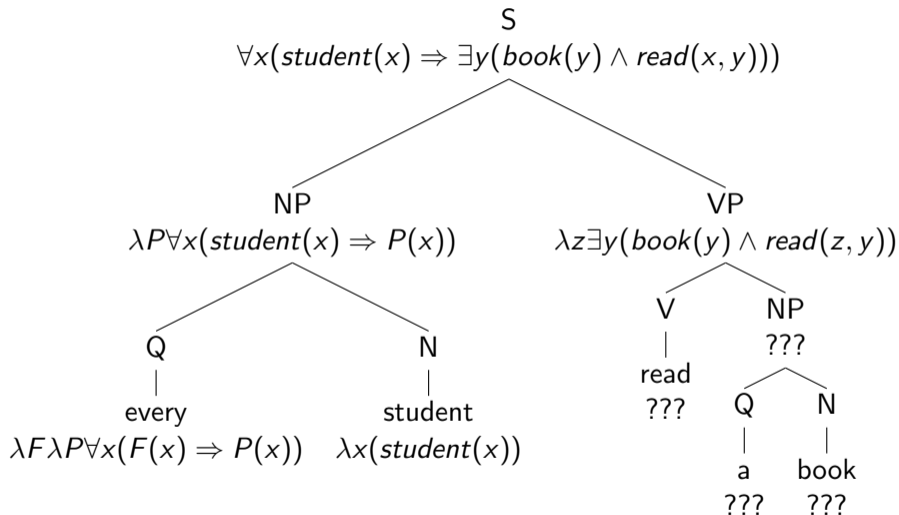
- Every student
 $\lambda P \forall x (student(x) \Rightarrow P(x))$
- read a book
 $\lambda z \exists y (book(y) \wedge read(z, y))$
- Every student read a book

$$\begin{aligned} & \lambda P \forall x (student(x) \Rightarrow P(x)) (\lambda z \exists y (book(y) \wedge read(z, y))) \\ \Leftrightarrow_{\beta} & \forall x (student(x) \Rightarrow \lambda z \exists y (book(y) \wedge read(z, y))(x)) \\ \Leftrightarrow_{\beta} & \forall x (student(x) \Rightarrow \exists y (book(y) \wedge read(x, y))) \end{aligned}$$

Beta Reduction: Practice



Beta Reduction: Practice



VP
 $\lambda z \exists y(book(y) \wedge read(z, y))$

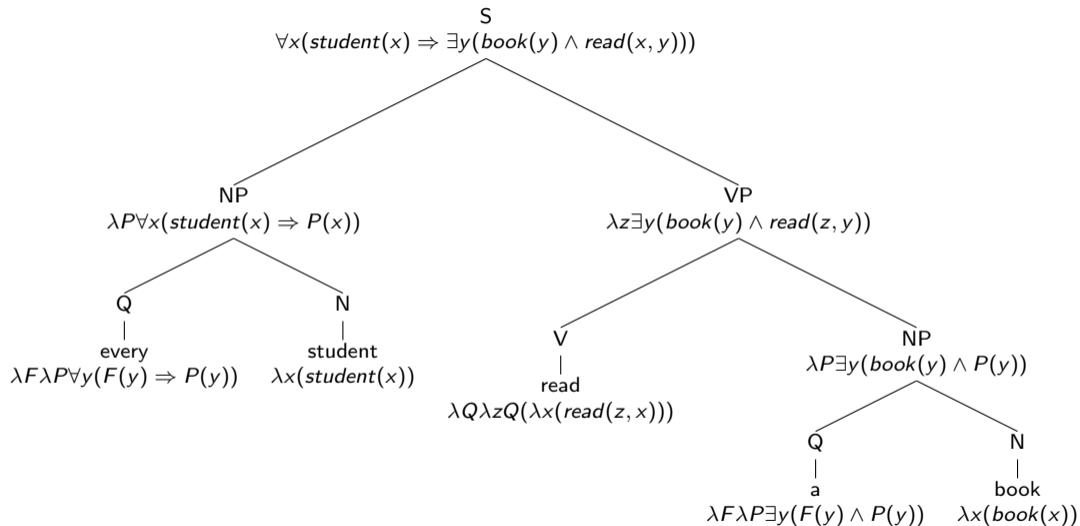
V
|
read
???

