

# **Social and Information Networks**

**CSCC46H, Fall 2019  
Lecture 4**

Prof. Ashton Anderson  
[ashton@cs.toronto.edu](mailto:ashton@cs.toronto.edu)



# Logistics

**A1 due next Monday @ 12pm on MarkUs**

**First letter of last name A-H? First blog post due this Friday**

# Today

**Signed networks**

**Empirical phenomena in networks**

# Positive and Negative Relationships

**So far, edges mostly interpreted positively**

- Friendship
- Interaction
- Collaboration

But relationships can be **negative** too

- Dislike
- Bad interaction
- Enemy

# Network Representation

**How would you model this?**

# Signed Networks

## Networks with **positive** and **negative** relationships

Consider an **undirected complete graph**

Label each edge as either:

**Positive:** friendship, trust, positive sentiment, ...

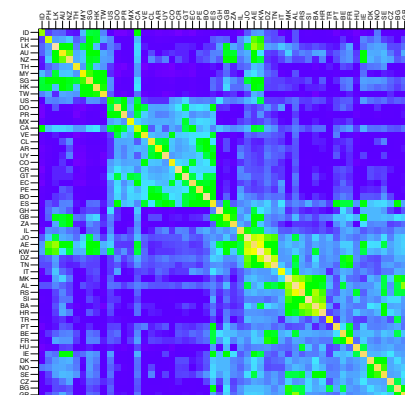
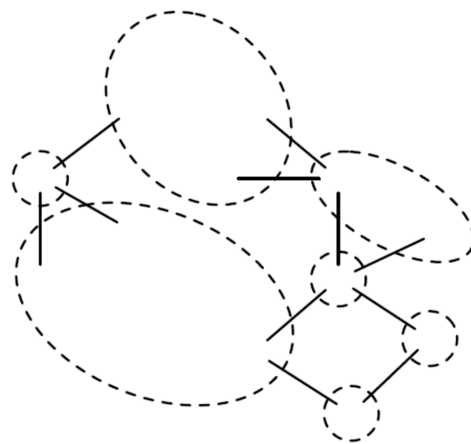
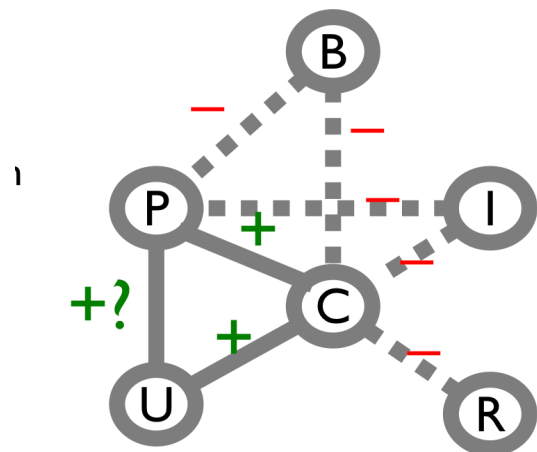
**Negative:** enemy, distrust, negative sentiment, ...

# Questions about Signed Networks

What are the **typical** patterns of interaction in signed networks?

How do we **reason** about **local and global structure** of positive and negative interactions?

What are the **patterns in empirical data**?

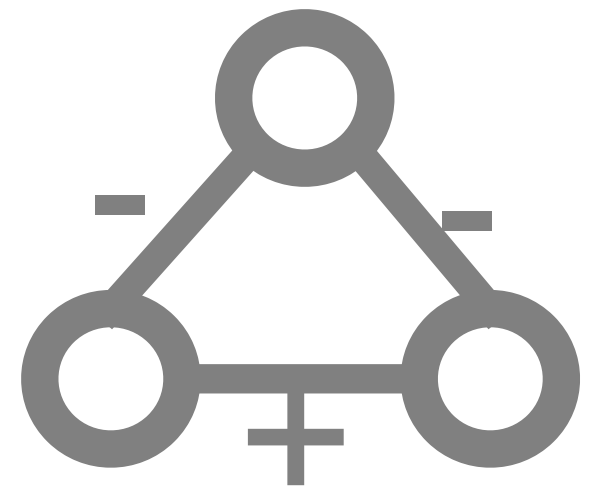


# Signed Networks

**Networks with positive and negative relationships**

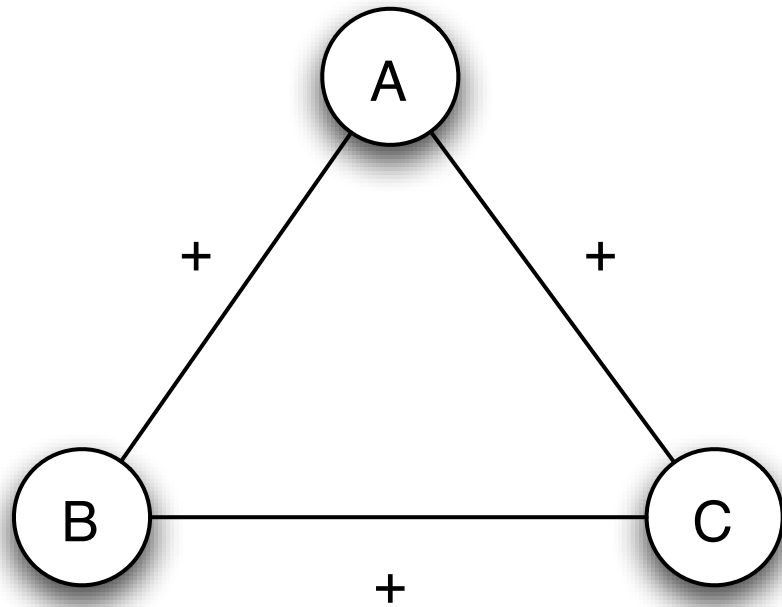
Our basic unit of investigation will be **signed triangles**

Focus on **undirected** networks

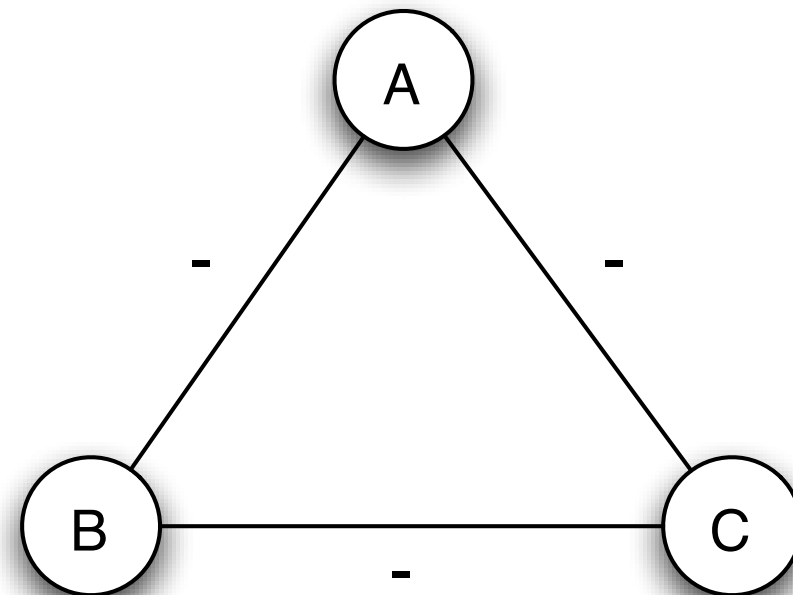
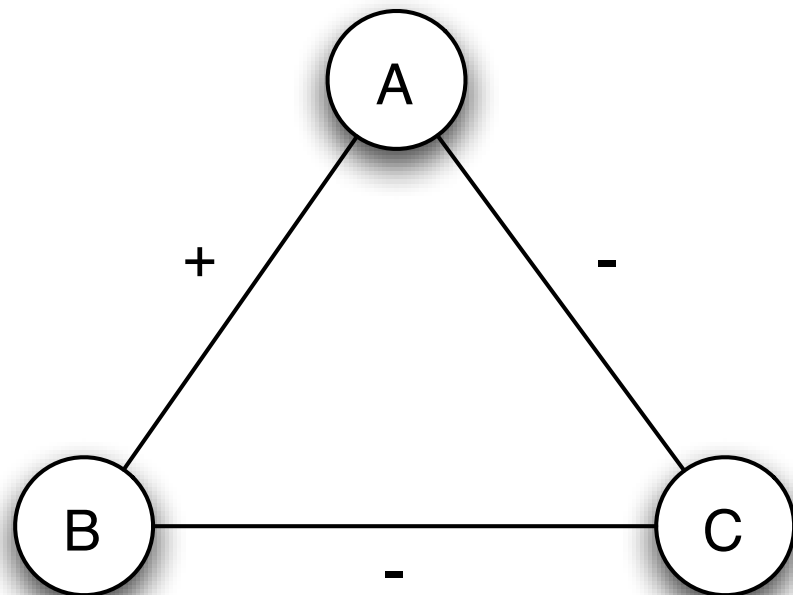
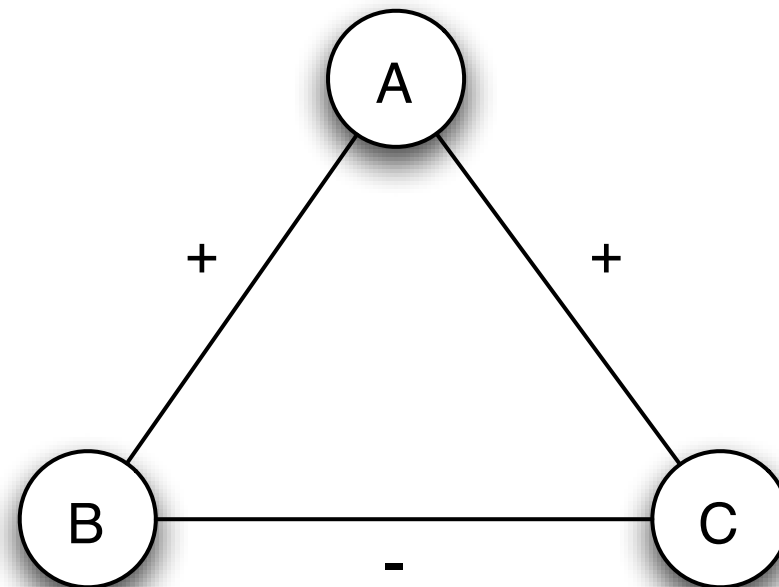
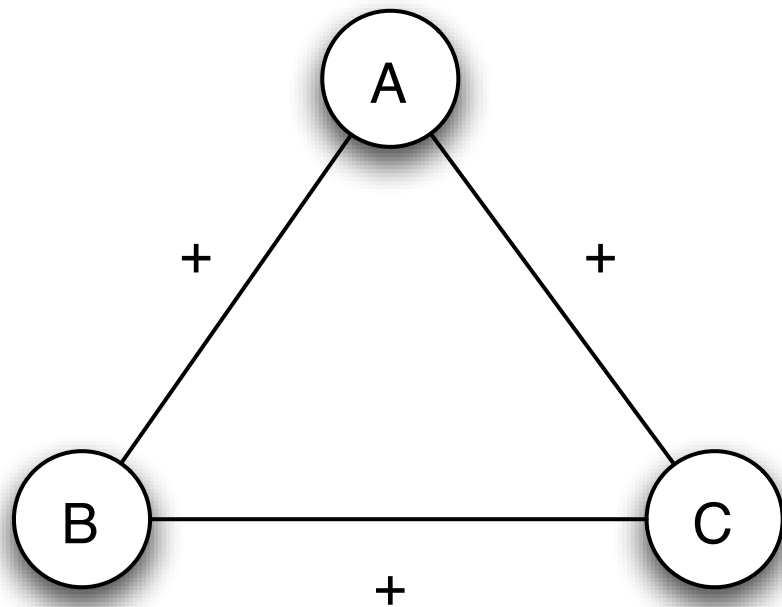




# Structural Balance

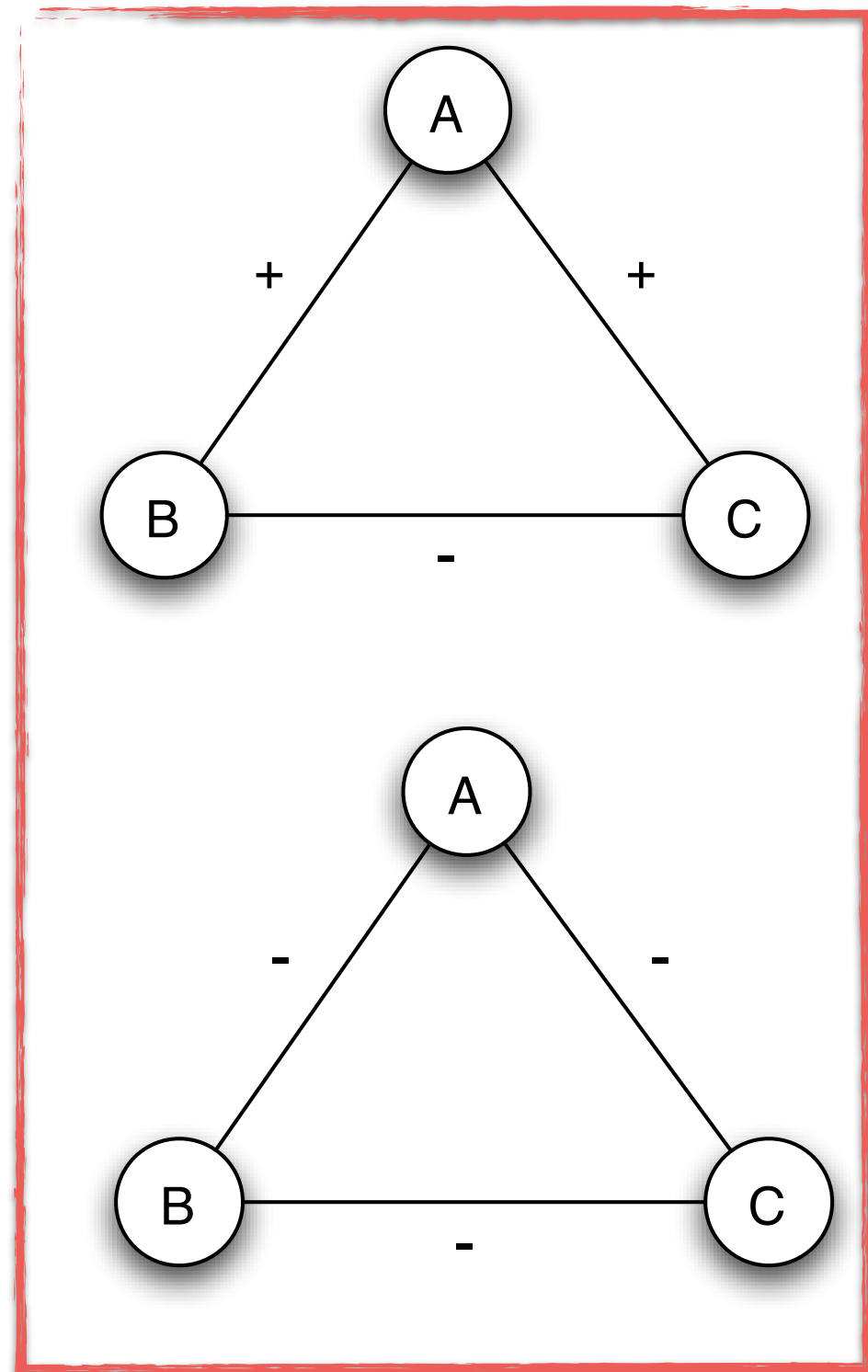
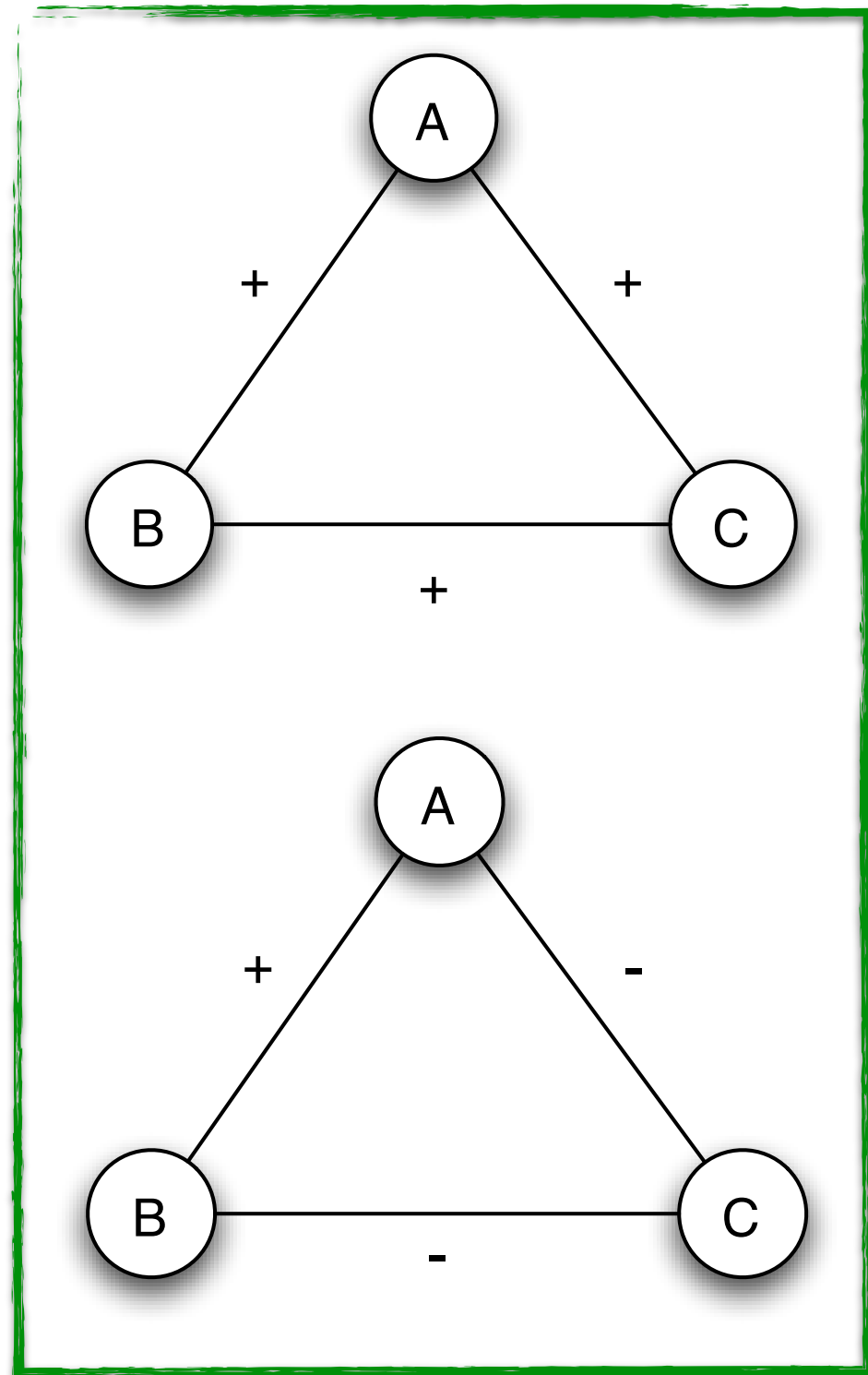


