Fine-Grained Complexity and Algorithm Design Boot Camp

Recent Advances in FPT and Exact Algorithms for NP-Complete Problems

Dániel Marx

Institute for Computer Science and Control,
Hungarian Academy of Sciences (MTA SZTAKI)
Budapest, Hungary

Simons Institute, Berkeley, CA
September 1, 2015
Overview

- **Today:**
  Introduction to FPT, classical and more recent examples.
  - Definition of FPT.
  - Simple classical examples.
  - Treewidth.
  - Algorithms and applications of treewidth.

- **Wednesday 3pm:**
  Parameterized reductions — negative evidence for FPT.

- **Thursday 3pm:**
  (Tight) lower bounds based on ETH.

- **Friday 3pm:**
  (Even tighter) lower bounds based on SETH.
Parameterized problems

Main idea

Instead of expressing the running time as a function $T(n)$ of $n$, we express it as a function $T(n, k)$ of the input size $n$ and some parameter $k$ of the input.

In other words: we do not want to be efficient on all inputs of size $n$, only for those where $k$ is small.

What can be the parameter $k$?

- The size $k$ of the solution we are looking for.
- The maximum degree of the input graph.
- The dimension of the point set in the input.
- The length of the strings in the input.
- The length of clauses in the input Boolean formula.
- ...
Parameterized problems

Main idea

Instead of expressing the running time as a function $T(n)$ of $n$, we express it as a function $T(n, k)$ of the input size $n$ and some parameter $k$ of the input.

In other words: we do not want to be efficient on all inputs of size $n$, only for those where $k$ is small.

What can be the parameter $k$?

- The size $k$ of the solution we are looking for.
- The maximum degree of the input graph.
- The dimension of the point set in the input.
- The length of the strings in the input.
- The length of clauses in the input Boolean formula.
- ...
Parameterized complexity

**Problem:**

**Input:** Graph $G$, integer $k$

**Question:** Is it possible to cover the edges with $k$ vertices?

**Complexity:** NP-complete

---

**Problem:**

**Input:** Graph $G$, integer $k$

**Question:** Is it possible to find $k$ independent vertices?

**Complexity:** NP-complete
Parameterized complexity

**Problem:**

**Input:**

**Question:**

**Vertex Cover**

Graph $G$, integer $k$

Is it possible to cover the edges with $k$ vertices?

**Independent Set**

Graph $G$, integer $k$

Is it possible to find $k$ independent vertices?

**Complexity:**

NP-complete

**Brute force:**

$O(n^k)$ possibilities

NP-complete

$O(n^k)$ possibilities
Parameterized complexity

**Problem:**

**Input:**

**Question:**

**Complexity:**

**Brute force:**

**Vertex Cover**

Graph $G$, integer $k$

Is it possible to cover the edges with $k$ vertices?

NP-complete

$O(n^k)$ possibilities

$O(2^k n^2)$ algorithm exists


**Independent Set**

Graph $G$, integer $k$

Is it possible to find $k$ independent vertices?

NP-complete

$O(n^k)$ possibilities

No $n^{o(k)}$ algorithm known
Bounded search tree method

Algorithm for **Vertex Cover:**

\[ e_1 = u_1 v_1 \]
Bounded search tree method

**Algorithm for Vertex Cover:**

\[ e_1 = u_1 v_1 \]

\[ \begin{align*} u_1 & \quad v_1 \\ \end{align*} \]
Bounded search tree method

Algorithm for **Vertex Cover**:

\[ e_1 = u_1 v_1 \]

\[ e_2 = u_2 v_2 \]
Bounded search tree method

Algorithm for **Vertex Cover:**

\[ e_1 = u_1 v_1 \]

\[ e_2 = u_2 v_2 \]
Bounded search tree method

Algorithm for **Vertex Cover**:

\[ e_1 = u_1 v_1 \]

\[ e_2 = u_2 v_2 \]

Height of the search tree \( \leq k \) \( \Rightarrow \) at most \( 2^k \) leaves \( \Rightarrow 2^k \cdot n^{O(1)} \) time algorithm.
Fixed-parameter tractability

Main definition

A parameterized problem is **fixed-parameter tractable (FPT)** if there is an $f(k)n^c$ time algorithm for some constant $c$. 