NP
???

Q
|
a
???

N
|
book
???

Beta Reduction: Solution to Practice



VP
 $\lambda z \exists y(book(y) \wedge read(z, y))$

V
 read
 $\lambda Q \lambda z Q(\lambda x(read(z, x)))$

NP
 $\lambda P \exists y(book(y) \wedge P(y))$

Q
 a
 $\lambda F \lambda P \exists y(F(y) \wedge P(y))$

N
 book
 $\lambda x(book(x))$

Quantifier Storage

Every student read a book

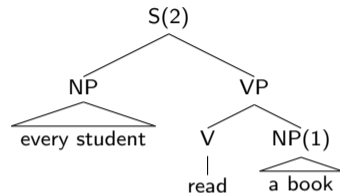
- $\forall x(\textit{student}(x) \Rightarrow \exists y(\textit{book}(y) \wedge \textit{read}(x, y)))$
- $\exists y(\textit{book}(y) \wedge \forall x(\textit{student}(x) \Rightarrow \textit{read}(x, y)))$

How to get the second reading?

Quantifier Storage

- Storage at (1)

Logic Statement	Quantifier Storage
$\lambda G \exists y (book(y) \wedge G(y))$	\emptyset
$\Rightarrow \lambda F.F(x)$	$[\langle x; \lambda G \exists y (book(y) \wedge G(y)) \rangle]$



- Retrieve at (2)

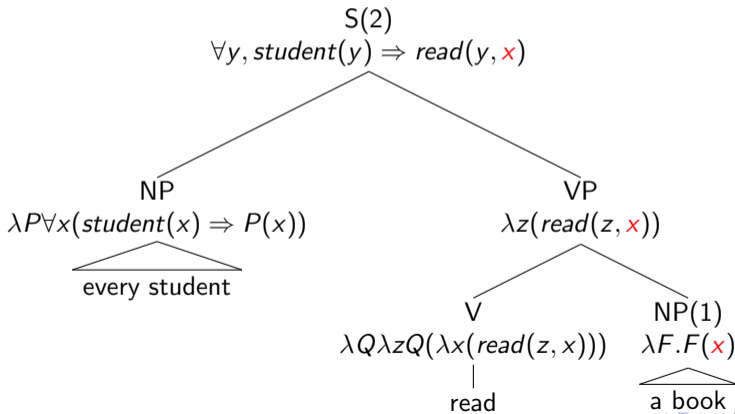
$\forall y, student(y) \Rightarrow read(y, x)$	$[\langle x; \lambda G \exists y (book(y) \wedge G(y)) \rangle]$
$\lambda G \exists z (book(z) \wedge G(z)) (\lambda x (\forall y, student(y) \Rightarrow read(y, x)))$	\square
$\exists z (book(z) \wedge (\lambda x (\forall y, student(y) \Rightarrow read(y, x)))(z))$	\square
$\exists z (book(z) \wedge (\forall y, student(y) \Rightarrow read(y, z)))$	\square

Quantifier Storage

- Storage at (1)

$$\lambda G \exists y (book(y) \wedge G(y)) \quad \emptyset$$

$$\Rightarrow \lambda F.F(x) \quad [\langle x; \lambda G \exists y (book(y) \wedge G(y)) \rangle]$$



Quantifier Storage

- Retrieve at (2)

$\forall y. student(y) \Rightarrow read(y, x)$

$[\langle x; \lambda G \exists y (book(y) \wedge G(y)) \rangle]$

$\lambda G \exists z (book(z) \wedge G(z)) (\lambda x (\forall y. student(y) \Rightarrow read(y, x)))$

□

$\exists z (book(z) \wedge (\lambda x (\forall y. student(y) \Rightarrow read(y, x)))(z))$

□

$\exists z (book(z) \wedge (\forall y. student(y) \Rightarrow read(y, z)))$

□

