# Predicate Logic: Syntax

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Lecture 12

## Outline

Learning goals

**Symbols** 

**Terms** 

**Formulas** 

Parse Trees

Revisiting the learning goals

# Learning goals

By the end of this lecture, you should be able to

- ▶ Define the set of terms inductively.
- ▶ Define the set of formulas inductively.
- ▶ Determine whether a variable in a formula is free or bound.
- Prove properties of terms and formulas by structural induction.
- ▶ Draw the parse tree of a formula.

# The Language of Predicate Logic

- Domain: a non-empty set of objects.
- Individuals: concrete objects in the domain.
- ▶ Variables: placeholders for concrete objects in the domain.
- ► Functions: takes objects in the domain as arguments and returns an object of the domain.
- Relations: takes objects in the domain as arguments and returns true or false. They describe properties of objects or relationships between objects.
- Quantifiers: for how many objects in the domain is the statement true?

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# Predicate Language L

#### Eight classes of symbols:

- ▶ Individual symbols: a, b, c.
- ▶ Relation symbols: F, G, H. A special equality symbol  $\approx$
- ▶ Function symbols: f, g, h.
- Free variable symbols: u, v, w.
- **b** Bound variable symbols: x, y, z.
- ▶ Connective symbols:  $\neg$ ,  $\land$ ,  $\lor$ ,  $\rightarrow$ ,  $\leftrightarrow$ .
- ▶ Quantifier symbol:  $\forall$ ,  $\exists$ .
- ▶ Punctuation symbols: '(', ')', and ','

## Free and Bound Variables

In a formula  $\forall x \ A(x)$  or  $\exists x \ A(x)$ , the scope of a quantifier is the formula A(x).

A quantifier binds its variable within its scope.

An occurrence of a variable in a formula

- is bound if it lies in the scope of some quantifier of the same variable.
- is free, otherwise.

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# Two Kinds of Expressions

#### Two kinds of expressions:

- ▶ A term refers to an object in the domain.
- ▶ A formula evaluates to 1 or 0.

#### **Terms**

## The set of terms Term(L) is defined below:

- 1. An individual symbol a standing alone is a term.
- 2. A free variable symbol u standing alone is a term.
- 3. If  $t_1,\ldots,t_n$  are terms and f is an n-ary function symbol, then  $f(t_1,\ldots,t_n)$  is a term.
- 4. Nothing else is a term.

## **Examples of Terms**

#### Terms:

- $\triangleright a, b, c, u, v, w$
- ightharpoonup f(b), g(a, f(b)), g(u, b), f(g(f(u), b))

A term with no free variable symbols is called a closed term. Which one(s) of the above are closed terms?

# CQ: Which expressions are terms?

Which of the following expressions is a term? If there are multiple correct answers, choose your favourite one.

- (A) w
- (B) g(a, u)
- (C) F(f(u,v),a)
- (D) f(u, g(v, w), a)
- (E) g(u, f(v, w), a)

Individual symbols: a

Relation symbols: F is a binary relation symbol.

Function symbols: f is a binary function symbol and g is a 3-ary

function symbol.

Free variable symbols: u, v, w.

# Defining the set of terms inductively

The set of terms can be inductively defined as follows:

- ightharpoonup The domain set X:
- ▶ The core set *C*:
- ▶ The set of operations *P*:

#### Structural induction on terms

Theorem: Every term has a property P.

Proof by structural induction:

- ▶ Base cases:
- Inductive cases:

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#### Atomic Formulas

The set of atomic formulas Atom(L) is defined below:

- ▶ If F is an n-ary relation symbol and  $t_1, \ldots, t_n$   $(n \ge 1)$  are terms, then  $F(t_1, \ldots, t_n)$  is an atomic formula.
- $\blacktriangleright$  If  $t_1,t_2$  are terms, then  $\approx (t_1,t_2)$  is an atomic formula.
- Nothing else is an atomic formula.