Examples of NP-hard problems that are FPT:
- Finding a vertex cover of size $k$.
- Finding a path of length $k$.
- Finding $k$ disjoint triangles.
- Drawing the graph in the plane with $k$ edge crossings.
- Finding disjoint paths that connect $k$ pairs of points.
A parameterized problem is fixed-parameter tractable (FPT) if there is an \( f(k) n^c \) time algorithm for some constant \( c \).

Examples of NP-hard problems that are FPT:

- Finding a vertex cover of size \( k \).
- Finding a path of length \( k \).
- Finding \( k \) disjoint triangles.
- Drawing the graph in the plane with \( k \) edge crossings.
- Finding disjoint paths that connect \( k \) pairs of points.

...
FPT techniques

Bounded-depth search trees

Kernelization

Algebraic techniques

Treewidth

Color coding

Iterative compression
Parameterized Algorithms

Marek Cygan, Fedor V. Fomin, Łukasz Kowalik, Daniel Lokshtanov, Dániel Marx, Marcin Pilipczuk, Michał Pilipczuk, Saket Saurabh
W[1]-hardness

Negative evidence similar to NP-completeness. If a problem is \textbf{W[1]-hard}, then the problem is not FPT unless FPT=W[1].

Some W[1]-hard problems:

- Finding a clique/independent set of size $k$.
- Finding a dominating set of size $k$.
- Finding $k$ pairwise disjoint sets.
- ... 

More about this on Wednesday at 3pm.
Games to play

- The FPT vs. W[1]-hard game
  Is the problem fixed-parameter tractable?
- The $f(k)$ game for FPT problems
  What is the best $f(k)$ dependence on the parameter?
- The exponent game for W[1]-hard problems
  What is the best possible dependence on $k$ in the exponent?

Significant progress on these questions in recent years, both from the algorithmic and from the complexity side.
Color coding
**Color Coding**

### $k$-Path

**Input:** A graph $G$, integer $k$.

**Find:** A simple path of length $k$.

**Note:** The problem is clearly NP-hard, as it contains the Hamiltonian Path problem.

**Theorem [Alon, Yuster, Zwick 1994]**

$k$-Path can be solved in time $2^{O(k)} \cdot n^{O(1)}$.

Previous best algorithms had running time $k^{O(k)} \cdot n^{O(1)}$. 
Color Coding

- Assign colors from $[k]$ to vertices $V(G)$ uniformly and independently at random.

![Graph Diagram]

Check if there is a path colored $1\rightarrow 2\rightarrow \cdots \rightarrow k$; output "YES" or "NO". If there is no $k$-path: no path colored $1\rightarrow 2\rightarrow \cdots \rightarrow k$ exists. If there is a $k$-path: the probability that such a path is colored $1\rightarrow 2\rightarrow \cdots \rightarrow k$ is $\frac{1}{k^k}$; thus the algorithm outputs "YES" with at least that probability.
Color Coding

- Assign colors from $[k]$ to vertices $V(G)$ uniformly and independently at random.

![Graph Diagram]
Assign colors from \([k]\) to vertices \(V(G)\) uniformly and independently at random.

Check if there is a path colored \(1 \rightarrow 2 \rightarrow \cdots \rightarrow k\); output “YES” or “NO”.

