# Predicate Logic: Syntax 

Alice Gao

Lecture 12

## Outline

## Learning goals

Symbols

Terms

Formulas

## Parse Trees

Revisiting the learning goals

## Learning goals

By the end of this lecture, you should be able to

- Define the set of terms inductively.
- Define the set of formulas inductively.
- Determine whether a variable in a formula is free or bound.
- Prove properties of terms and formulas by structural induction.
- Draw the parse tree of a formula.


## The Language of Predicate Logic

- Domain: a non-empty set of objects.
- Individuals: concrete objects in the domain.
- Variables: placeholders for concrete objects in the domain.
- Functions: takes objects in the domain as arguments and returns an object of the domain.
- Relations: takes objects in the domain as arguments and returns true or false. They describe properties of objects or relationships between objects.
- Quantifiers: for how many objects in the domain is the statement true?


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## Predicate Language $L$

Eight classes of symbols:

- Individual symbols: $a, b, c$.
- Relation symbols: $F, G, H$.

A special equality symbol $\approx$

- Function symbols: $f, g, h$.
- Free variable symbols: $u, v, w$.
- Bound variable symbols: $x, y, z$.
- Connective symbols: $\neg, \wedge, \vee, \rightarrow, \leftrightarrow$.
- Quantifier symbol: $\forall, \exists$.
- Punctuation symbols: '(', ')', and ','


## Free and Bound Variables

In a formula $\forall x A(x)$ or $\exists x A(x)$, the scope of a quantifier is the formula $A(x)$.

A quantifier binds its variable within its scope.
An occurrence of a variable in a formula

- is bound if it lies in the scope of some quantifier of the same variable.
- is free, otherwise.


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## Two Kinds of Expressions

Two kinds of expressions:

- A term refers to an object in the domain.
- A formula evaluates to 1 or 0 .


## Terms

The set of terms Term $(L)$ is defined below:

1. An individual symbol $a$ standing alone is a term.
2. A free variable symbol $u$ standing alone is a term.
3. If $t_{1}, \ldots, t_{n}$ are terms and $f$ is an $n$-ary function symbol, then $f\left(t_{1}, \ldots, t_{n}\right)$ is a term.
4. Nothing else is a term.

## Examples of Terms

Terms:

- $a, b, c, u, v, w$
-f(b),g(a,f(b)),g(u,b),f(g(f(u),b))) (1)
A term with no free variable symbols is called a closed term. Which one(s) of the above are closed terms?


## CQ: Which expressions are terms?

Which of the following expressions is a term?
If there are multiple correct answers, choose your favourite one.
(A) $w$
(B) $g(a, u)$
(C) $F(f(u, v), a)$
(D) $f(u, g(v, w), a)$
(E) $g(u, f(v, w), a)$

Individual symbols: $a$
Relation symbols: $F$ is a binary relation symbol.
Function symbols: $f$ is a binary function symbol and $g$ is a 3-ary function symbol.
Free variable symbols: $u, v, w$.

## Defining the set of terms inductively

The set of terms can be inductively defined as follows:

- The domain set $X$ :
- The core set $C$ :
- The set of operations $P$ :


## Structural induction on terms

Theorem: Every term has a property $P$.
Proof by structural induction:

- Base cases:
- Inductive cases:


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## Atomic Formulas

The set of atomic formulas $\operatorname{Atom}(L)$ is defined below:

- If $F$ is an n -ary relation symbol and $t_{1}, \ldots, t_{n}(n \geq 1)$ are terms, then $F\left(t_{1}, \ldots, t_{n}\right)$ is an atomic formula.
- If $t_{1}, t_{2}$ are terms, then $\approx\left(t_{1}, t_{2}\right)$ is an atomic formula.
- Nothing else is an atomic formula.