# Structural Balance



Four signed triads: which are **stable**?

# Structural Balance



Four signed triads: which are **stable**?

# Theory of Structural Balance

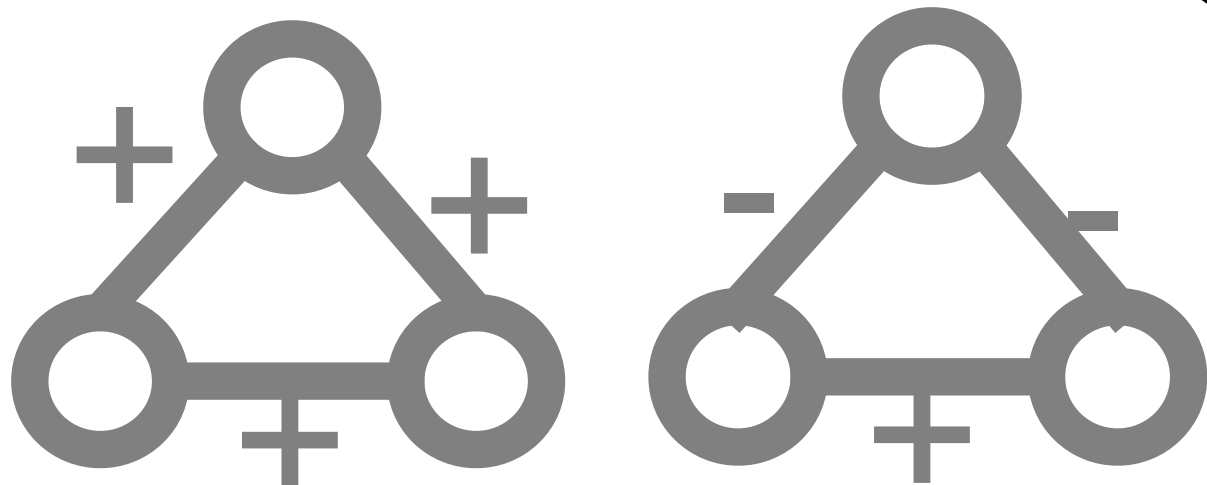
**Start with the intuition** [Heider '46]:

Friend of my friend is my friend

Enemy of enemy is my friend

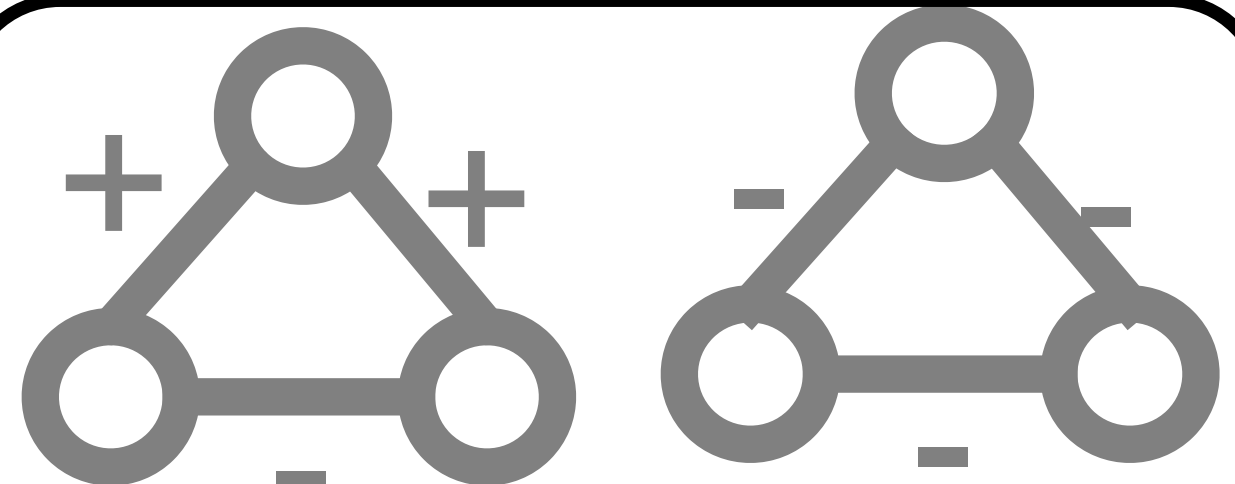
Enemy of friend is my enemy

Look at connected triples of nodes:



Balanced

Consistent with “friend of a friend” or  
“enemy of the enemy” intuition

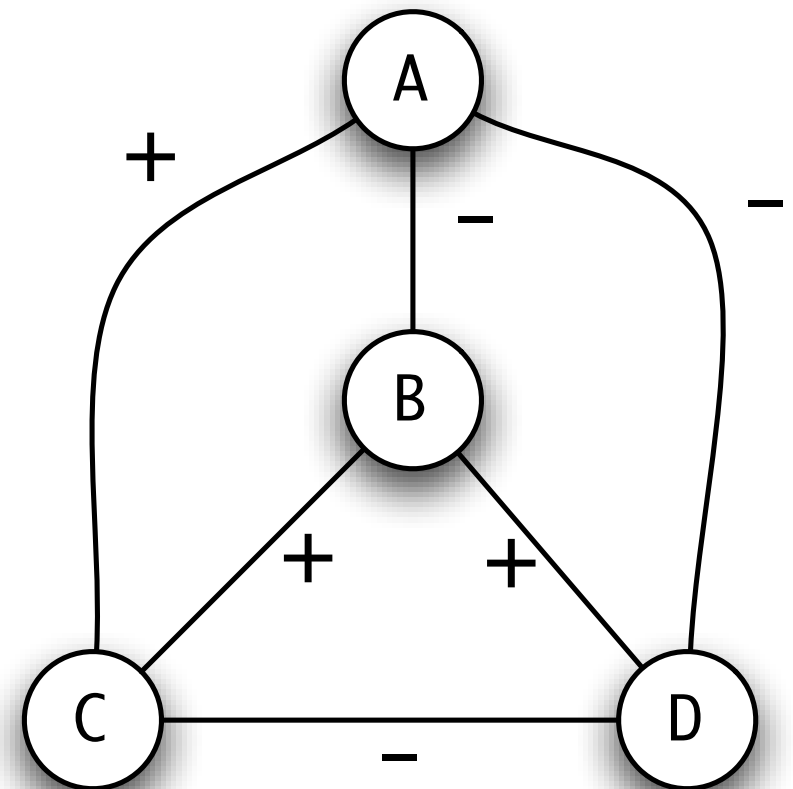
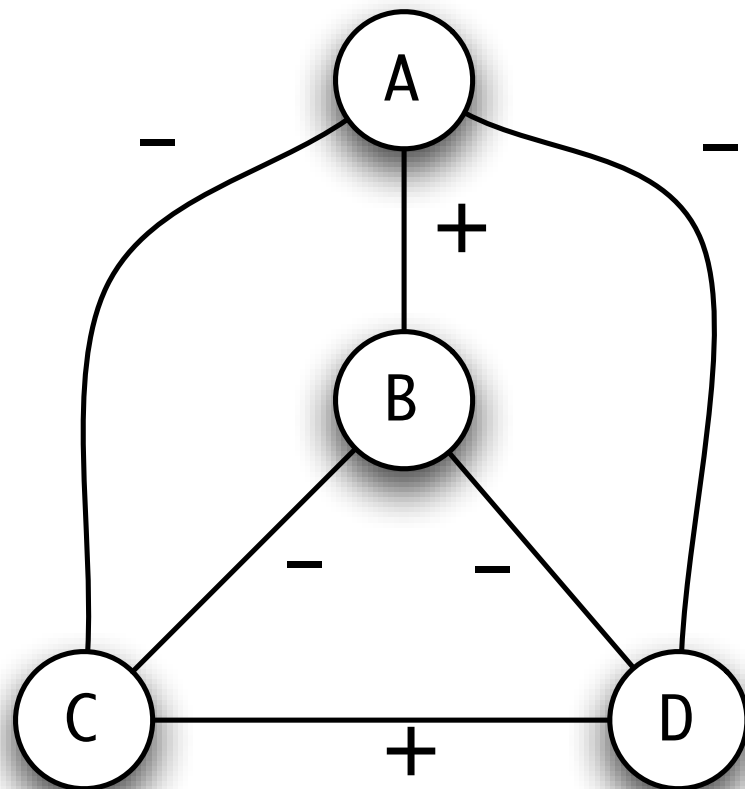


Unbalanced

Inconsistent with the “friend of a friend” or  
“enemy of the enemy” intuition



# Structural Balance

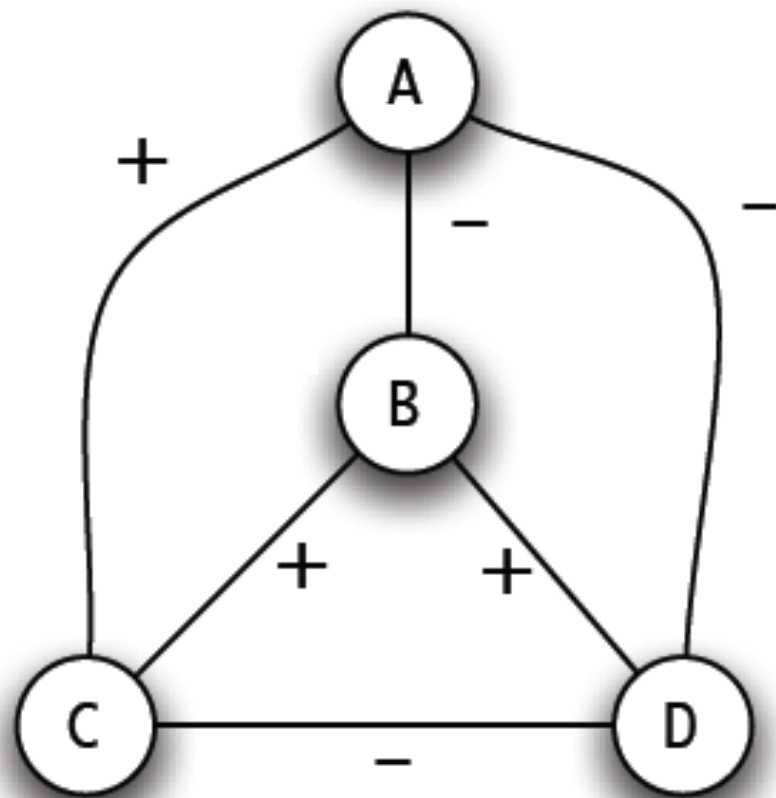


Which network is balanced?