- If there is no \(k\)-path: no path colored \(1 \rightarrow 2 \rightarrow \cdots \rightarrow k\) exists ⇒ “NO”.
- If there is a \(k\)-path: the probability that such a path is colored \(1 \rightarrow 2 \rightarrow \cdots \rightarrow k\) is \(k^{-k}\) thus the algorithm outputs “YES” with at least that probability.
Error probability

**Useful fact**

If the probability of success is at least $p$, then the probability that the algorithm does not say “YES” after $1/p$ repetitions is at most

$$(1 - p)^{1/p} < (e^{-p})^{1/p} = 1/e \approx 0.38$$
Error probability

**Useful fact**
If the probability of success is at least $p$, then the probability that the algorithm does not say “YES” after $1/p$ repetitions is at most

$$(1 - p)^{1/p} < (e^{-p})^{1/p} = 1/e \approx 0.38$$

- Thus if $p > k^{-k}$, then error probability is at most $1/e$ after $k^k$ repetitions.
- Repeating the whole algorithm a constant number of times can make the error probability an arbitrary small constant.
- For example, by trying $100 \cdot k^k$ random colorings, the probability of a wrong answer is at most $1/e^{100}$. 
Finding a path colored $1 - 2 - \cdots - k$

- Edges connecting nonadjacent color classes are removed.
- The remaining edges are directed towards the larger class.
- All we need to check if there is a directed path from class 1 to class $k$. 
Finding a path colored 1 – 2 – ⋅⋅⋅ – k

Edges connecting nonadjacent color classes are removed.
The remaining edges are directed towards the larger class.
All we need to check if there is a directed path from class 1 to class k.
Finding a path colored $1 - 2 - \cdots - k$

- Edges connecting nonadjacent color classes are removed.
- The remaining edges are directed towards the larger class.
- All we need to check if there is a directed path from class 1 to class $k$. 
Finding a path colored $1 - 2 - \cdots - k$

- Edges connecting nonadjacent color classes are removed.
- The remaining edges are directed towards the larger class.
- All we need to check if there is a directed path from class 1 to class $k$. 
Finding a path colored $1 - 2 - \cdots - k$

- Edges connecting nonadjacent color classes are removed.
- The remaining edges are directed towards the larger class.
- All we need to check if there is a directed path from class $1$ to class $k$. 
Color Coding

Finding a 1 - 2 - \cdots - k colored path

polynomial-time solvable

\textit{k-PATH}

Color Coding

success probability:

\( k^{-k} \)
Improved Color Coding

- Assign colors from $[k]$ to vertices $V(G)$ uniformly and independently at random.

- Check if there is a **colorful** path where each color appears exactly once on the vertices; output “YES” or “NO”.

![Graph with colored vertices](image)
Improved Color Coding

Assign colors from $[k]$ to vertices $V(G)$ uniformly and independently at random.

Check if there is a colorful path where each color appears exactly once on the vertices; output “YES” or “NO”.

If there is no $k$-path: no colorful path exists $\Rightarrow$ “NO”.
If there is a $k$-path: the probability that it is colorful is

$$\frac{k!}{k^k} > \frac{\left(\frac{k}{e}\right)^k}{k^k} = e^{-k},$$

thus the algorithm outputs “YES” with at least that probability.
Improved Color Coding

- Assign colors from $[k]$ to vertices $V(G)$ uniformly and independently at random.

![Graph with vertex colors]

- Repeating the algorithm $100e^k$ times decreases the error probability to $e^{-100}$.

How to find a colorful path?
- Try all permutations ($k! \cdot n^{O(1)}$ time)
- Dynamic programming ($2^k \cdot n^{O(1)}$ time)
Finding a colorful path

Subproblems:
We introduce \( 2^k \cdot |V(G)| \) Boolean variables:

\[
x(v, C) = \text{TRUE} \text{ for some } v \in V(G) \text{ and } C \subseteq [k]
\]

\[\uparrow\]

There is a path \( P \) ending at \( v \) such that each color in \( C \) appears on \( P \) exactly once and no other color appears.

Answer:
There is a colorful path \( \iff x(v, [k]) = \text{TRUE} \text{ for some vertex } v \).

