CSC 2401F 2007, Assignment 2 Due: November 22, 1:10 PM

- 1. Let G = (V, E) be a directed graph and let  $ODDSTCON = \{(G, u, v) | \exists$  a path from u to v and all paths from uto v have odd length}. Show that ODDSTCON is complete for  $\mathcal{NL}$  wrt  $\leq_{log-space}$  by showing:
  - (a) ODDSTCON is in  $\mathcal{NL}$ .
  - (b) ODDSTCON is hard for  $\mathcal{NL}$  wrt  $\leq_{log-space}$ .

10 points

10 points

- 2. Do one of the following
  - Consider regular expressions with a "concatenation squaring operator"  $R^2$  such that  $L[(R)^2] = L[R] \circ L[R]$ . Call such expressions Sregular expressions. Show that  $\{R|R \text{ is a Sregular expression} over \Sigma$  such that  $L[R] \neq \Sigma^*\}$  is hard for exponential space.

10 points

20 points

- Consider extended regular expressions which are regular expressions with the addition of an intersection operator. Show that  $\{R|R \text{ is an extended regular expression over } \Sigma \text{ such that } L[R] \neq \Sigma^*\}$  is hard for exponential space.
- 3. (a) Show that for every k, there is a language L in PH such that L ∉ SIZE(O(n<sup>k</sup>)).
  Hint: by a counting argument show that there exists a language L such that for sufficiently large n, L is in SIZE(n<sup>k+2</sup>) but not in SIZE(n<sup>k+1</sup>) and hence not in SIZE(O(n<sup>k</sup>).
  20 points
  - (b) Using the first part of this question, show that  $P \subseteq SIZE(O(n^k))$  for some fixed k implies  $P \neq NP$ .

10 points