## Examples of Atomic Formulas

Terms:
$>a, b, c, u, v, w$

- $f(b), g(a, f(b)), g(u, b), f(g(f(u), b))$

Atomic formulas:

- $F(a, u, f(b), f(w), g(v, f(a)))$
- $\approx(b, w)$


## Well-Formed Formulas

The set of well-formed formulas Form $(L)$ is defined below:

1. An atomic formula is a well-formed formula.
2. If $A$ is a well-formed formula, then $(\neg A)$ is a well-formed formula.
3. If $A$ and $B$ are well-formed formulas and $\star$ is one of $\wedge, \vee, \rightarrow$, and $\leftrightarrow$, then $(A \star B)$ is a well-formed formula.
4. If $A(u)$ is a well-formed formula and $x$ does not occur in $A(u)$, then $\forall x A(x)$ and $\exists x A(x)$ are well-formed formulas.
5. Nothing else is a well-formed formula.

## Explaining Case 4 of Formulas

If $A(u)$ is a well-formed formula and $x$ does not occur in $A(u)$, then $\forall x A(x)$ and $\exists x A(x)$ are well-formed formulas.

- $A(u)$ is a well-formed formula where $u$ is a free variable in the formula. We want to quantify $u$.
- In order to do so, we need to choose a symbol for a bound variable, e.g. $x$. We need to make sure that our choice of the bound variable symbol does not already occur in $A(u)$.


## Examples for Case 4

- We are allowed to generate the formula $\forall y F(y, y)$. Start with $F(u, u)$. If we quantify $u$ by replacing it with $y$, we get $\forall y F(y, y)$.
- We are not allowed to generate the formula $\exists y \forall y F(y, y)$. Start with $\forall y F(y, y)$. If we want to add the $\exists$ quantifier, we will need to choose a bound variable symbol that is not $y$ because $y$ already appears in $\forall y F(y, y)$. So, there is no way for us to generate $\exists y \forall y F(y, y)$.
- We are allowed to generate the formula $\exists x G(x) \vee \forall x H(x)$. Start with $G(u)$ and $H(v)$ separately. We can quantify $u$ by replacing it with $x$ since $x$ does not appear in $G(u)$ ). We get $\exists x G(x)$. We can quantify $v$ by replacing it with $x$ since $x$ does not appear in $H(v)$. We get $\forall x H(x)$. Connecting the two formulas using $\vee$, we get $\exists x G(x) \vee \forall x H(x)$.


## Examples of Formulas

Well-Formed Formulas:

- $F(a, b), \forall y F(a, y), \exists x \forall y F(x, y)$
- $F(u, v), \exists y F(u, y)$

A formula with no free variable symbols is called a closed formula or a sentence.
Which formulas above are closed formulas?

## Determine whether a formula is well-formed

Which of the following is a well-formed formula?
(A) $f(u) \rightarrow F(u, v)$
(B) $\forall x F(m, f(x))$
(C) $F(u, v) \rightarrow G(G(u))$
(D) $G(m, f(m))$
(E) $F(m, f(G(u, v)))$

Individual symbols: $m$.
Free Variable Symbols: $u, v$.
Bound Variable symbols: $x$.
Relation symbols: $F$ and $G$ are binary relation symbols.
Function symbols: $f$ is a unary function.

## Defining the set of formulas inductively

The set of formulas can be inductively defined as follows:

- The domain set $X$ :
- The core set $C$ :
- The set of operations $P$ :


## Structural induction on formulas

Theorem: Every formula has a property $P$.
Proof by structural induction:

- Base cases:
- Inductive cases:


## Comparing the Definitions of Well-Formed Formulas

Let's compare the set of predicate formulas to the set of propositional formulas.

Questions to think about:

- Which parts of the two definitions are the same?
- Which parts of the two definitions are different?


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## Parse Trees of Predicate Formulas

- The leaves are atomic formulas.
- Every quantifier has exactly one child (namely the formula which is its scope).
Example: $\forall x(F(b) \rightarrow \exists y(\forall z G(y, z) \vee H(f(u), x, y)))$


## Revisiting the learning goals

By the end of this lecture, you should be able to

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