Initialization & Recurrence:
Exercise.
Improved Color Coding

$k$-PATH

Color Coding
success probability:

$e^{-k}$

Finding a colorful path

Solvable in time

$2^k \cdot n^{O(1)}$
Derandomized Color Coding

\[ k \]-PATH

\[ k \]-perfect family
\[ 2^{O(k)} \log n \] functions

Finding a colorful path

Solvable in time
\[ 2^k \cdot n^{O(1)} \]
Treewidth
Generalizing trees

How could we define that a graph is “treelike”? 
Generalizing trees

How could we define that a graph is “treelike”?

1. Number of cycles is bounded.

- good
- bad
- bad
- bad
Generalizing trees

How could we define that a graph is “treelike”? 

1. Number of cycles is bounded.
   - good
   - bad
   - bad
   - bad

2. Removing a bounded number of vertices makes it acyclic.
   - good
   - good
   - bad
   - bad
Generalizing trees

How could we define that a graph is “treelike”? 

1. Number of cycles is bounded.

2. Removing a bounded number of vertices makes it acyclic.

The Party Problem

**Party Problem**

**Problem:** Invite some colleagues for a party.

**Maximize:** The total fun factor of the invited people.

**Constraint:** Everyone should be having fun.

Input: At each vertex, weights on the vertices.

Task: Find an independent set of maximum weight.
The Party Problem

**Party Problem**

**Problem:** Invite some colleagues for a party.

**Maximize:** The total fun factor of the invited people.

**Constraint:** Everyone should be having fun.

Do not invite a colleague and his direct boss at the same time!

Input: With weights on the vertices.

Task: Find an independent set of maximum weight.
The Party Problem

**Party Problem**

**Problem:** Invite some colleagues for a party.

**Maximize:** The total fun factor of the invited people.

**Constraint:** Everyone should be having fun.

Do not invite a colleague and his direct boss at the same time!

**Input:** A tree with weights on the vertices.

**Task:** Find an independent set of maximum weight.

![Diagram of a tree with weights on vertices](image)
The Party Problem

**Party Problem**

**Problem:** Invite some colleagues for a party.

**Maximize:** The total fun factor of the invited people.

**Constraint:** Everyone should be having fun. Do not invite a colleague and his direct boss at the same time!

**Input:** A tree with weights on the vertices.

**Task:** Find an independent set of maximum weight.
Dynamic programming paradigm:
We solve a large number of subproblems that depend on each other. The answer is a single subproblem.

Subproblems:
\( T_v \): the subtree rooted at \( v \).

\( A[v] \): max. weight of an independent set in \( T_v \)

\( B[v] \): max. weight of an independent set in \( T_v \) that does not contain \( v \)

Goal: determine \( A[r] \) for the root \( r \).
Solving the Party Problem

Subproblems:
- $T_v$: the subtree rooted at $v$.
- $B[v]$: max. weight of an independent set in $T_v$ that does not contain $v$.

Recurrence:
Assume $v_1, \ldots, v_k$ are the children of $v$. Use the recurrence relations

$$B[v] = \sum_{i=1}^{k} A[v_i]$$
$$A[v] = \max\{B[v], w(v) + \sum_{i=1}^{k} B[v_i]\}$$

The values $A[v]$ and $B[v]$ can be calculated in a bottom-up order (the leaves are trivial).
Treewidth — a measure of “tree-likeness”

**Tree decomposition:** Vertices are arranged in a tree structure satisfying the following properties:

1. If $u$ and $v$ are neighbors, then there is a bag containing both of them.
2. For every $v$, the bags containing $v$ form a connected subtree.

**Width of the decomposition:** largest bag size $-1$.

**Treewidth:** width of the best decomposition.
Treewidth — a measure of “tree-likeness”

**Tree decomposition:** Vertices are arranged in a tree structure satisfying the following properties:

1. If \( u \) and \( v \) are neighbors, then there is a bag containing both of them.
2. For every \( v \), the bags containing \( v \) form a connected subtree.