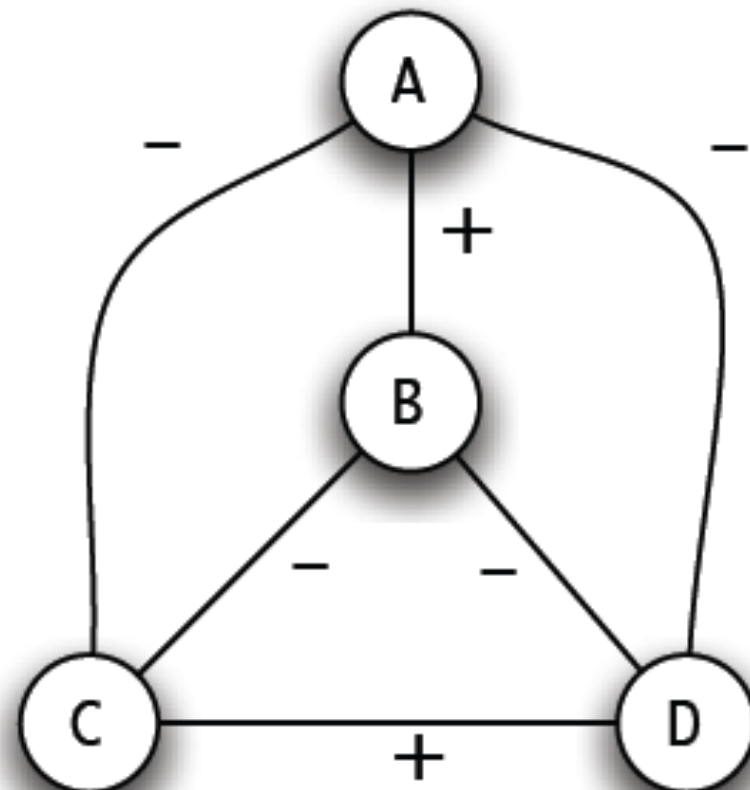
# Balanced/Unbalanced Networks

Define: A complete graph is *balanced* if every connected triple of nodes has:

All 3 edges labeled + **or** Exactly 1 edge labeled +

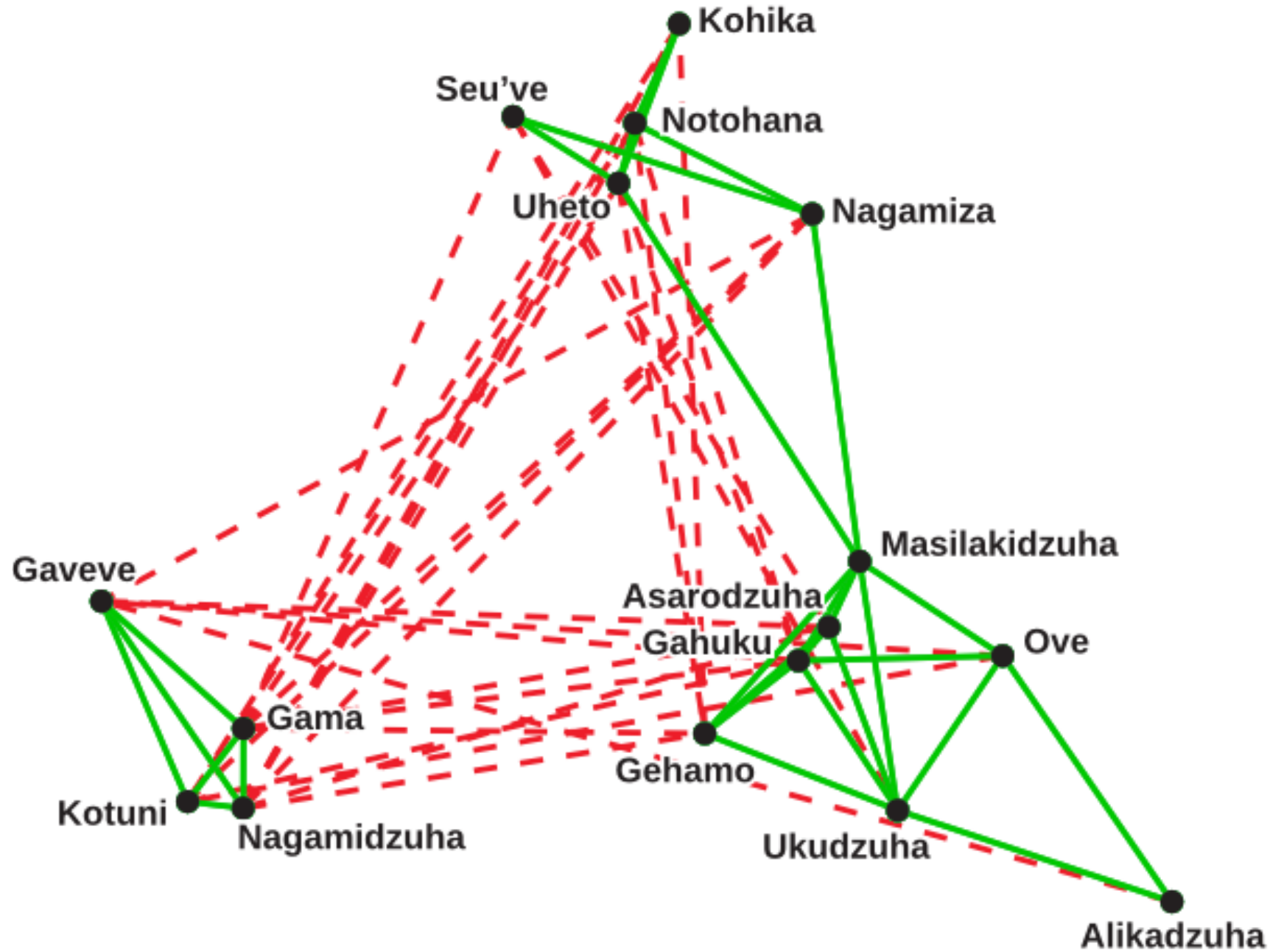


Unbalanced

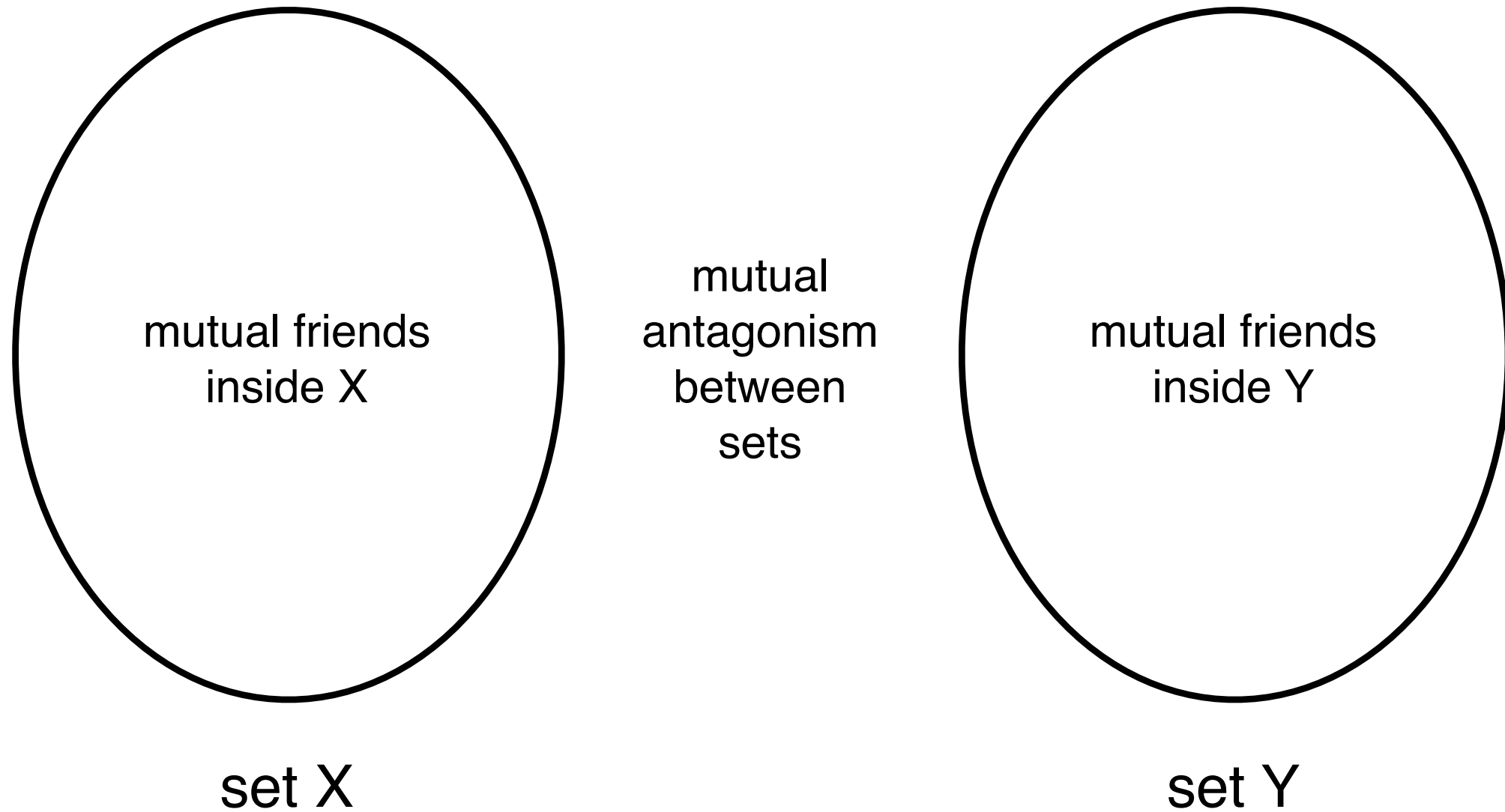


Balanced

# The Tribes of Eastern Central Highlands of New Guinea



# How general is this?





# Local Balance $\rightarrow$ Global Factions

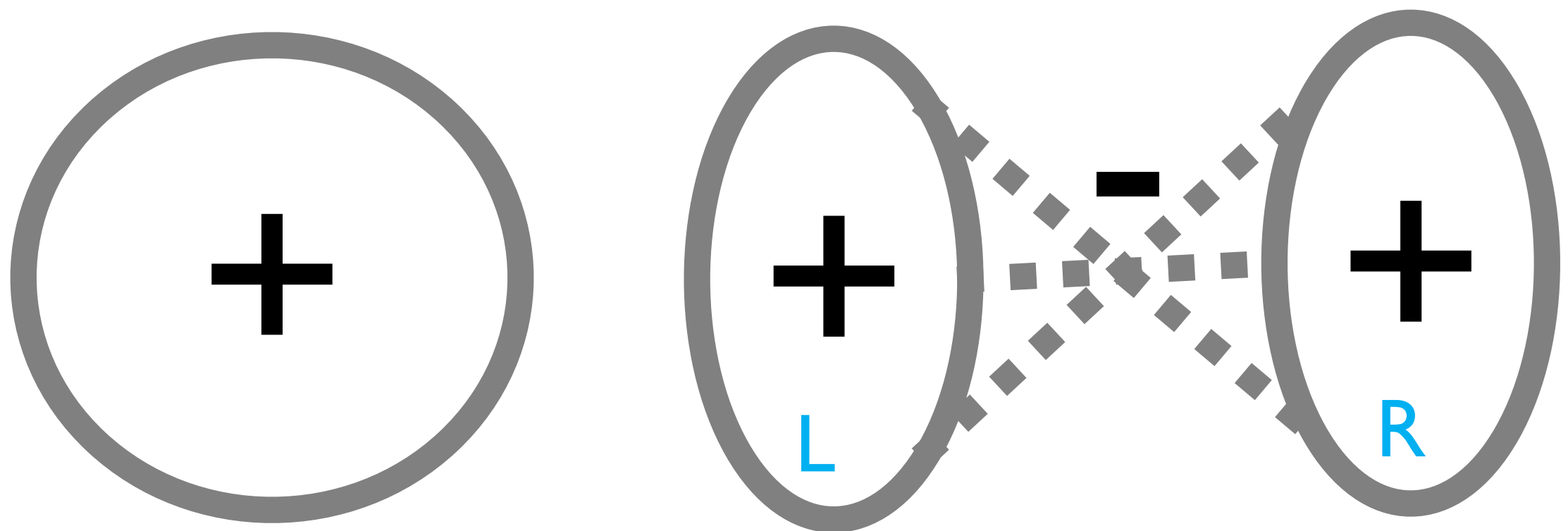
**The Balance Theorem: Balance implies global coalitions**

[Cartwright-Harary]

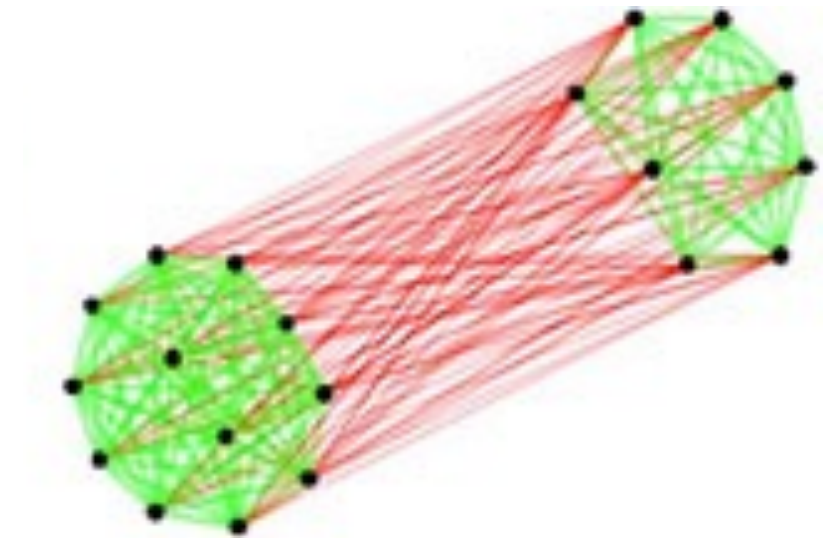
If **all triangles are balanced**, then either:

A) The network **contains only positive edges**, or

B) The network **can be split into two factions**: Nodes can be split into 2 sets where negative edges only point between the sets



# Balance Theorem



**Global coalitions  $\Rightarrow$  balance**

**Straightforward**

Every complete graph that looks like “this” is balanced

**Balance  $\Rightarrow$  Global coalitions**

**Less straightforward**

Every complete graph that’s balanced looks like “this”?

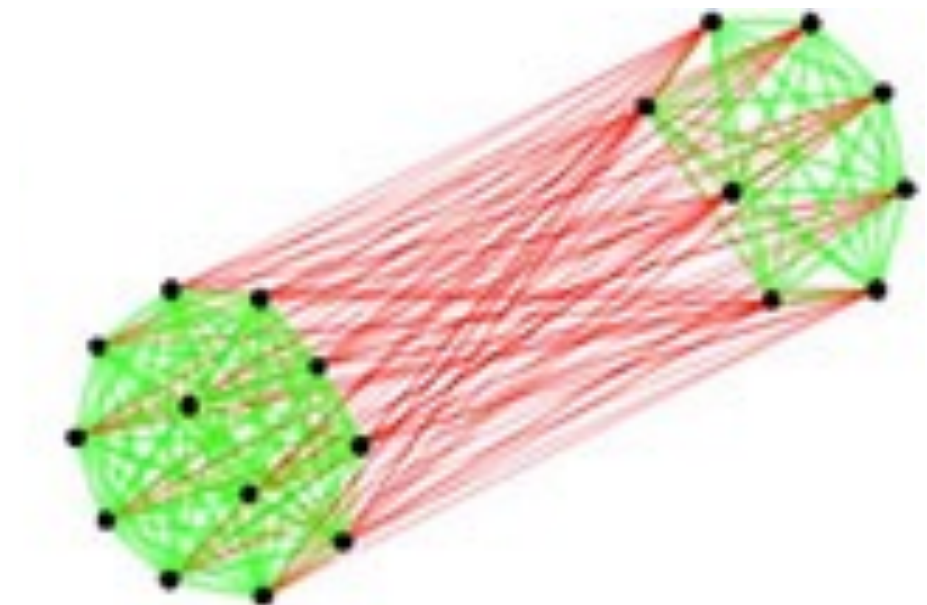
# Balance Theorem

**Global coalitions  $\Rightarrow$  balance:**

Any triangle is one of two types:

- A) All 3 nodes in one of the partitions
- B) 2 nodes in one partition, 1 in the other

- A): all 3 edges are +  $\longrightarrow$  balanced
- B): 2 nodes in one partition are +,  
other 2 edges are -  $\longrightarrow$  balanced