# **Examples of Atomic Formulas**

#### Terms:

- $\triangleright$  a, b, c, u, v, w
- $\blacktriangleright \ f(b)\text{, } g(a,f(b))\text{, } g(u,b)\text{, } f(g(f(u),b))$

#### Atomic formulas:

- ightharpoonup F(a, u, f(b), f(w), g(v, f(a)))
- ightharpoonup pprox (b,w)

## Well-Formed Formulas

The set of well-formed formulas Form(L) is defined below:

- 1. An atomic formula is a well-formed formula.
- 2. If A is a well-formed formula, then  $(\neg A)$  is a well-formed formula.
- 3. If A and B are well-formed formulas and  $\star$  is one of  $\land$ ,  $\lor$ ,  $\rightarrow$ , and  $\leftrightarrow$ , then  $(A \star B)$  is a well-formed formula.
- 4. If A(u) is a well-formed formula and x does not occur in A(u), then  $\forall x\,A(x)$  and  $\exists x\,A(x)$  are well-formed formulas.
- 5. Nothing else is a well-formed formula.

# Explaining Case 4 of Formulas

If A(u) is a well-formed formula and x does not occur in A(u), then  $\forall x\,A(x)$  and  $\exists x\,A(x)$  are well-formed formulas.

- ightharpoonup A(u) is a well-formed formula where u is a free variable in the formula. We want to quantify u.
- ▶ In order to do so, we need to choose a symbol for a bound variable, e.g. x. We need to make sure that our choice of the bound variable symbol does not already occur in A(u).

## Examples for Case 4

- ▶ We are allowed to generate the formula  $\forall y F(y,y)$ . Start with F(u,u). If we quantify u by replacing it with y, we get  $\forall y F(y,y)$ .
- ▶ We are not allowed to generate the formula  $\exists y \forall y F(y,y)$ . Start with  $\forall y F(y,y)$ . If we want to add the  $\exists$  quantifier, we will need to choose a bound variable symbol that is not y because y already appears in  $\forall y F(y,y)$ . So, there is no way for us to generate  $\exists y \forall y F(y,y)$ .
- ▶ We are allowed to generate the formula  $\exists xG(x) \lor \forall xH(x)$ . Start with G(u) and H(v) separately. We can quantify u by replacing it with x since x does not appear in G(u). We get  $\exists xG(x)$ . We can quantify v by replacing it with x since x does not appear in H(v). We get  $\forall xH(x)$ . Connecting the two formulas using  $\lor$ , we get  $\exists xG(x) \lor \forall xH(x)$ .

## **Examples of Formulas**

#### Well-Formed Formulas:

- $\blacktriangleright$  F(a,b),  $\forall y F(a,y)$ ,  $\exists x \forall y F(x,y)$
- ightharpoonup F(u,v),  $\exists y F(u,y)$

A formula with no free variable symbols is called a closed formula or a sentence.

Which formulas above are closed formulas?

## Determine whether a formula is well-formed

Which of the following is a well-formed formula?

- (A)  $f(u) \rightarrow F(u, v)$
- (B)  $\forall x \ F(m, f(x))$
- (C)  $F(u,v) \rightarrow G(G(u))$
- (D) G(m, f(m))
- (E) F(m, f(G(u, v)))

Individual symbols: m.

Free Variable Symbols: u, v.

Bound Variable symbols: x.

Relation symbols: F and G are binary relation symbols.

Function symbols: f is a unary function.

# Defining the set of formulas inductively

The set of formulas can be inductively defined as follows:

- ightharpoonup The domain set X:
- ▶ The core set *C*:
- ▶ The set of operations *P*:

## Structural induction on formulas

Theorem: Every formula has a property P.

Proof by structural induction:

- ▶ Base cases:
- Inductive cases:

# Comparing the Definitions of Well-Formed Formulas

Let's compare the set of predicate formulas to the set of propositional formulas.

Questions to think about:

- ▶ Which parts of the two definitions are the same?
- ▶ Which parts of the two definitions are different?

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## Parse Trees of Predicate Formulas

- ▶ The leaves are atomic formulas.
- Every quantifier has exactly one child (namely the formula which is its scope).

Example:  $\forall x (F(b) \rightarrow \exists y (\forall z G(y, z) \lor H(f(u), x, y)))$ 

# Revisiting the learning goals

By the end of this lecture, you should be able to

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