**Width of the decomposition:** largest bag size \(-1\).

**treewidth:** width of the best decomposition.

A subtree communicates with the outside world only via the root of the subtree.
**Theorem**

Given a tree decomposition of width $w$, **Weighted Max Independent Set** can be solved in time $O(2^w \cdot w^{O(1)} \cdot n)$.

$B_x$: vertices appearing in node $x$.
$V_x$: vertices appearing in the subtree rooted at $x$.

Generalizing our solution for trees:

Instead of computing 2 values $A[v], B[v]$ for each vertex of the tree, we compute $2^{|B_x|} \leq 2^{w+1}$ values for each bag $B_x$.

$M[x, S]$: the max. weight of an independent set $I \subseteq V_x$ with $I \cap B_x = S$. 
Weighted Max Independent Set and treewidth

**Theorem**

Given a tree decomposition of width \( w \), Weighted Max Independent Set can be solved in time \( O(2^w \cdot w^{O(1)} \cdot n) \).

- \( B_x \): vertices appearing in node \( x \).
- \( V_x \): vertices appearing in the subtree rooted at \( x \).

Generalizing our solution for trees:

Instead of computing 2 values \( A[v], B[v] \) for each vertex of the tree, we compute \( 2^{|B_x|} \leq 2^{w+1} \) values for each bag \( B_x \).

**\( M[x, S] \):**

the max. weight of an independent set \( I \subseteq V_x \) with \( I \cap B_x = S \).

How to determine \( M[x, S] \) if all the values are known for the children of \( x \)?
3-Coloring and tree decompositions

Theorem
Given a tree decomposition of width \( w \), 3-Coloring can be solved in time \( 3^w \cdot w^{O(1)} \cdot n \).

\( B_x \): vertices appearing in node \( x \).
\( V_x \): vertices appearing in the subtree rooted at \( x \).

For every node \( x \) and coloring \( c : B_x \rightarrow \{1, 2, 3\} \), we compute the Boolean value \( E[x, c] \), which is true if and only if \( c \) can be extended to a proper 3-coloring of \( V_x \).

Claim:
We can determine \( E[x, c] \) if all the values are known for the children of \( x \).
Tree decompositions and dynamic programming

**General scheme:** Define subproblems for each subtree and solve them in a bottom up manner.

Number of subproblems:

- **3-Coloring:** $3^{w+1}$
  (number of 3-colorings of the bag)
- **Independent Set:** $2^{w+1}$
  (each vertex of the bag is either in the solution or not)
- **Dominating Set:** $3^{w+1}$
  (each vertex of the bag is either (1) in the solution, (2) not in the solution, but dominated, (3) not in the solution and not yet dominated)
- **Hamiltonian Cycle:** $w^{O(w)} = 2^{O(w \log w)}$
  (number of ways the paths of the partial solution can match vertices of the bag).
Number of subproblems for Hamiltonian Cycle

To describe a partial solution, we need to describe the matching of the bag formed by the paths in the partial solution.

Number of matchings: $w^{O(w)} \Rightarrow$ the textbook dynamic programming algorithm has running time $w^{O(w)} \cdot n^{O(1)}$. 

29
To describe a partial solution, we need to describe the matching of the bag formed by the paths in the partial solution.

Number of matchings: $w^{O(w)} \Rightarrow$ the textbook dynamic programming algorithm has running time $w^{O(w)} \cdot n^{O(1)}$. 
Number of subproblems for Hamiltonian Cycle

To describe a partial solution, we need to describe the matching of the bag formed by the paths in the partial solution.

Number of matchings: $w^{O(w)} \Rightarrow$ the textbook dynamic programming algorithm has running time $w^{O(w)} \cdot n^{O(1)}$.

But, surprisingly, it is possible to solve Hamiltonian Cycle in time $2^{O(w)} \cdot n^{O(1)}$!
A very powerful technique for many problems on graphs of bounded-treewidth.