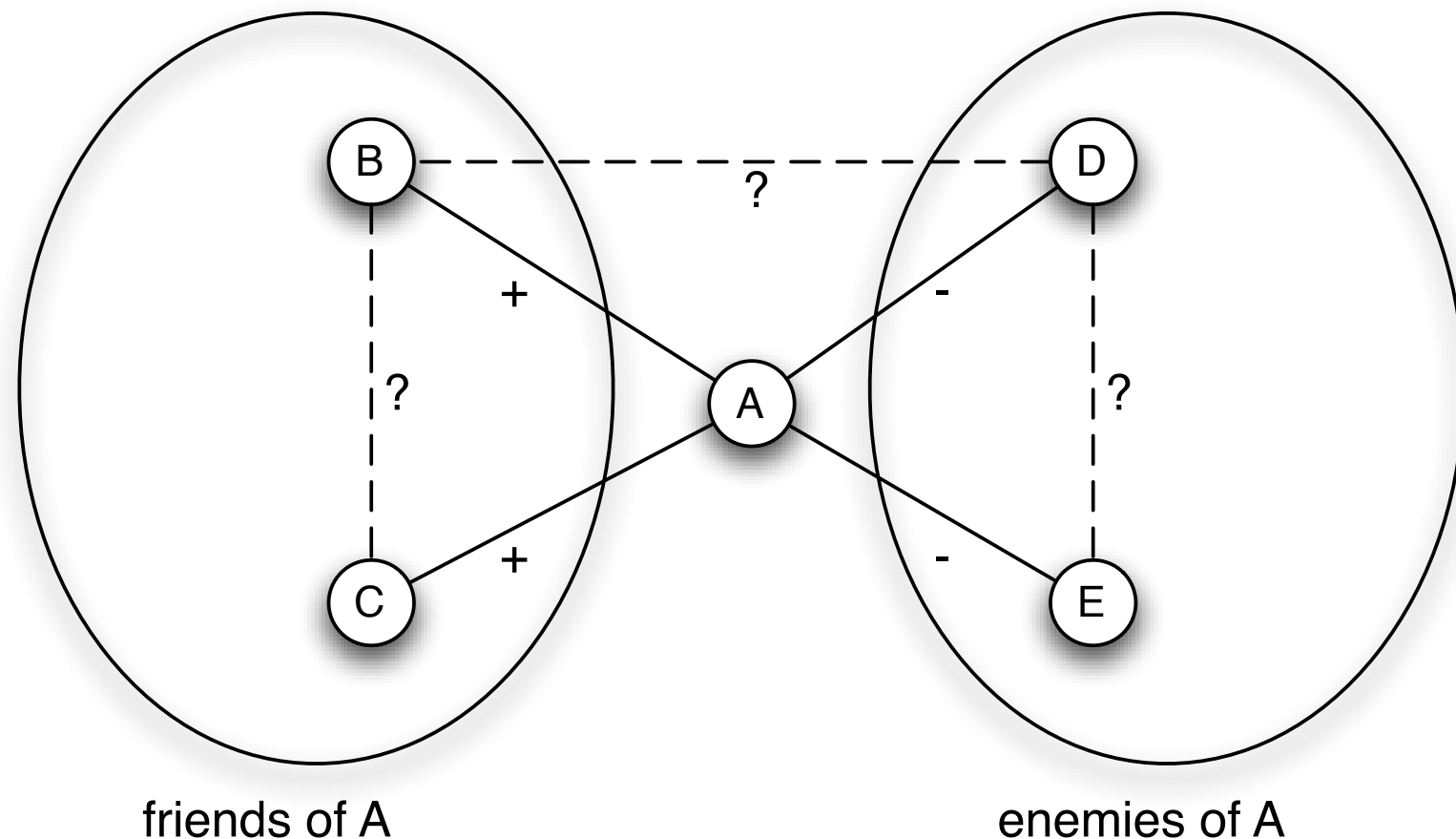
# Proof of Balance Theorem

**Balance  $\Rightarrow$  Global coalitions:**

Pick a node **A**.

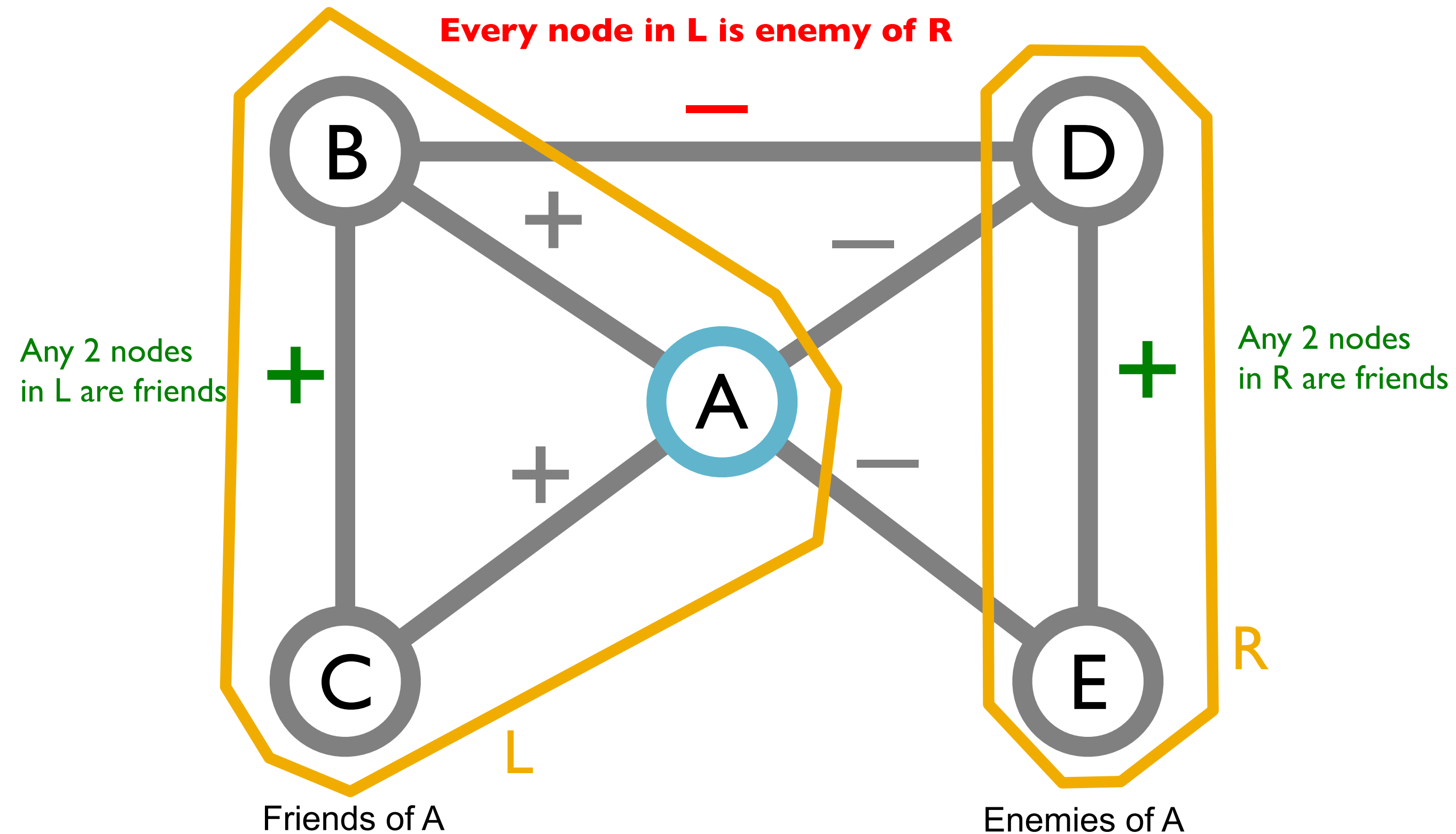
Because it's a complete graph, **A** is either friends or enemies with each person.

Now **check 3 cases:**

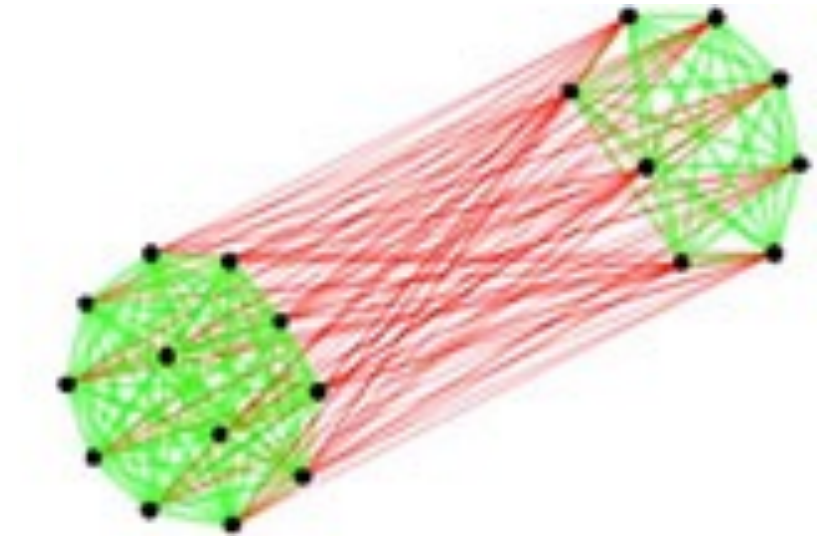




# Proof of Balance Theorem



# Balance Theorem



**Global coalitions  $\Rightarrow$  balance**

**Straight-forward**

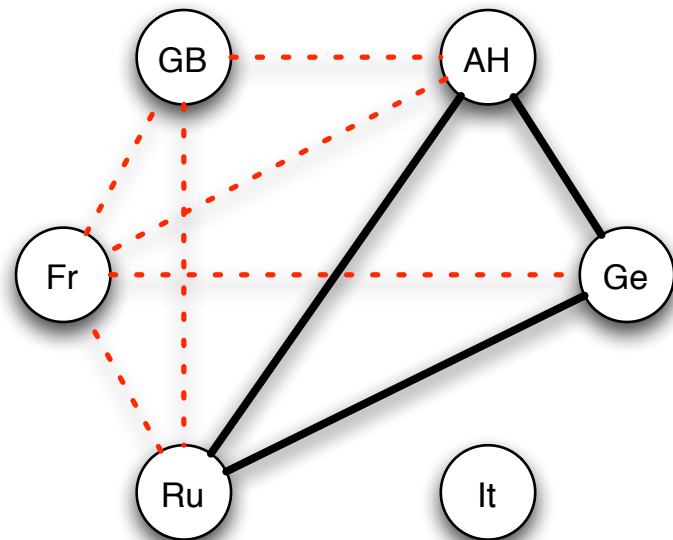
Every complete graph partitioned into two friendly coalitions that dislike either other is balanced

**Balance  $\Rightarrow$  Global coalitions**

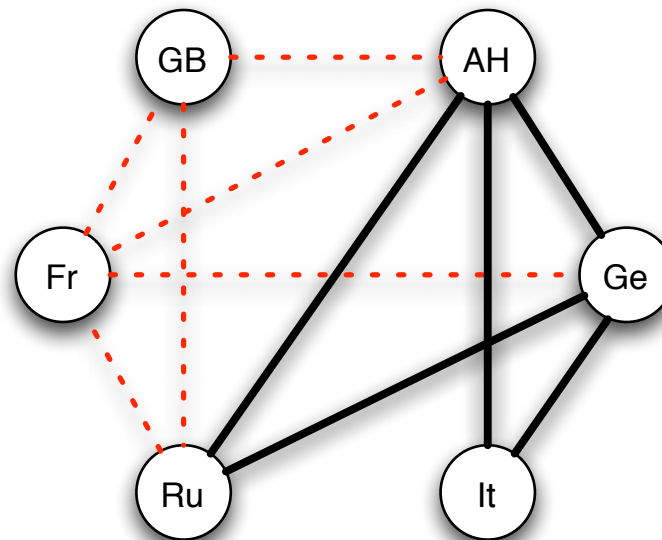
**Less straight-forward**

Every complete graph that's balanced can be partitioned into two friendly coalitions that dislike either other

