

CSC 2125

Homework Operational Semantics

$[\text{ass}_{\text{ns}}]$	$\langle x := a, s \rangle \rightarrow s[x \mapsto \mathcal{A}[a]]s$
$[\text{skip}_{\text{ns}}]$	$\langle \text{skip}, s \rangle \rightarrow s$
$[\text{comp}_{\text{ns}}]$	$\frac{\langle S_1, s \rangle \rightarrow s', \langle S_2, s' \rangle \rightarrow s''}{\langle S_1; S_2, s \rangle \rightarrow s''}$
$[\text{if}_{\text{ns}}^{\text{tt}}]$	$\frac{\langle S_1, s \rangle \rightarrow s'}{\langle \text{if } b \text{ then } S_1 \text{ else } S_2, s \rangle \rightarrow s'} \text{ if } \mathcal{B}[b]s = \text{tt}$
$[\text{if}_{\text{ns}}^{\text{ff}}]$	$\frac{\langle S_2, s \rangle \rightarrow s'}{\langle \text{if } b \text{ then } S_1 \text{ else } S_2, s \rangle \rightarrow s'} \text{ if } \mathcal{B}[b]s = \text{ff}$
$[\text{while}_{\text{ns}}^{\text{tt}}]$	$\frac{\langle S, s \rangle \rightarrow s', \langle \text{while } b \text{ do } S, s' \rangle \rightarrow s''}{\langle \text{while } b \text{ do } S, s \rangle \rightarrow s''} \text{ if } \mathcal{B}[b]s = \text{tt}$
$[\text{while}_{\text{ns}}^{\text{ff}}]$	$\langle \text{while } b \text{ do } S, s \rangle \rightarrow s \text{ if } \mathcal{B}[b]s = \text{ff}$

Table 2.1: Natural semantics for **While**

1. Consider following statement
repeat S until b

- a. Extend the natural operational (“big-step”) semantics of the **WHILE** language (Table 2.1 from [1]) by a rule for relation \rightarrow for the **repeat**-construct. (The semantics for the repeat-construct should not rely on the existence of a while-construct)

- b. Two statements in a natural semantic are considered equivalent if for all states s and s' :

$$\langle S_1, s \rangle \rightarrow s' \text{ iff } \langle S_2, s \rangle \rightarrow s'$$

How can you show that the repeat construct is semantically equivalent to

S; while $\neg b$ do S.

Why does this lead to the conclusion that the extended semantics is deterministic?

$[\text{ass}_{\text{sos}}]$	$\langle x := a, s \rangle \Rightarrow s[x \mapsto \mathcal{A}[a]]s$
$[\text{skip}_{\text{sos}}]$	$\langle \text{skip}, s \rangle \Rightarrow s$
$[\text{comp}_{\text{sos}}^1]$	$\frac{\langle S_1, s \rangle \Rightarrow \langle S'_1, s' \rangle}{\langle S_1; S_2, s \rangle \Rightarrow \langle S'_1; S_2, s' \rangle}$
$[\text{comp}_{\text{sos}}^2]$	$\frac{\langle S_1, s \rangle \Rightarrow s'}{\langle S_1; S_2, s \rangle \Rightarrow \langle S_2, s' \rangle}$
$[\text{if}_{\text{sos}}^{\text{tt}}]$	$\langle \text{if } b \text{ then } S_1 \text{ else } S_2, s \rangle \Rightarrow \langle S_1, s \rangle \text{ if } \mathcal{B}[b]s = \text{tt}$
$[\text{if}_{\text{sos}}^{\text{ff}}]$	$\langle \text{if } b \text{ then } S_1 \text{ else } S_2, s \rangle \Rightarrow \langle S_2, s \rangle \text{ if } \mathcal{B}[b]s = \text{ff}$
$[\text{while}_{\text{sos}}]$	$\langle \text{while } b \text{ do } S, s \rangle \Rightarrow$ $\langle \text{if } b \text{ then } (S; \text{while } b \text{ do } S) \text{ else skip}, s \rangle$

Table 2.2: Structural operational semantics for **While**

2. Consider following statement

repeat S until b

- a. Define the structural operational (“small-step”) semantics as in Table 2.2 from [1] for the **repeat**-construct. (The semantics for the **repeat**-construct should not rely on the existence of a **while**-construct)
 - b. How must the notion of semantic equivalence be defined for structural operational semantics?
3. What distinguishes the two notions of semantic equivalence in 1) and 2)?

[1] Nielson, H., Nielson, F.: “Semantics with Applications: A Formal Introduction”, Wiley Professional Computing, 1992.