**Classical result:**

**Theorem [textbook algorithm]**

Given a tree decomposition of width $w$, Hamiltonian Cycle can be solved in time $w^{O(w)} \cdot n^{O(1)} = 2^{O(w \log w)} \cdot n^{O(1)}$.

**Improved algorithm:**

**Theorem [Cygan, Nederlof, Pilipczuk, Pilipczuk, van Rooij, Wojtaszczyk 2011]**

Given a tree decomposition of width $w$, Hamiltonian Cycle can be solved in time $4^w \cdot n^{O(1)}$. 
Isolation Lemma

Let $\mathcal{F}$ be a nonempty family of subsets of $U$ and assign a weight $w(u) \in [N]$ to each $u \in U$ uniformly and independently at random. The probability that there is a unique $S \in \mathcal{F}$ having minimum weight is at least

$$1 - \frac{|U|}{N}.$$
Isolation Lemma

Let $F$ be a nonempty family of subsets of $U$ and assign a weight $w(u) \in [N]$ to each $u \in U$ uniformly and independently at random. The probability that there is a unique $S \in F$ having minimum weight is at least

$$1 - \frac{|U|}{N}.$$ 

Let $U = E(G)$ and $F$ be the set of all Hamiltonian cycles.

- By setting $N := |V(G)|^{O(1)}$, we can assume that there is a unique minimum weight Hamiltonian cycle.
- If $N$ is polynomial in the input size, we can guess this minimum weight.
- So we are looking for a Hamiltonian cycle of weight exactly $C$, under the assumption that there is a unique such cycle.
Cycle covers

- **Cycle cover:** A subgraph having degree exactly two at each vertex.
Cycle covers

- **Cycle cover:** A subgraph having degree exactly two at each vertex.

- A Hamiltonian cycle is a cycle cover, but a cycle cover can have more than one component.
Cycle covers

- **Cycle cover**: A subgraph having degree exactly two at each vertex.

- A Hamiltonian cycle is a cycle cover, but a cycle cover can have more than one component.

- **Colored cycle cover**: each component is colored black or white.
Cycle covers

- **Cycle cover**: A subgraph having degree exactly two at each vertex.

![Graphs](image)

- A Hamiltonian cycle is a cycle cover, but a cycle cover can have more than one component.

- **Colored cycle cover**: each component is colored black or white.
Cycle covers

- **Cycle cover**: A subgraph having degree exactly two at each vertex.

- A Hamiltonian cycle is a cycle cover, but a cycle cover can have more than one component.

- **Colored cycle cover**: each component is colored black or white.

- A cycle cover with $k$ components gives rise to $2^k$ colored cycle covers.
Cycle covers

- **Cycle cover**: A subgraph having degree exactly two at each vertex.

- A Hamiltonian cycle is a cycle cover, but a cycle cover can have more than one component.

- **Colored cycle cover**: each component is colored black or white.

- A cycle cover with $k$ components gives rise to $2^k$ colored cycle covers.

  - If there is no weight-$C$ Hamiltonian cycle: the number of weight-$C$ colored cycle covers is $0 \mod 4$.
  - If there is a unique weight-$C$ Hamiltonian cycle: the number of weight-$C$ colored cycle covers is $2 \mod 4$. 
Cycle covers

- **Cycle cover**: A subgraph having degree exactly two at each vertex.

- A Hamiltonian cycle is a cycle cover, but a cycle cover can have more than one component.

- **Colored cycle cover**: each component is colored black or white.

- A cycle cover with $k$ components gives rise to $2^k$ colored cycle covers.

  - If there is no weight-$C$ Hamiltonian cycle: the number of weight-$C$ colored cycle covers is $0 \mod 4$.
  - If there is a unique weight-$C$ Hamiltonian cycle: the number of weight-$C$ colored cycle covers is $2 \mod 4$. 
Cycle covers

- **Cycle cover**: A subgraph having degree exactly two at each vertex.