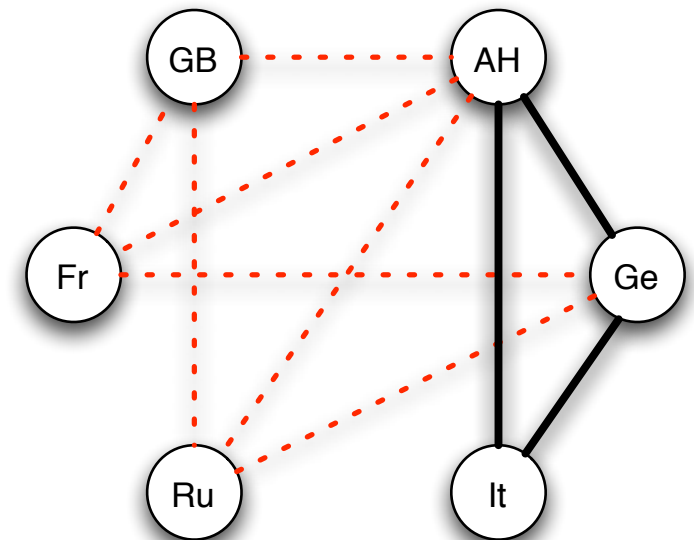
# European alliances, pre-WWI



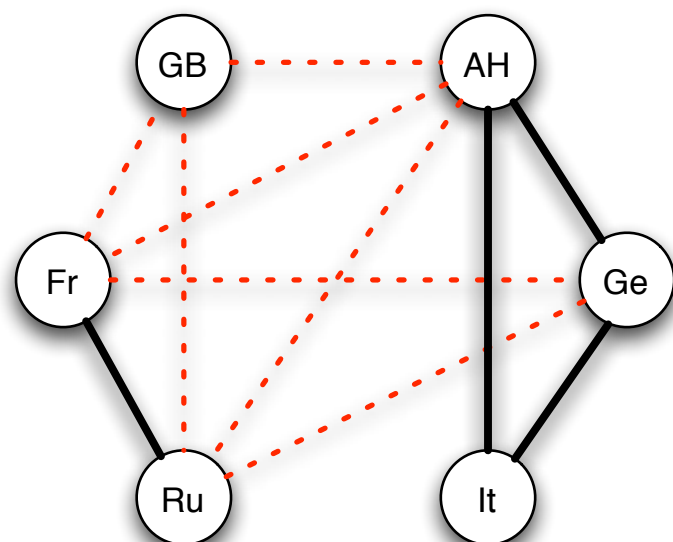
(a) *Three Emperors' League 1872-81*



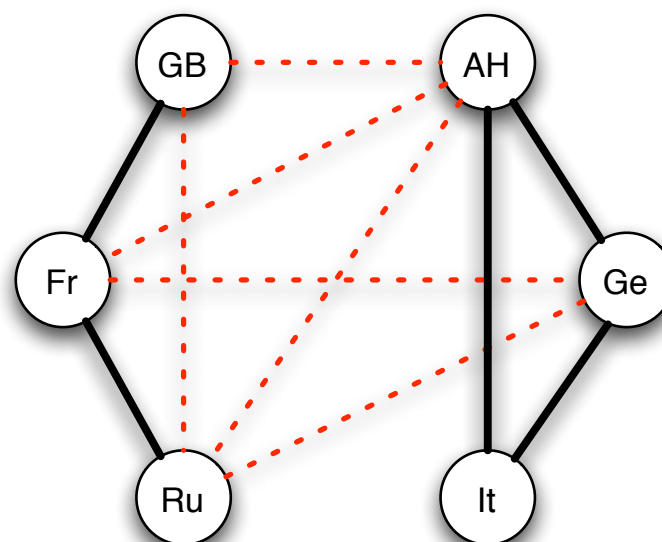
(b) *Triple Alliance 1882*



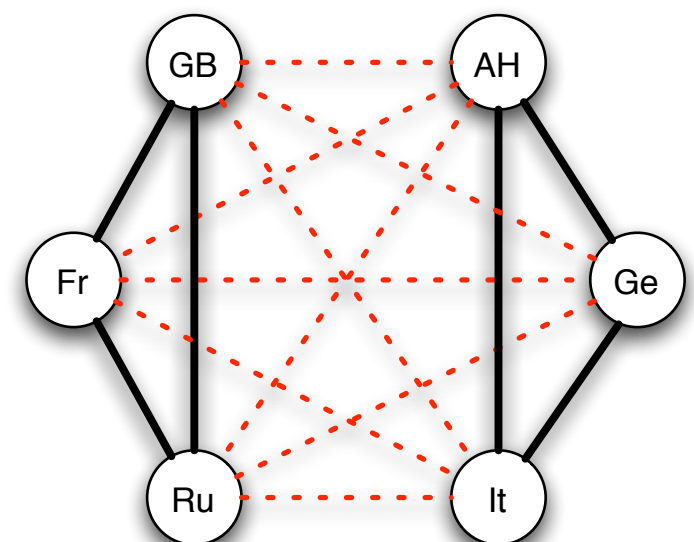
(c) *German-Russian Lapse 1890*



(d) *French-Russian Alliance 1891-94*



(e) *Entente Cordiale 1904*



(f) *British Russian Alliance 1907*

# Example: International Relations

## International relations:

**Positive edge:** alliance

**Negative edge:** animosity

Separation of Bangladesh from Pakistan in 1971: US supports Pakistan. **Why?**

USSR was the enemy of China

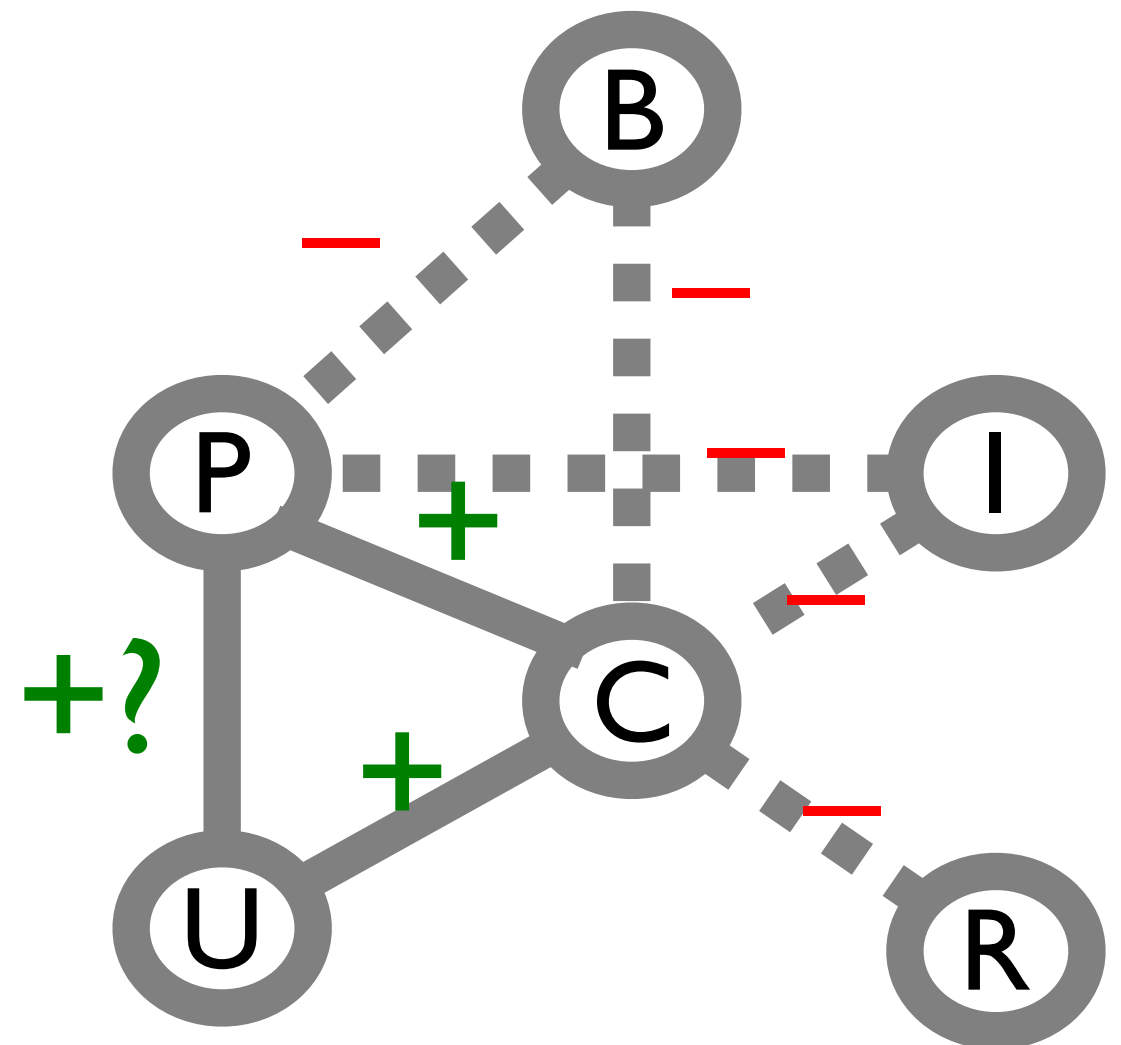
China was the enemy of India

India was the enemy of Pakistan

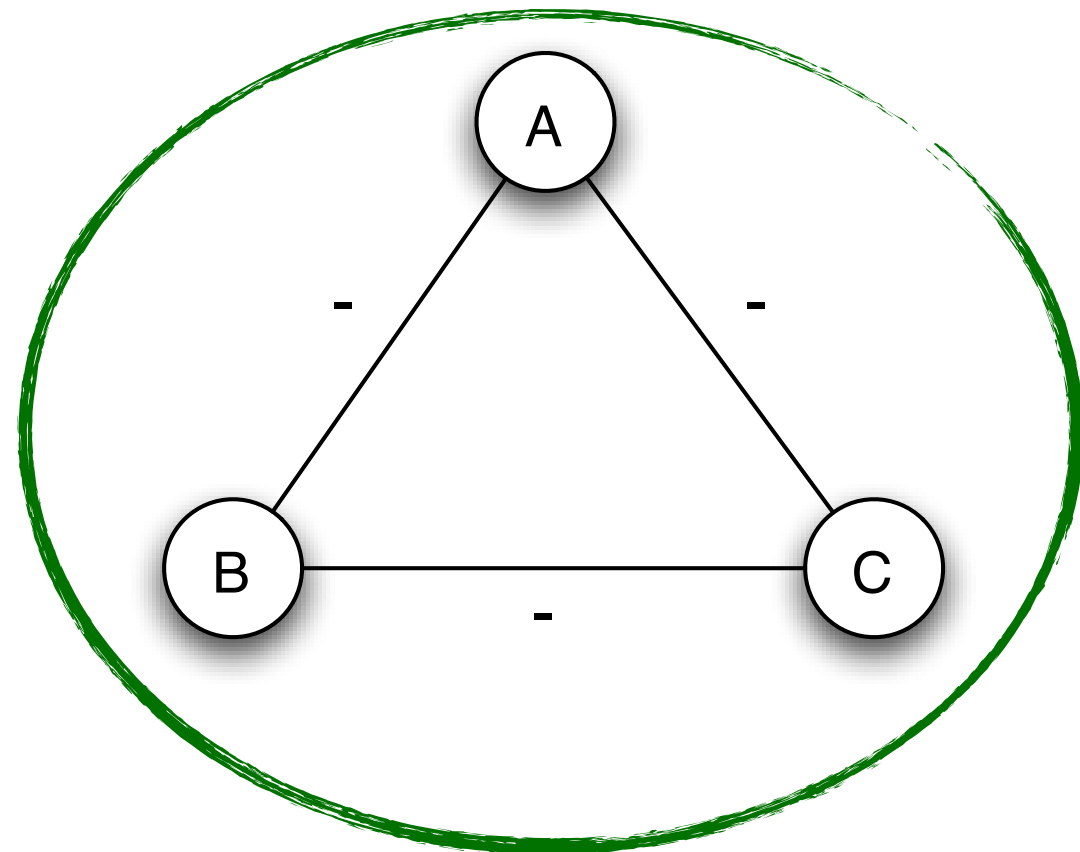
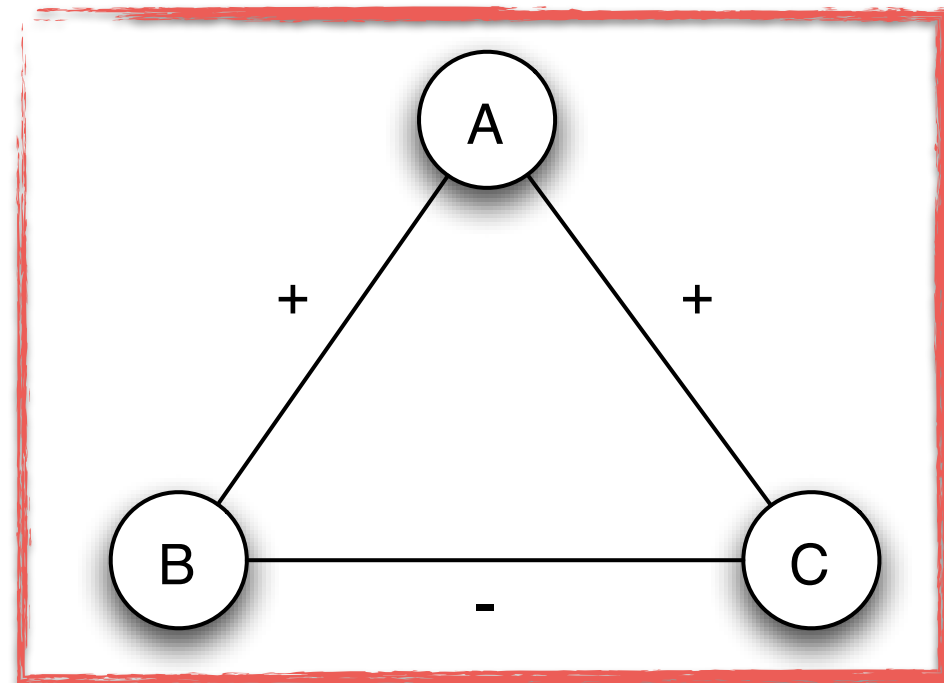
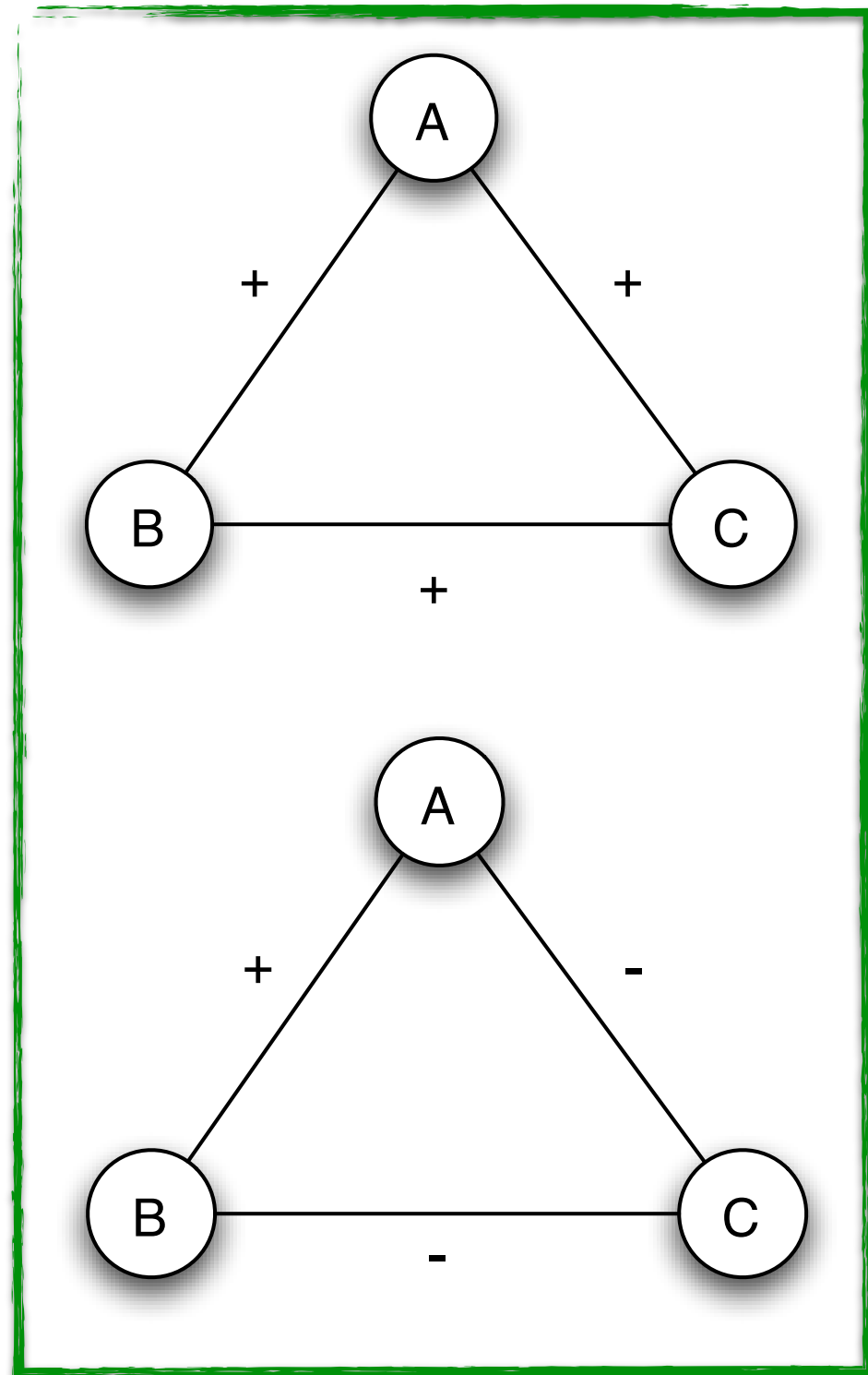
US was friendly with China

China vetoed

Bangladesh from U.N.



# Structural Balance



**What if we allow three mutual enemies?**

# Weak Structural Balance → Many Global Factions

Define: A complete network is *weakly balanced* if there is no triangle with exactly 2 positive edges and 1 negative edge.

## Characterization of Weakly Balanced Networks:

If a labeled complete graph is weakly balanced, then its nodes can be **partitioned**

(divided into groups such that two nodes belonging to the same group are friends, and every two nodes belonging to different groups are enemies)

**Global picture: same thing as before, but with many factions, not necessarily two**

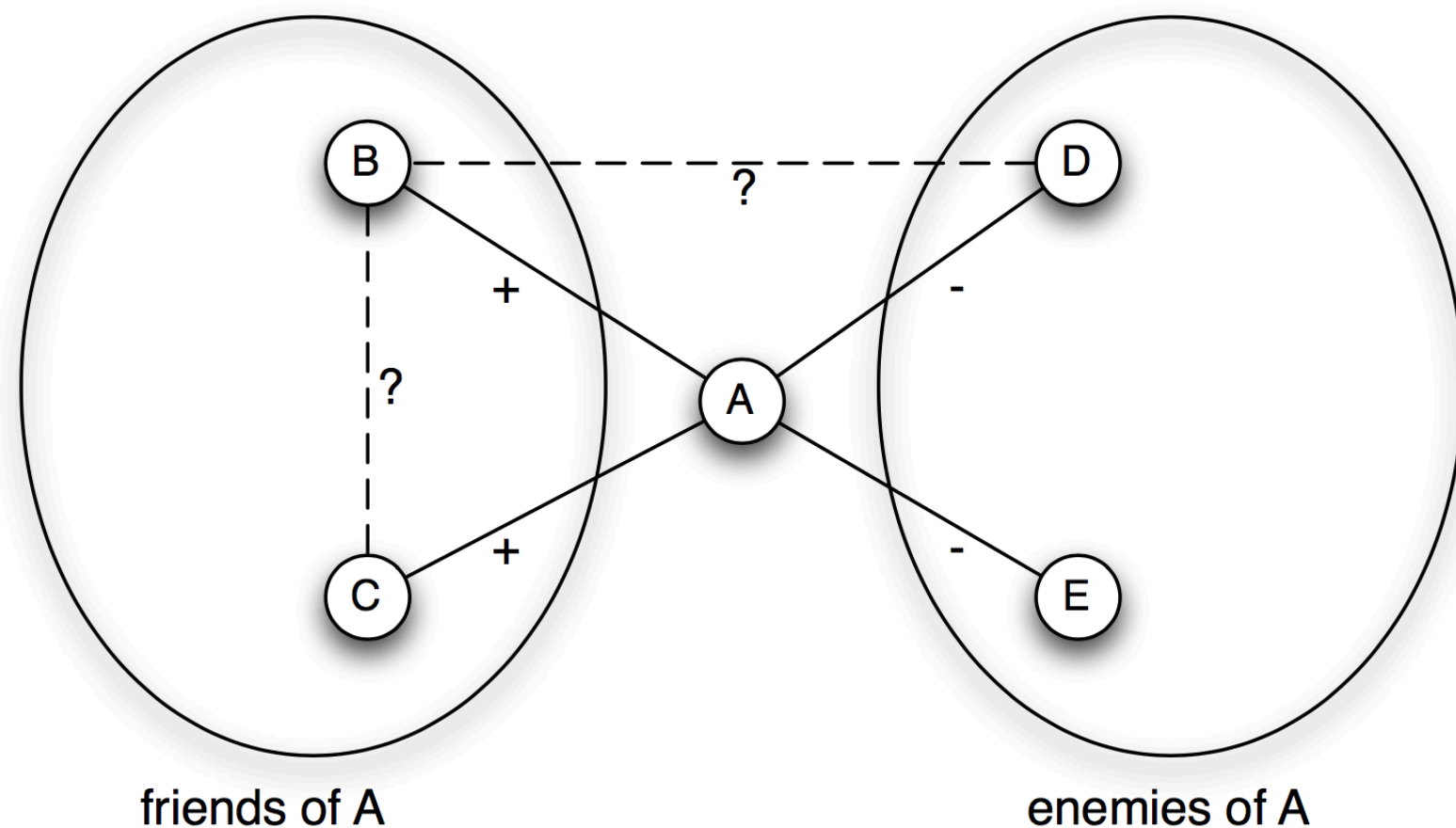


# Proof of Characterization

Pick a node **A**.

