

## Evaluating a predicate formula.

Oct 27

Finally, we are ready to evaluate formulas with free and bound variables.

How do we evaluate a formula  $(\forall x \alpha)$  or  $(\exists x \alpha)$ ?

- For  $(\forall x \alpha)$ , we need to verify that  $\alpha$  is true for every possible value of  $x$  in the domain.
- For  $(\exists x \alpha)$ , we need to verify that  $\alpha$  is true for at least one value of  $x$  in the domain.

Formally, the values of  $(\forall x \alpha)$  and  $(\exists x \alpha)$  are.

- $(\forall x \alpha)^{(I, E)} = T$  if  $\alpha^{(I, E[x \mapsto d])} = T$  for every  $d \in \text{dom}(I)$ .  
and it is F otherwise.
- $(\exists x \alpha)^{(I, E)} = T$  if  $\alpha^{(I, E[x \mapsto d])} = T$  for at least one  $d \in \text{dom}(I)$ .  
and it is F otherwise.

$E[x \mapsto d]$  is a new environment created by making a change to  $E$ .

$E[x \mapsto d]$  means that "keep all the mappings in  $E$  intact EXCEPT re-assign  $x$  to  $d \in \text{dom}(I)$ ".

Example: Consider  $E$  defined below.

$$E(x) = 3, \quad E(y) = 3, \quad E(z) = 1. \quad \text{dom}(I) = \{1, 2, 3\}.$$

- ①  $E[x \mapsto 2](x) = 2$      $E[x \mapsto 2](y) = 3$      $E[x \mapsto 2](z) = 1$ .
- ②  $E[x \mapsto 2][y \mapsto 2](x) = 2$      $E[x \mapsto 2][y \mapsto 2](y) = 2$   
 $E[x \mapsto 2][y \mapsto 2](z) = 1$

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Our language: Variable symbols:  $x, y$   
Predicate symbols:  $Q^{(2)}$

Interpretation  $I$ :  $\text{dom}(I) = \{1, 2\}$   $Q^I = \emptyset$   
Environment  $E$ :  $E(x) = 1$   $E(y) = 2$

Examples:

- ① Give an interpretation  $I_1$  and an environment  $E_1$  such that  $I_1 \models_{E_1} \alpha$  where  $\alpha = (\exists x Q(x, y))$ .
- $y$  is free in  $\alpha$ , so its value is given by  $E_1$ .
  - To make  $\alpha$  true, we need there to be at least one tuple in  $Q^{I_1}$  and the second value in the tuple must be 2 because  $E_1(y) = 2$ .

Let  $I_1 = I$  and  $E_1 = E$  except that  $Q^{I_1} = \{ \langle 1, 2 \rangle \}$ .

$$Q(x, y) \stackrel{(I_1, E_1[x \mapsto 1])}{=} T \text{ because } \langle 1, 2 \rangle \in Q^{I_1}.$$

Hence,  $(\exists x Q(x, y)) \stackrel{(I_1, E_1)}{=} T$  and  $I_1 \models_{E_1} \alpha$ .

- ② Give  $I_2$  and  $E_2$  such that  $I_2 \models_{E_2} \alpha$  where  $\alpha = (\forall x Q(x, y))$ .
- $y$  is free in  $\alpha$ , so its value is given by  $E_2$ .
  - To make  $\alpha$  true, we need 2 tuples in  $Q^{I_2}$ , each starting with a possible value for  $x$  and ending with 2 because  $E_2(y) = 2$ .

Let  $I_2 = I$  and  $E_2 = E$  except that  $Q^{I_2} = \{ \langle 1, 2 \rangle, \langle 2, 2 \rangle \}$

$$\text{Case } [x \mapsto 1] \quad Q(x, y) \stackrel{(I_2, E_2[x \mapsto 1])}{=} T \text{ because } \langle 1, 2 \rangle \in Q^{I_2}.$$

$$\text{Case } [x \mapsto 2] \quad Q(x, y) \stackrel{(I_2, E_2[x \mapsto 2])}{=} T \text{ because } \langle 2, 2 \rangle \in Q^{I_2}.$$

Therefore,  $(\forall x Q(x, y)) \stackrel{(I_2, E_2)}{=} T$  and  $I_2 \models_{E_2} \alpha$ .

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Environment  $E$ :  $E(x) = 1$   $E(y) = 2$ .

Examples:

③ Give  $I_3$  and  $E_3$  such that  $I_3 \models_{E_3} \alpha$  where  $\alpha = (\exists x (\forall y Q(x, y)))$ .

Start with  $I_3 = I$  and  $E_3 = E$ .

(To make  $\alpha$  true, we need at least two tuples in  $Q^{I_3}$ . Both tuples need to start with the same value for  $x$ . The second values of the tuples need to enumerate all possible domain values for  $y$ .)

Let  $Q^{I_3} = \{ \langle 1, 1 \rangle, \langle 1, 2 \rangle \}$ .

To show that  $\alpha^{(I_3, E_3)} = T$ , it is sufficient to show that

$$(\forall y Q(x, y))^{(I_3, E_3[x \mapsto 1])} = T.$$

Consider all possible values for  $y$ :

$$\text{Case } [y \mapsto 1]: Q(x, y)^{(I_3, E_3[x \mapsto 1][y \mapsto 1])} = T.$$

$$\text{Case } [y \mapsto 2]: Q(x, y)^{(I_3, E_3[x \mapsto 1][y \mapsto 2])} = T.$$

Therefore,  $(\forall y Q(x, y))^{(I_3, E_3[x \mapsto 1])} = T$  and  $I_3 \models_{E_3} \alpha$ .

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Our language: variable symbols:  $x, y$  predicate symbols:  $Q^{(2)}$

Interpretation  $I$ :  $\text{dom}(I) = \{1, 2\}$   $Q^I = \emptyset$ .

Environment  $E$ :  $E(x) = 1$ ,  $E(y) = 2$ .

Examples:

④ Show that  $I \not\models_E \alpha$  where  $\alpha = (\exists x (\forall y Q(x, y)))$ .

We need to show that  $(\exists x (\forall y Q(x, y)))^{(I, E)} = F$ , which is equivalent to

$$(\neg (\exists x (\forall y Q(x, y))))^{(I, E)} = T.$$

$$\text{or } (\forall x (\exists y (\neg Q(x, y))))^{(I, E)} = T.$$

Consider all possible values for  $x$ :

Case  $[x \mapsto 1]$ :  $Q(x, y)^{(I, E[x \mapsto 1][y \mapsto 1])} = F$  because  $\langle 1, 1 \rangle \notin Q^I$ .

Thus,  $(\exists y (\neg Q(x, y)))^{(I, E[x \mapsto 1])} = T$ .

Case  $[x \mapsto 2]$ :  $Q(x, y)^{(I, E[x \mapsto 2][y \mapsto 1])} = F$  because  $\langle 2, 1 \rangle \notin Q^I$ .

Thus  $(\exists y (\neg Q(x, y)))^{(I, E[x \mapsto 2])} = T$ .

Therefore,  $(\forall x (\exists y (\neg Q(x, y))))^{(I, E)} = T$  and  $I \not\models_E \alpha$ .