![Diagram of a cycle cover]

- A Hamiltonian cycle is a cycle cover, but a cycle cover can have more than one component.

- **Colored cycle cover**: each component is colored black or white.

- A cycle cover with $k$ components gives rise to $2^k$ colored cycle covers.

  - If there is no weight-$C$ Hamiltonian cycle: the number of weight-$C$ colored cycle covers is $0 \mod 4$.
  - If there is a unique weight-$C$ Hamiltonian cycle: the number of weight-$C$ colored cycle covers is $2 \mod 4$. 
Assign random weights $\leq 2|E(G)|$ to the edges.

If there is a Hamiltonian cycle, then with probability $1/2$, there is a $C$ such that there is a unique weight-$C$ Hamiltonian cycle.

Try all possible $C$.

Count the number of weight-$C$ colored cycle covers: can be done in time $4^w \cdot n^{O(1)}$ if a tree decomposition of width $w$ is given.

Answer YES if this number is $2 \text{ mod } 4$. 
Cut and Count

HAMILTONIAN CYCLE

Random weights success probability: 1/2

Counting weighted colored cycle covers

$4^k \cdot n^{O(1)}$ time
There are two ways in which we can encounter bounded-treewidth graphs:

1. Designing algorithms for graphs of bounded treewidth.
   - Which problems can be solved efficiently on such graphs?
   - What is the best possible dependence of the running time on treewidth?

2. Using bounded-treewidth algorithms as subroutines.
   - Most notably for planar graphs.
Planar graphs
Subexponential algorithm for $3$-COLORING

**Theorem** [textbook dynamic programming]

$3$-COLORING can be solved in time $2^{O(w)} \cdot n^{O(1)}$ on graphs of treewidth $w$.

+$+$

**Theorem** [Robertson and Seymour]

A planar graph on $n$ vertices has treewidth $O(\sqrt{n})$. 
Subexponential algorithm for 3-COLORING

**Theorem** [textbook dynamic programming]

3-COLORING can be solved in time $2^{O(w)} \cdot n^{O(1)}$ on graphs of treewidth $w$.

+ 

**Theorem** [Robertson and Seymour]

A planar graph on $n$ vertices has treewidth $O(\sqrt{n})$.

⇓

**Corollary**

3-COLORING can be solved in time $2^{O(\sqrt{n})}$ on planar graphs.

Textbook algorithm + combinatorial bound

⇓

Subexponential algorithm
We need only the following basic facts:

**Treewidth**

1. If a graph $G$ has treewidth $w$, then many classical NP-hard problems can be solved in time $2^{O(w)} \cdot n^{O(1)}$ or $2^{O(w \log w)} \cdot n^{O(1)}$ on $G$.

2. A planar graph on $n$ vertices has treewidth $O(\sqrt{n})$.

This immediately gives subexponential-time ($2^{O(\sqrt{n})}$ or $2^{O(\sqrt{n} \log n)}$) algorithms for many problems on planar graphs.

- 3-Coloring
- Hamiltonian Cycle
- Independent Set
- Vertex Cover
- ...
Subexponential planar algorithms using treewidth

We need only the following basic facts:

<table>
<thead>
<tr>
<th>Treewidth</th>
</tr>
</thead>
<tbody>
<tr>
<td>1. If a graph $G$ has treewidth $w$, then many classical NP-hard problems can be solved in time $2^{O(w)} \cdot n^{O(1)}$ or $2^{O(w \log w)} \cdot n^{O(1)}$ on $G$.</td>
</tr>
<tr>
<td>2. A planar graph on $n$ vertices has treewidth $O(\sqrt{n})$.</td>
</tr>
</tbody>
</table>

Next:

What about parameterized problems? Can we make $f(k)$ subexponential for \textsc{Vertex Cover} or \textsc{k-Path} on planar graphs?