Because it's a complete graph, **A** is either friends or enemies with each person.

Now check 2 cases:

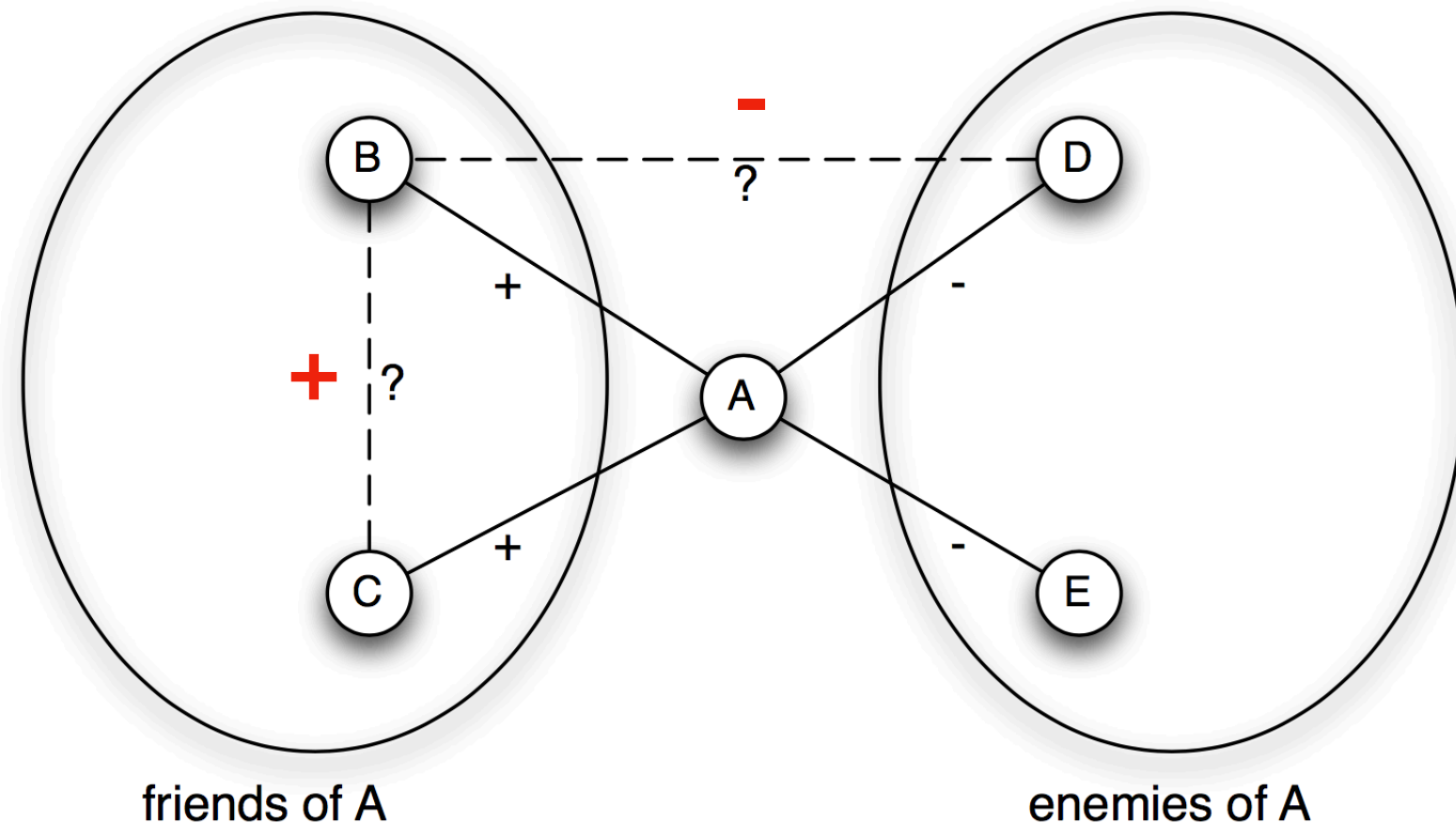


# Proof of Characterization

All of A's friends are friends with each other and are enemies with all of A's enemies

Remove A and his friends from the graph and **recurse!**

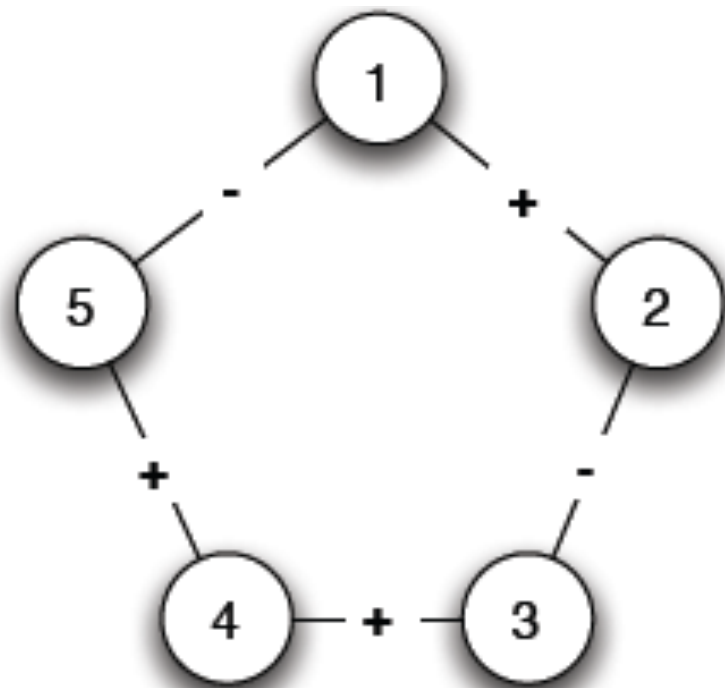
Graph still weakly balanced, find a second group, same argument applies, recurse until we've found all factions



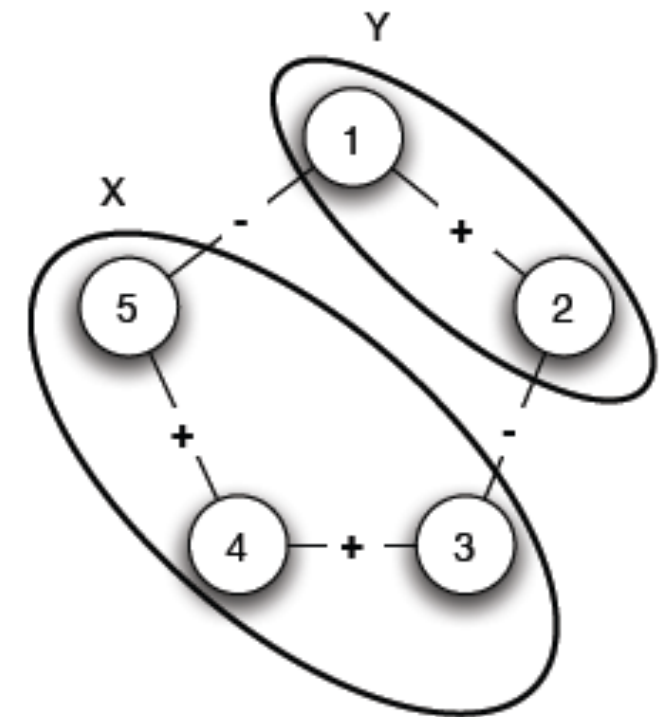
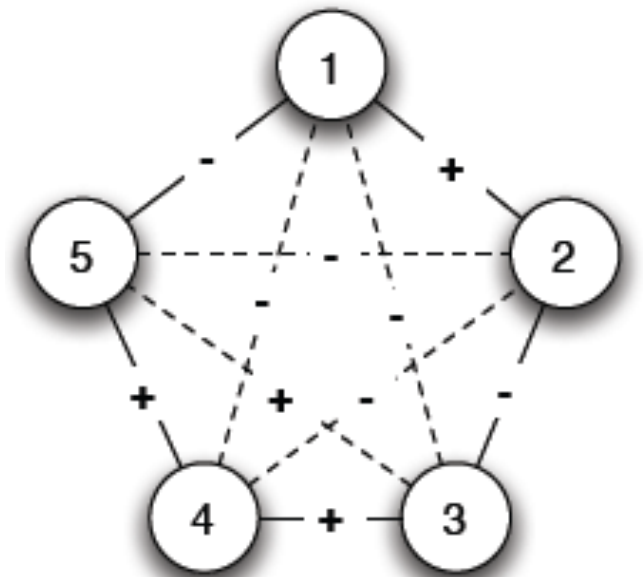
# Balance in General Networks

So far we've talked about complete graphs

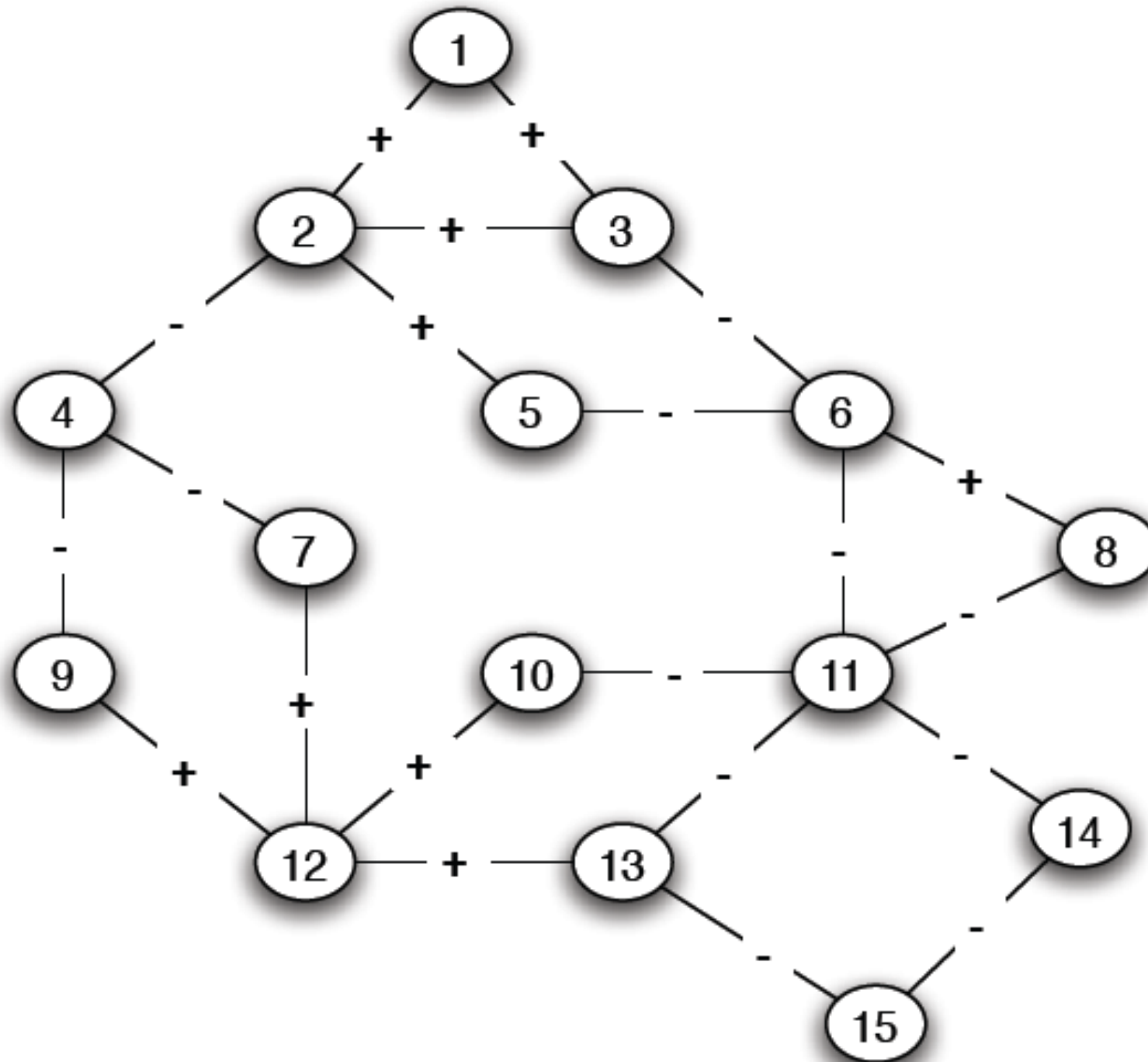
What about incomplete graphs?



Balanced?



# Signed Graph: Is it Balanced?

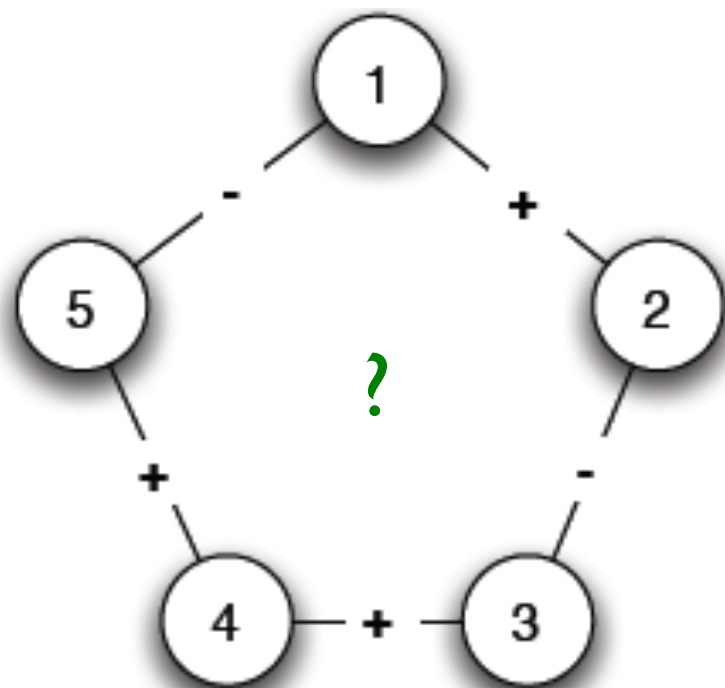


# Balance in General Networks

So far we talked about complete graphs

## Def 1: Local view

Fill in the missing edges to achieve balance



Balanced?

