

## CSC236 winter 2020, quiz week 11<sub>1</sub>

### Question.

Recall that for a language  $L$ , the Kleene star operation can be defined as  $L^* = L^0 \cup L^1 \cup L^2 \cup \dots$ . We will define a similar operation  $Evenstar(L) = L^0 \cup L^2 \cup L^4 \cup \dots$ , i.e. the union of all the even powers of  $L$ .

Describe a procedure for constructing an FSA that accepts  $Evenstar(\mathcal{L}(M))$  given an arbitrary DFSA  $M$ . (You may use non-determinism.)

### Solution.

Observe that, by definition,  $Evenstar(L) = (LL)^*$ . We can construct a machine that accepts  $\mathcal{L}(M)\mathcal{L}(M)$  using the construction for concatenation shown in figure 7.18 of the course notes (i.e. we juxtapose two copies of  $M$ , make the accepting states of the first copy non-accepting, and join them to the start state of the second copy by  $\varepsilon$  transitions). We can then apply the construction from figure 7.19 of the course notes (marking the initial state as accepting, and adding epsilon transitions from accepting states to the start state) to *that* machine, to get an NFSA that accepts the Kleene star of the concatenated language, as required.