But first, let’s see the reason why an $n$-vertex planar graph has treewidth $O(\sqrt{n})$. 
Minors

**Definition**

Graph $H$ is a **minor** of $G$ ($H \leq G$) if $H$ can be obtained from $G$ by deleting edges, deleting vertices, and contracting edges.

![Graph](https://via.placeholder.com/150)

**Note:** length of the longest path in $H$ is at most the length of the longest path in $G$. 
Planar Excluded Grid Theorem

**Theorem** [Robertson, Seymour, Thomas 1994]

Every planar graph with treewidth at least $5k$ has a $k \times k$ grid minor.

*Note:* for general graphs, treewidth at least $k^{100}$ or so guarantees a $k \times k$ grid minor [Chekuri and Chuzhoy 2013]!
Bidimensionality for \( k \)-Path

**Observation:** If the treewidth of a planar graph \( G \) is at least \( 5\sqrt{k} \)

\( \Rightarrow \) It has a \( \sqrt{k} \times \sqrt{k} \) grid minor (Planar Excluded Grid Theorem)

\( \Rightarrow \) The grid has a path of length at least \( k \).

\( \Rightarrow \) \( G \) has a path of length at least \( k \).
Bidimensionality for $k$-Path

**Observation:** If the treewidth of a planar graph $G$ is at least $5\sqrt{k}$

⇒ It has a $\sqrt{k} \times \sqrt{k}$ grid minor (Planar Excluded Grid Theorem)

⇒ The grid has a path of length at least $k$.

⇒ $G$ has a path of length at least $k$.

We use this observation to find a path of length at least $k$ on planar graphs:

- Set $w := 5\sqrt{k}$.
- Find an $O(1)$-approximate tree decomposition.
  - If treewidth is at least $w$: we answer “there is a path of length at least $k$.”
  - If we get a tree decomposition of width $O(w)$, then we can solve the problem in time $2^{O(w \log w)} \cdot n^{O(1)} = 2^{O(\sqrt{k} \log k)} \cdot n^{O(1)}$. 
Bidimensionality

**Definition**

A graph invariant $x(G)$ is **minor-bidimensional** if

- $x(G') \leq x(G)$ for every minor $G'$ of $G$, and
- If $G_k$ is the $k \times k$ grid, then $x(G_k) \geq ck^2$ (for some constant $c > 0$).

**Examples:** minimum vertex cover, length of the longest path, feedback vertex set are minor-bidimensional.
Bidimensionality

**Definition**

A graph invariant $x(G)$ is **minor-bidimensional** if

- $x(G') \leq x(G)$ for every minor $G'$ of $G$, and
- If $G_k$ is the $k \times k$ grid, then $x(G_k) \geq ck^2$ (for some constant $c > 0$).

**Examples:** minimum vertex cover, length of the longest path, feedback vertex set are minor-bidimensional.
Bidimensionality

Definition

A graph invariant $x(G)$ is **minor-bidimensional** if

- $x(G') \leq x(G)$ for every minor $G'$ of $G$, and
- If $G_k$ is the $k \times k$ grid, then $x(G_k) \geq ck^2$ (for some constant $c > 0$).

Examples: minimum vertex cover, length of the longest path, feedback vertex set are minor-bidimensional.
Square root phenomenon for planar graphs

- Simple $2^{O(\sqrt{n})}$ time algorithms for planar graphs by using that planar graphs have treewidth $O(\sqrt{n})$.
- Simple $2^{O(\sqrt{k})} \cdot n^{O(1)}$ time parameterized algorithms using bidimensionality.
- More complicated and problem-specific algorithms for problems where bidimensionality does not work (Steiner Tree, Subset TSP).
- $n^{O(\sqrt{k})}$ time algorithms for W[1]-hard problems.

In many cases, these algorithms are optimal. More about this on Thursday at 3pm...
Wrap up

- The FPT vs. W[1]-hard game
- The $f(k)$ game for FPT problems
- The exponent game for W[1]-hard problems

We have seen that many nontrivial positive results were obtained for these questions.

**Next:** what about negative results?