Chris J. Maddison University of Toronto Daniel Tarlow Microsoft Research Tom Minka Microsoft Research

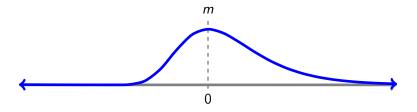
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Goal: Given an unnormalized log density $\phi(x)$, produce independent samples x_1, \ldots, x_n from the Gibbs distribution $p(x) \propto \exp(\phi(x))$.

The Gumbel Distribution

$G \sim \text{Gumbel}(m)$ is Gumbel distributed with location m, if its density is

$$p(g) = \exp(-g + m) \exp(-\exp(-g + m))$$



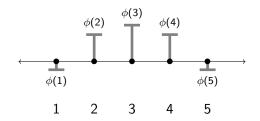
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The Gumbel Distribution

The Gumbel distribution is max-stable. If $G_i \sim \text{Gumbel}(0)$ IID, then

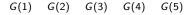
 $\max{G_1, G_2} \sim \text{Gumbel}(\log 2)$

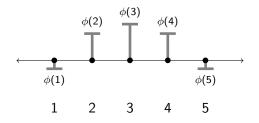
 $p(i) \propto \exp(\phi(i))$ for $i \in \{1, 2, 3, 4, 5\}$



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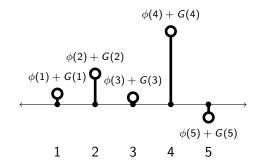
$$p(i) \propto \exp(\phi(i))$$
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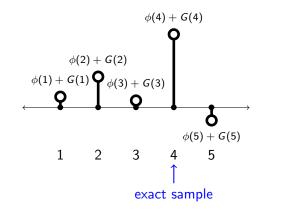
 $G(i) \sim \text{Gumbel}(0) \text{ IID}$

$$p(i) \propto \exp(\phi(i))$$
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$$p(i) \propto \exp(\phi(i))$$
 for $i \in \{1, 2, 3, 4, 5\}$



The Gumbel-Max Trick (well-known, see Yellott 1977)

More formally for any subset B of the indices.

$$\begin{aligned} \operatorname*{argmax}_{i \in B} G(i) + \phi(i) &\sim \frac{\exp(\phi(i))\mathbf{1}(i \in B)}{\sum_{i \in B} \exp(\phi(i))} \\ \max_{i \in B} G(i) + \phi(i) &\sim \operatorname{Gumbel}(\log \sum_{i \in B} \exp(\phi(i))) \end{aligned}$$

What about continuous space?

What about continuous space?

1. Is there an analogous process for perturbing infinite spaces?

2. Can we define practical algorithms for optimizing it?

Now we are interested in

$$p(x) \propto \exp(\phi(x))$$
 for $x \in \mathbb{R}^d$
 $\mu(B) = \int_{x \in B} \exp(\phi(x))$ for $B \subseteq \mathbb{R}^d$

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For this talk just look at \mathbb{R} .

A Quick Re-Frame

We produced a sequence of Gumbels and locations

$$(G(1) + \phi(1), 1)$$
 ... $(G(5) + \phi(5), 5)$

such that

$$\max\{G(i) + \phi(i) \mid i \in B\} \sim \mathsf{Gumbel}(\log \sum_{i \in B} \exp(\phi(i)))$$
$$\arg\max\{G(i) + \phi(i) \mid i \in B\} \sim \frac{\exp(\phi(i))\mathbf{1}(i \in B)}{\sum_{i \in B} \exp(\phi(i))}$$

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By analogy, we want a sequence (G_k, X_k) for $k \to \infty$ such that

$$\max\{G_k \mid X_k \in B\} \sim \mathsf{Gumbel}(\log \mu(B))$$
$$\arg\max\{G_k \mid X_k \in B\} \sim \frac{\exp(\phi(x))\mathbf{1}(x \in B)}{\int_{i \in B} \exp(\phi(x))}$$

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 $\textbf{bottom-up:}\xspace$ instantiate noise \rightarrow find maxes

• Generating infinitely many random variables, *then* finding maxes is a non-starter.