If the graph is “Balance-able”,  
then call it balanced

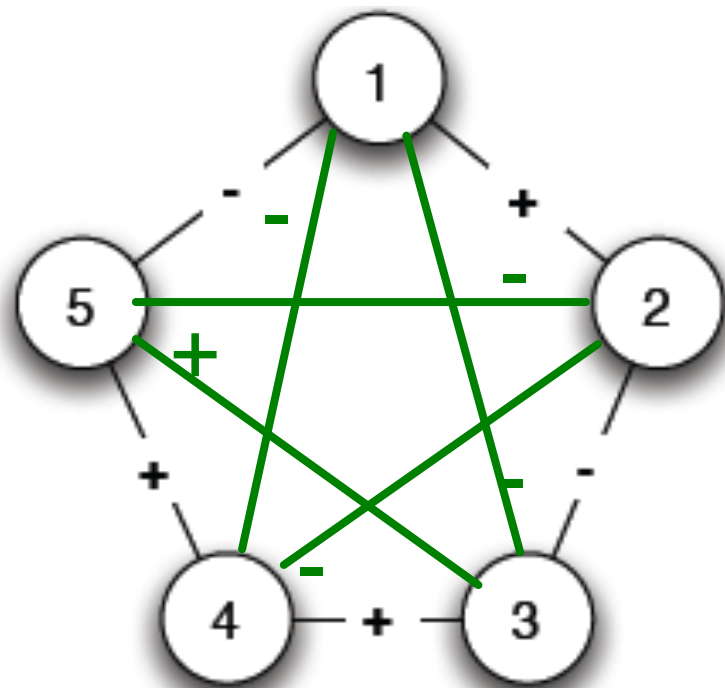
# Balance in General Networks

So far we talked about complete graphs

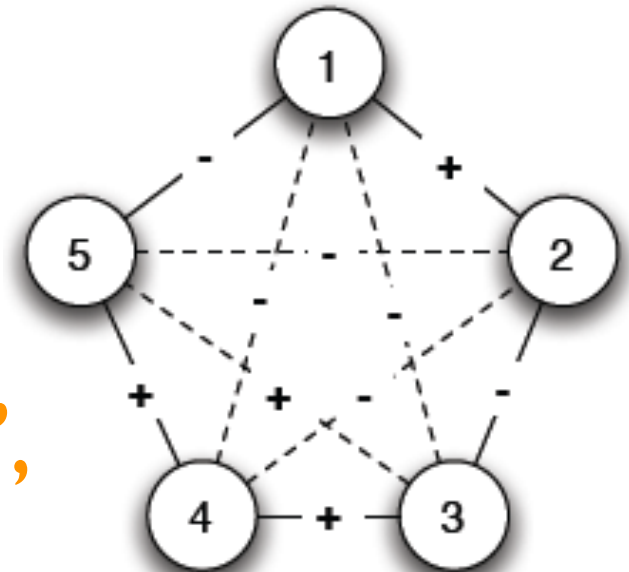
## Def 1: Local view

Fill in the missing edges to achieve balance

If the graph is “Balance-able”, then call it balanced



Balanced?





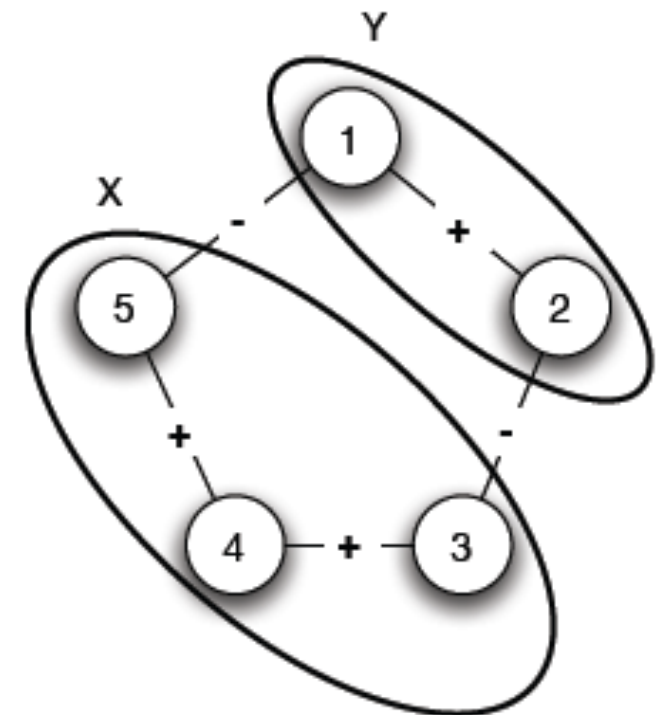
# Balance in General Networks

So far we talked about complete graphs

## Def 2: Global view

Divide the graph into two coalitions

If you can separate the graph into coalitions as before, call it balanced

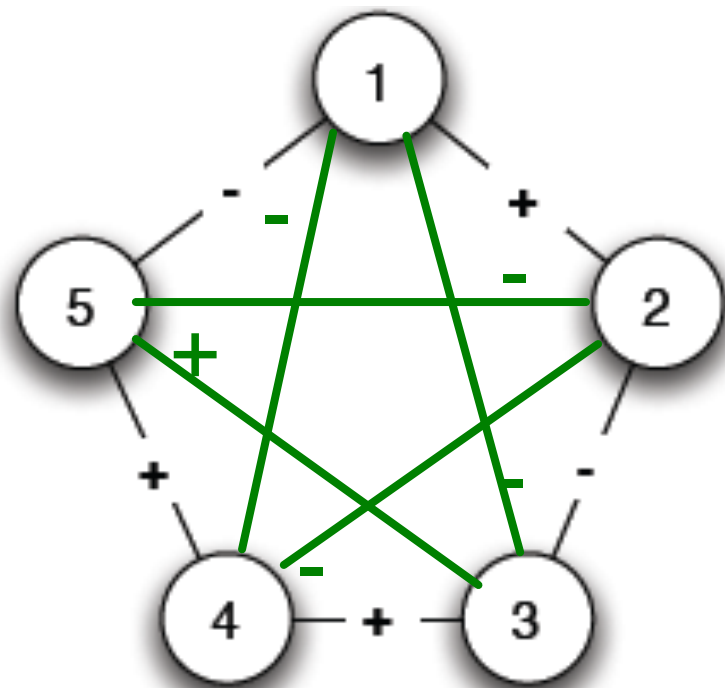


# Balance in General Networks

So far we talked about complete graphs

## Def 1: Local view

Fill in the missing edges to achieve balance

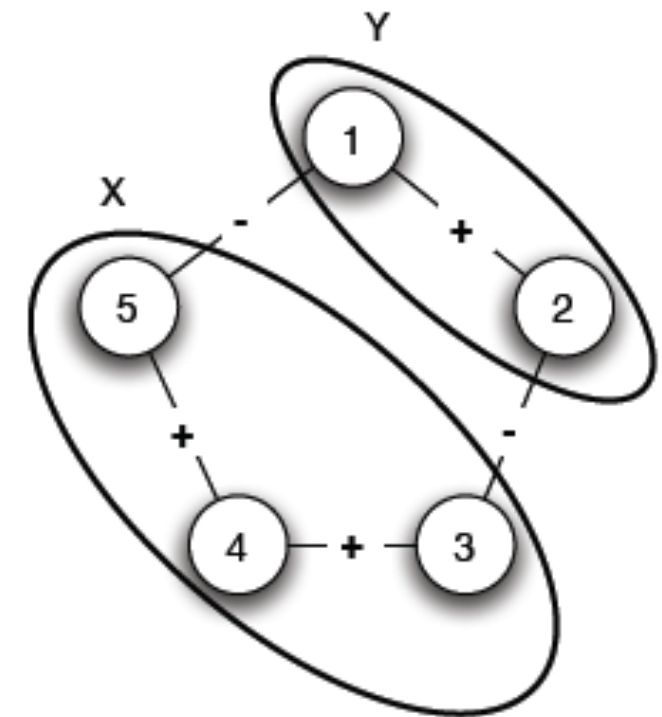
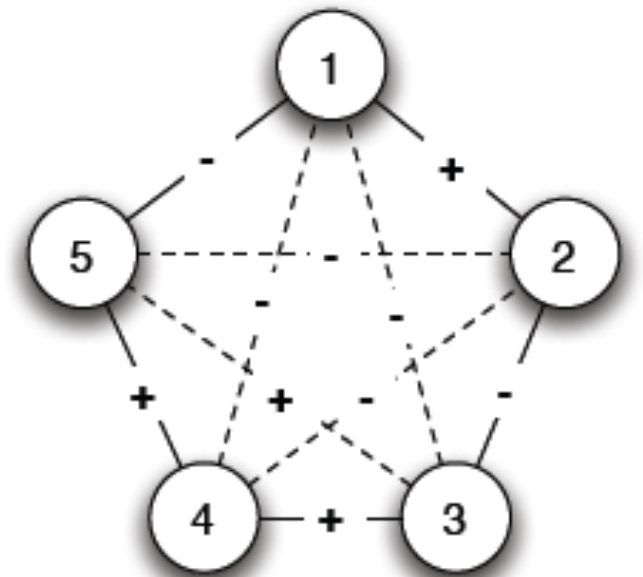


Balanced?

## Def 2: Global view

Divide the graph into two coalitions

The 2 definitions are **equivalent!**



# Balance in General Networks

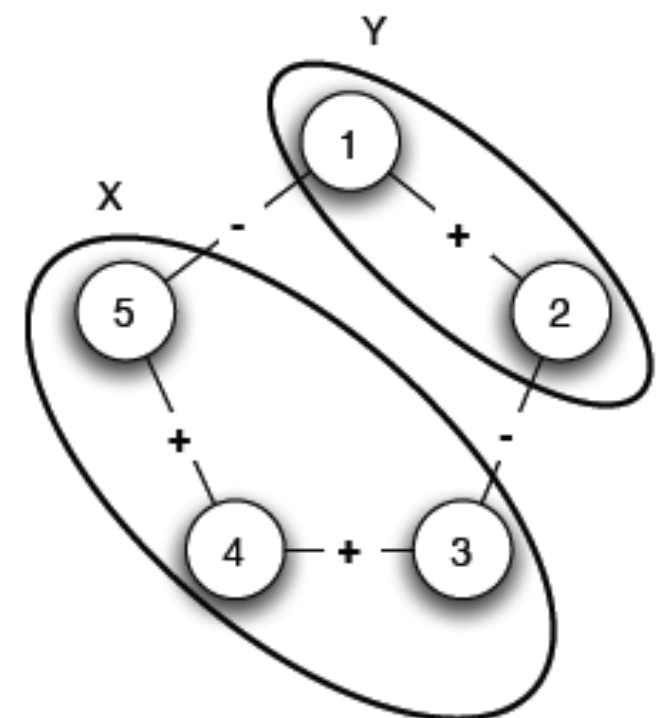
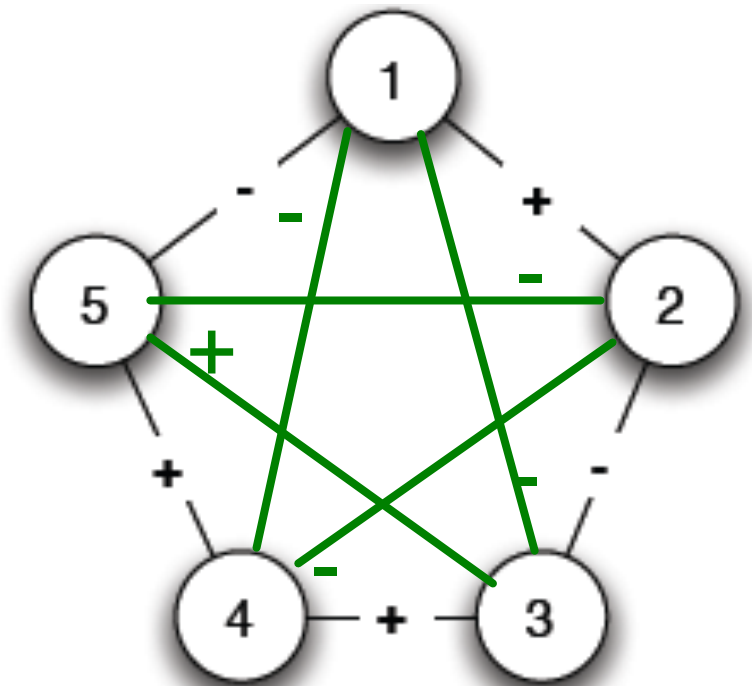
Claim: in general (not necessarily complete) networks, the **local** and **global** definitions of balance are equivalent

## Def 1: Local view

Fill in the missing edges to achieve balance

## Def 2: Global view

Divide the graph into two coalitions

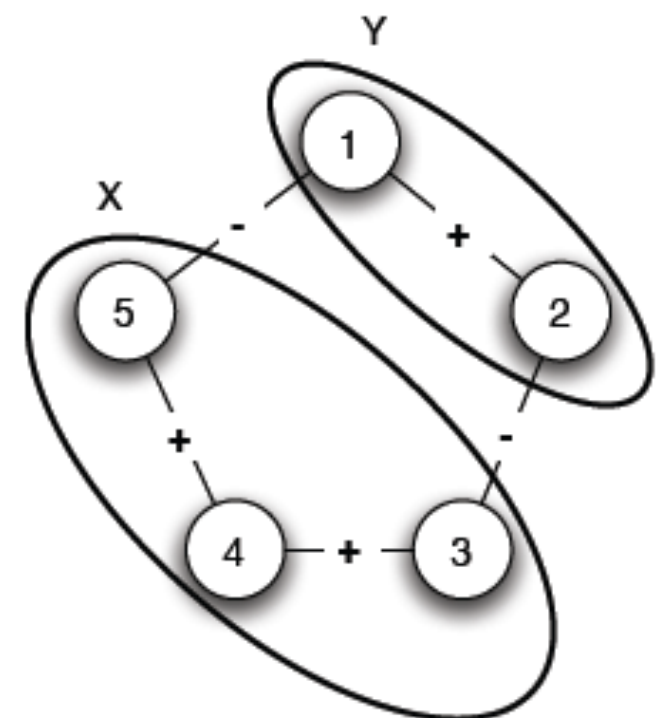
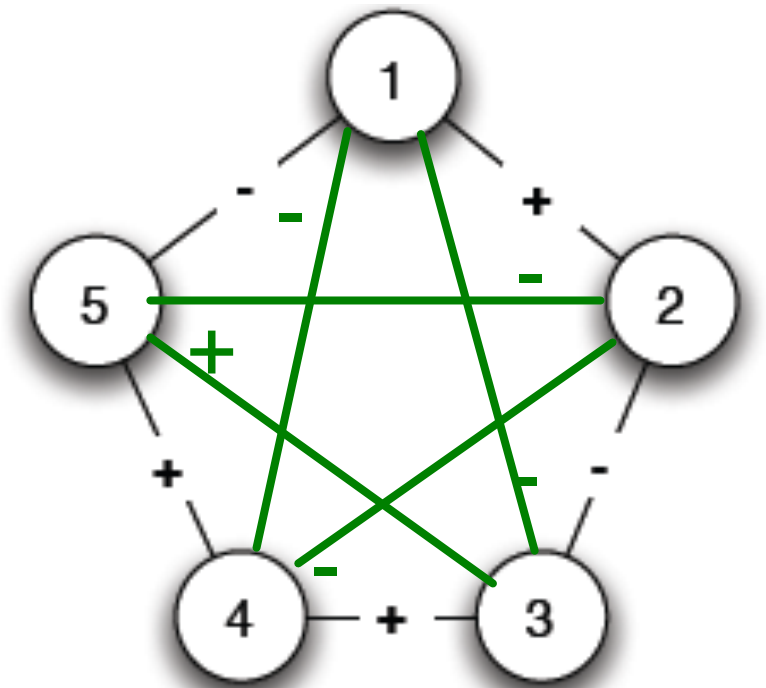


# Balance in General Networks

**Actually easy to see:**

**Local  $\Rightarrow$  global:** (if you can fill in edges such that the resulting complete graph is balanced, then it can be divided into coalitions)

