CSC236 tutorial exercises, Week #8 sample solutions

For each of the algorithms in questions 1-3, prove termination. If proving termination requires a loop invariant, you may state it without proof. (Though you should be confident that your invariant actually holds, and comfortable with how it *could* be proved, if necessary.)

1.

```
1 def ssum(A):
2
       """Pre: A is a list of non-negative ints
      Post: return the sum of A
3
      WARNING: A may be irrerversibly altered!
4
       .....
5
       i = 0
6
7
       s = 0
8
      while i < len(A):
9
          if A[i] == 0:
10
               i += 1
11
           else:
12
               s += 1
13
               A[i] -= 1
14
      return s
```

Solution:

At the end of a given iteration j, let

- Σ_i be the sum of elements in A_j
- m_i be $\Sigma_i + (\operatorname{len}(A_i) i_i)$ (this will be my loop measure)

I claim the following as loop invariants for any iteration j:

- $\Sigma_j \in \mathbb{N}$.
- $i_j \in \mathbb{N}$
- $i_i \leq \operatorname{len}(A)$

It follows that $m_i \in \mathbb{N}$.

Furthermore, the sequence m_0, m_1, \ldots is decreasing. To see why, observe by the code that based on the condition checked on line 9, exactly one of the following happens in each iteration:

• *i* is increased

• Σ is decreased

In either case, m decreases.

```
2.
1 def binsearch(A, x):
       """Pre: A is a sorted list of numbers. x is a number.
2
       Post: return i such that A[i] = x, or -1 if x is
3
             not an element of A.
4
       0.0.0
5
       1o = 0
6
       hi = len(A) - 1
7
8
       while lo <= hi:</pre>
           m = lo + (hi - lo) // 2
9
          mid = A[m]
10
           if mid == x:
11
12
               return m
           elif mid < x:</pre>
13
14
               lo = m + 1
15
           else:
16
                hi = m - 1
17
       return -1
```

Solution:

In this answer I will assume the invariant $hi_j, lo_j \in \mathbb{N}$. I will also assume the invariant $lo_j \leq hi_j + 1$

Lemma 1. For every iteration j > 0, $lo_{j-1} \le m_j \le hi_{j-1}$

Proof. By the loop condition, $lo_{j-1} \leq hi_{j-1}$.

It follows that $(hi_{j-1} - lo_{j-1}) // 2$ is non-negative, which means that $lo_{j-1} \leq m_j$.

To see why the other inequality holds, observe that $hi_{j-1} = lo_{j-1} + (hi_{j-1} - lo_{j-1})$. It follows that $lo_{j-1} + (hi_{j-1} - lo_{j-1}) // 2 \le hi_{j-1}$, since dividing the second term by 2 makes it no larger. \Box

(Note: It would also have been acceptable to simply treat the above lemma as another assumed loop invariant, rather than proving it.)

For a given iteration j, define loop measure $q_j = (hi_j - lo_j) + 1$. It follows from the previous lemma that q_j decreases with each iteration, since (if the loop doesn't exit early), either hi is decreased, or lo is increased. Either action decreases q_j .

That $q_j \in \mathbb{N}$ follows from the invariants I assumed initially.

```
3.
```

```
1 def perambulate(A):
      """Pre: A is a non-empty list of non-negative ints
2
      0.0.0
3
4
      seen = []
5
      curr = A[0]
      i = 0
6
7
      while curr not in seen:
          i = (i + curr) % len(A)
8
          seen.append(curr)
9
```

```
10 curr = A[i]
11 return curr
```

Solution:

For each iteration j, define $m_j = \text{len}(A) - \text{len}(\text{seen}_j)$. It is clear that this decreases with each iteration, since we always increase the length of seen on line 9.

For sake of contradiction, assume there is some iteration j such that $m_j < 0$. I will use the invariant that $seen_j$ contains only elements from A. In this scenario, it follows from the pigeonhole principle that $seen_j$ contains at least one duplicate element x. Consider the iteration k + 1 in which x was last appended to seen. Then $curr_k = x$, and $seen_k$ is a list containing at least one instance of x. But in order for a k + 1th iteration to execute, the loop condition required that $curr_k = x$ was not an element of $seen_k$, a contradiction. Therefore, by contradiction, $m_j \ge 0$ for all j.

Thus m_0, m_1, \ldots is a decreasing sequence of natural numbers, and the loop exits.

Note: Once again, we could have replaced some of the argument above with another assumed loop invariant, such as that $seen_j$ contains no duplicate elements. A proof of this invariant would use the same line of reasoning as the above.

4. Identify the logical flaw in the following "proof" of termination of the function mean.

```
1 def mean(a, b):
2 """Pre: a and b are ints, a < b
3 Post: return the arithmetic mean of a and b
4 """
5 while a != b:
6 a += 1
7 b -= 1
8 return a</pre>
```

Proof of termination Define loop measure $m_j = b_j - a_j$.

 $m_i \in \mathbb{N}$, since a < b, by the precondition.

It remains to show that m decreases with each iteration. For an arbitrary iteration j > 0, $m_j = m_{j-1}-2$ (since we increase a by 1 and decrease b by 1).

Thus $\langle m_0, m_1, \ldots \rangle$ is a decreasing sequence of natural numbers, and therefore finite, so mean terminates.

Solution:

The second line of the proof is not valid. The precondition only applies to a_0 and b_0 , the initial values of a and b. It ensures that $m_0 \in \mathbb{N}$, but when we modify a and b inside the loop, we might reach a situation where $a_j > b_j$. In fact, this happens for any inputs with an odd difference. A call such as mean (2, 3) will loop infinitely.