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top-down: pick max \rightarrow generate the rest

• Generate maxes over increasingly refined subsets of space.

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bottom-up: instantiate noise \rightarrow find maxes

• Generating infinitely many random variables, *then* finding maxes is a non-starter.

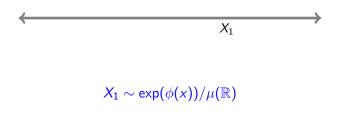
top-down: pick max \rightarrow generate the rest

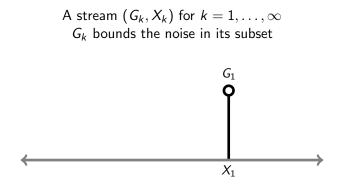
• Generate maxes over increasingly refined subsets of space.

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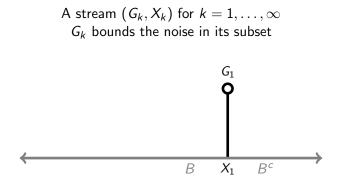
With Gumbel noise, these two directions are equivalent.

A stream (G_k, X_k) for $k = 1, ..., \infty$ G_k bounds the noise in its subset



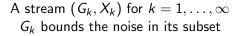


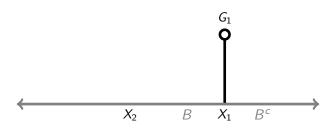
$G_1 \sim \operatorname{Gumbel}(\log \mu(\mathbb{R}))$



split space on X_1

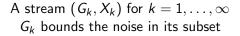
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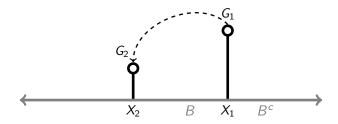




$X_2 \sim \exp(\phi(x))\mathbf{1}(x \in B) / \mu(B)$

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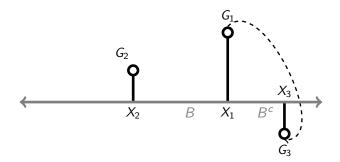




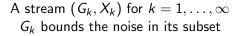
$G_2 \sim \text{TruncGumbel}(\log \mu(B), G_1)$

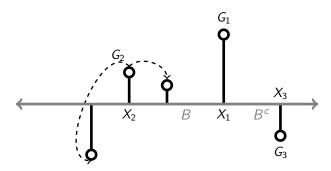
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A stream (G_k, X_k) for $k = 1, ..., \infty$ G_k bounds the noise in its subset



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recursively subdivide space and generate regional maxes

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$$\begin{array}{l} \mathsf{For}\; B \subseteq \mathbb{R} \\ \max\{G_k \,|\, X_k \in B\} \sim \mathsf{Gumbel}(\log \mu(B)) \\ \arg\max\{G_k \,|\, X_k \in B\} \sim \frac{\exp(\phi(x)) \mathbf{1}(x \in B)}{\mu(\mathbb{R})} \end{array}$$

Call $\{\max\{G_k | X_k \in B\} | B \subseteq \mathbb{R}\}$ a *Gumbel Process*.

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- 1. We want to draw independent samples
- 2. We found a process whose optima are samples
- 3. But the procedure for generating it assumes we can draw independent samples

How to practically optimize a Gumbel process without assuming you can tractably sample from p(x) and compute $\mu(B)$.

Like in rejection sampling, decompose $\phi(x)$ into a tractable and boundable component

$$\phi(x)=i(x)+o(x)$$

where for region *B* we can tractably sample and compute volumes from $q(x) \propto \exp(i(x))$ and bound $o(x) \leq M_B$.

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We can also decompose the Gumbel Process

We can take (G_k^q, X_k^q) , a stream of values from the Gumbel process for q(x) and transform it into a realization of a Gumbel process for p(x) by adding o(x).

