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You were convinced about an Assertion



You were convinced about an Assertion

You didn't learn any other information apart from the fact that the assertion is true.





Verifier V









Verifier V

Statement: *x*

Verifier V

Statement: *x*

Verifier V

Statement: *x*





Verifier V

Statement: *x*

Verifier V

Statement: *x*

Here is a witness *w*



Verifier V

Here is a witness *w*

Statement: *x*

• Completeness: If $x \in \mathcal{L}$, then prover can send a witness that makes the verifier accept.

Interactive Proofs



Verifier V

Here is a witness *w*

Statement: *x*

- verifier accept.

Interactive Proofs

Prover P

• Completeness: If $x \in \mathcal{L}$, then prover can send a witness that makes the

• Soundness: If $x \notin \mathcal{L}$, then whatever string prover sends, the verifier rejects.



Statement: *x*

- verifier accept.

Interactive Proofs

Here is a witness *w*

Prover P

• Completeness: If $x \in \mathcal{L}$, then prover can send a witness that makes the

• Soundness: If $x \notin \mathcal{L}$, then whatever string prover sends, the verifier rejects.

Can we convince the verifier that $x \in \mathscr{L}$ without leaking any information about the witness?

Verifier V

Verifier V

Statement: *x*

Verifier V

Statement: *x*

Prover P

Witness *w*
Zero-Knowledge Proofs

Verifier V

Statement: *x*

Prover P

Zero-Knowledge Proofs

Verifier V

Accepts/rejects

Statement: *x*

Prover P



Accepts/rejects

• Completeness: If $x \in \mathcal{L}$, then verifier accepts at the end of the interaction.

Zero-Knowledge Proofs





Accepts/rejects

Zero-Knowledge Proofs

Statement: *x*



• Completeness: If $x \in \mathcal{L}$, then verifier accepts at the end of the interaction.

• Soundness: If $x \notin \mathcal{L}$, then verifier accepts with probability at most 1/3.



Accepts/rejects

- Soundness: If $x \notin \mathcal{L}$, then verifier accepts with probability at most 1/3.
- Zero-knowledge: The verifier should learn no information about the witness.



• Completeness: If $x \in \mathcal{L}$, then verifier accepts at the end of the interaction.

Statement: *x*

Verifier V

Prover P

Statement: *x*

Verifier V

Witness w

Prover P



Statement: *x*

Prover P



Statement: *x*

Prover P

Verifier V



Verifier V





Verifier V





Verifier V

















3-colorability is NP-complete















Verifier V





Step-1: Prover chooses a random permutation of the colors.



Step-1: Prover chooses a colors.







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Step-2: Prover places each color behind a black-box



Step-3: Verifier chooses a random edge



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Step-4: Prover opens the colors on the endpoints of this edge.



Step-5: Verifier checks if the endpoints are colored differently.





If the graph is 3-colorable, then the verifier always accepts.

Completeness





If the graph is not 3-colorable, then there exists at least one edge that has both endpoints colored using the same color.

Soundness





If the graph is not 3-colorable, then there exists at least one edge that has both endpoints colored using the same color.

Soundness




Verifier V



Zero-Knowledge



Verifier V



Zero-Knowledge