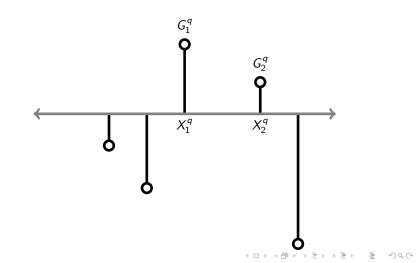
$$G_k^q + o(X_k^q) = G_k$$

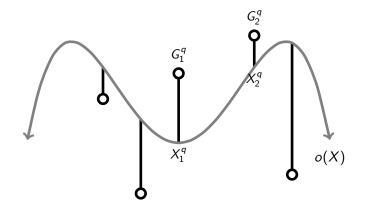
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Take stream
$$(G_k^q, X_k^q)$$
 for $q(x)$, then

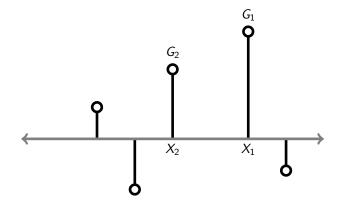
$$\max\{G_k^q + o(X_k^q) | X_k^q \in B\} \sim \text{Gumbel}(\log \mu(B))$$

$$\arg\max\{G_k^q + o(X_k^q) | X_k^q \in B\} \sim \frac{\exp(\phi(x))\mathbf{1}(x \in B)}{\mu(B)}$$





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To draw a sample we want to find

$$\arg\max\{G_k^q + o(X_k^q)\}$$

This decomposition is useful because we can bound

- contribution from the noise of q Gumbel process
- contribution of o(x) this community is good at bounding these functions

$$\max\{G_k^q + o(X_k^q) \,|\, X_k^q \in B\} \le \max\{G_k^q \,|\, X_k^q \in B\} + M_B$$

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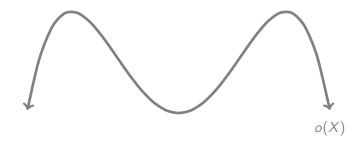
Core Idea: Use A* search to find the optimum.

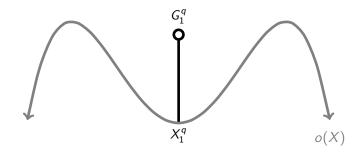
A* Sampling — Ingredients

- The stream of values (G_k^q, X_k^q)
 - G_k^q bounds the noise in its subset.
- Upper bounds on a subset B, $G_k^q + M_B$
- Lower bounds on a subset B, $G_k^q + o(X_k^q)$

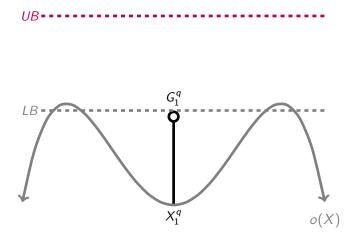
Generally, the two expensive operations are computing M_B and $o(X_k^q)$

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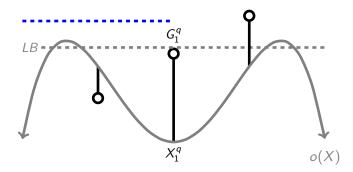




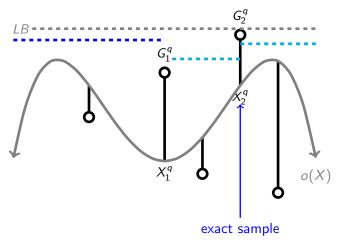
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 G_2^q **A A A** G_1^q χ_2^q X_1^q o(X)



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Come see us at the poster

- Experiments relating A* sampling to other samplers
- Analysis relating A* sampling to adaptive rejection type samplers
 - A* sampling couples which regions are refined and where the sample is more efficient use of bounds and likelihood.

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- Whenever you might sit down to implement slice sampling or rejection sampling for low dimensional non-trivial distributions consider A* sampling.
 - e.g. for the conditionals of a Gibbs sampler
 - In many cases more efficient that alternatives
- We do not solve the problem of high dimensions scales poorly in the worst case.

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• Not surprising, because general & exact.

Conclusions

- Extended the Gumbel-Max trick to continuous spaces.
- Defined A* Sampling, a practical algorithm that optimizes a Gumbel process with A*.

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• Result is new generic sampling algorithm and a useful perspective on the sampling problem.

Acknowledgments

Special thanks to: James Martens Radford Neal Elad Mezuman Roger Grosse







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