After filling in, we have a complete network as before, the Balance Theorem applies

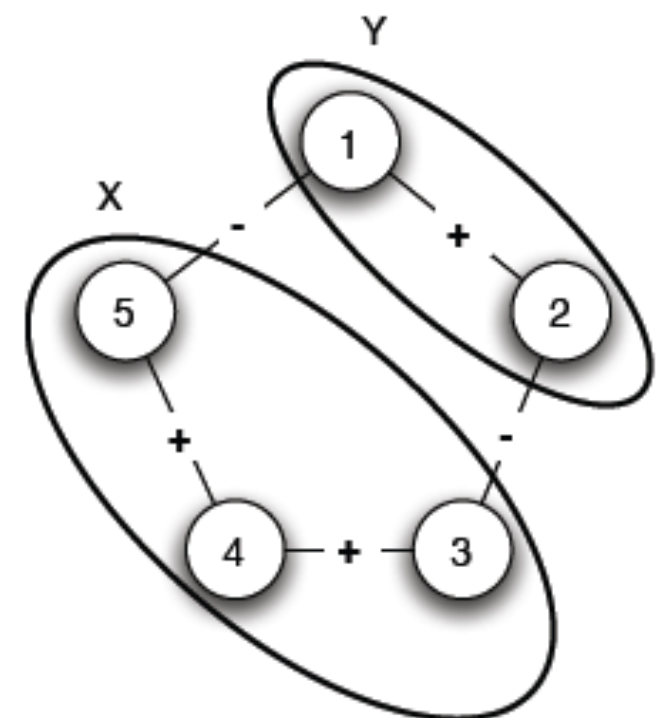
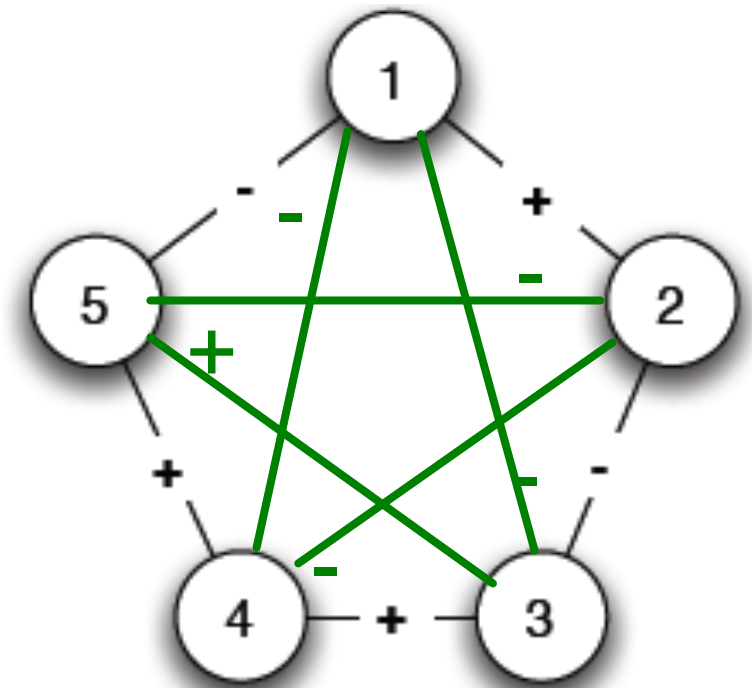


# Balance in General Networks

**Actually easy to see:**

**Global => local:** (if the graph can be divided into coalitions, then you can fill in edges that results in a complete balanced graph )

Fill in edges within and between coalitions as before: positive edges within the coalitions and negative edges between them



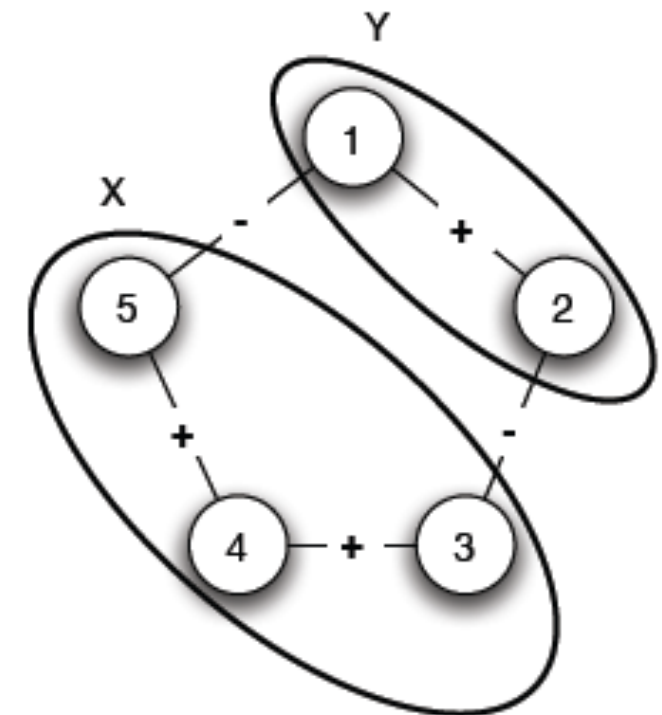
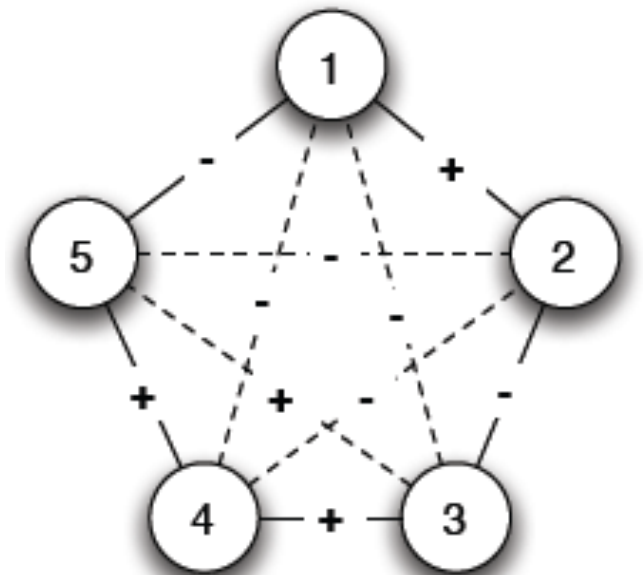
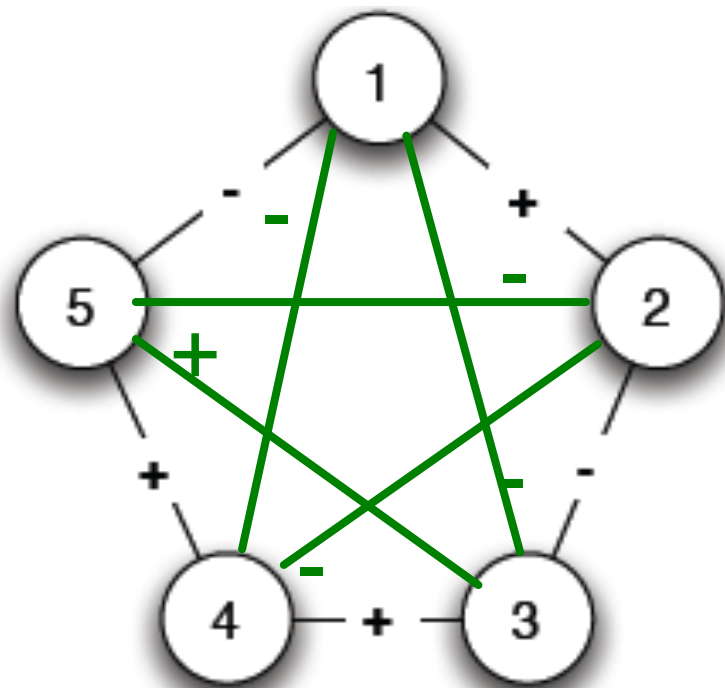
# Balance in General Networks

**Actually easy to see:**

**Local => global:** after filling in, result in complete network as before

**Global => local:** fill in edges within and between coalitions as before

**Done!**

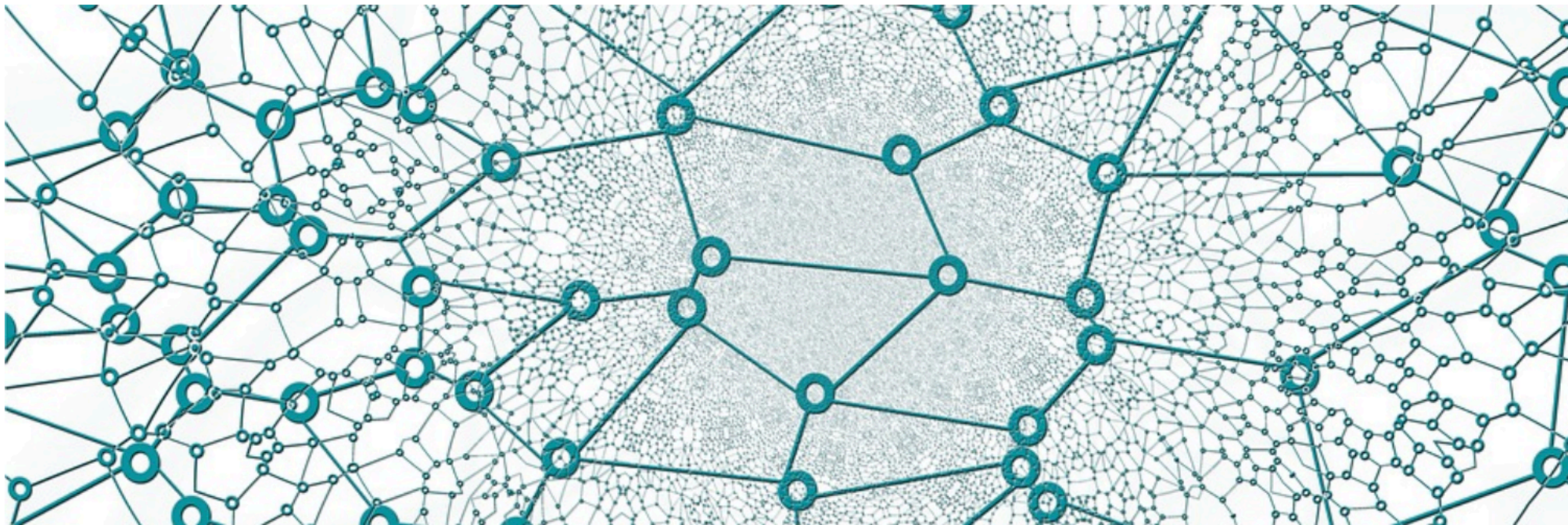




# Balance in General Networks

We have a natural definition for **balance** in general signed networks

“**Natural**” because we arrived at it **two different ways** that **turn out to be equivalent**

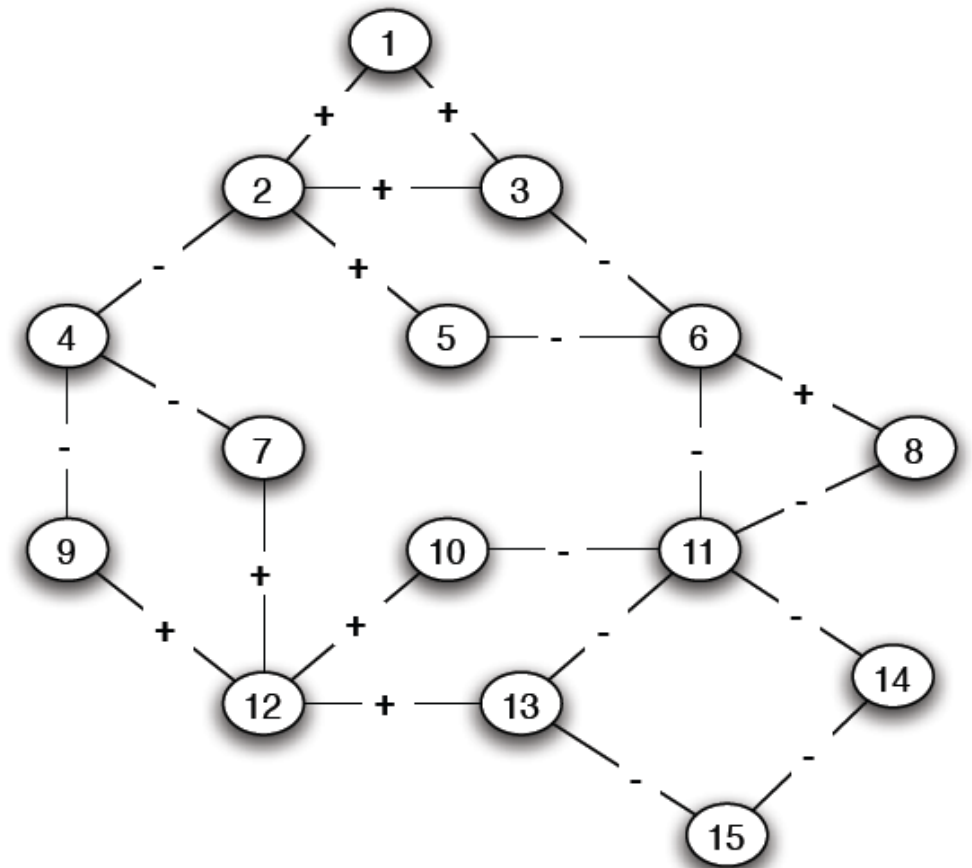


# Balance in General Networks

We have a natural definition for **balance** in general signed networks

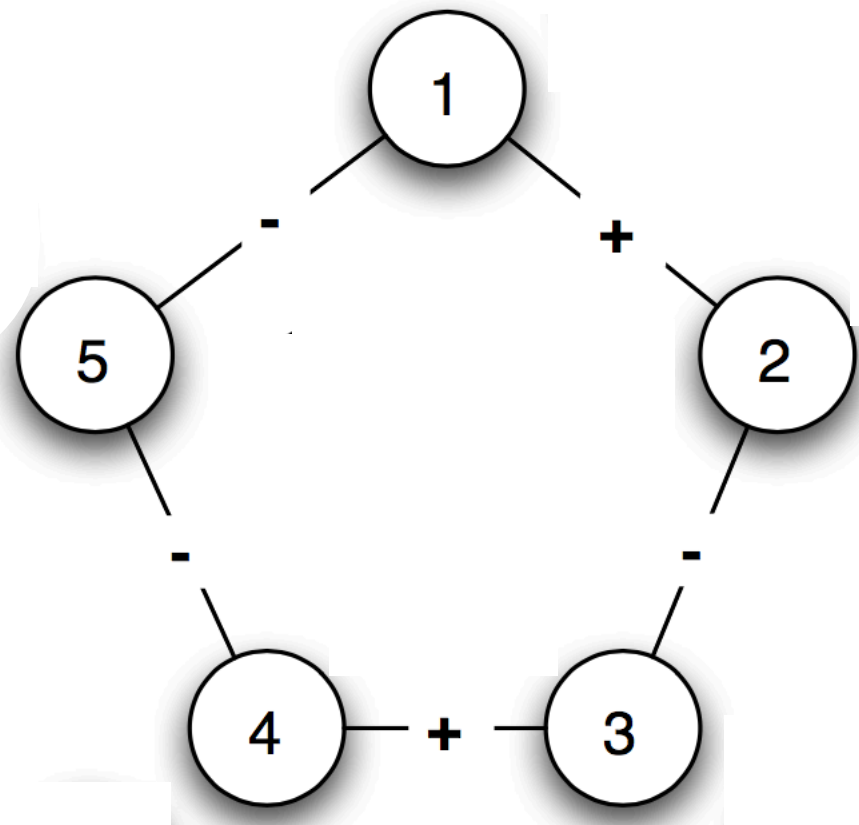
“**Natural**” because we arrived at it **two different ways** that **turn out to be equivalent**

But, there's a problem: **how to actually check if a network is balanced in this way?**



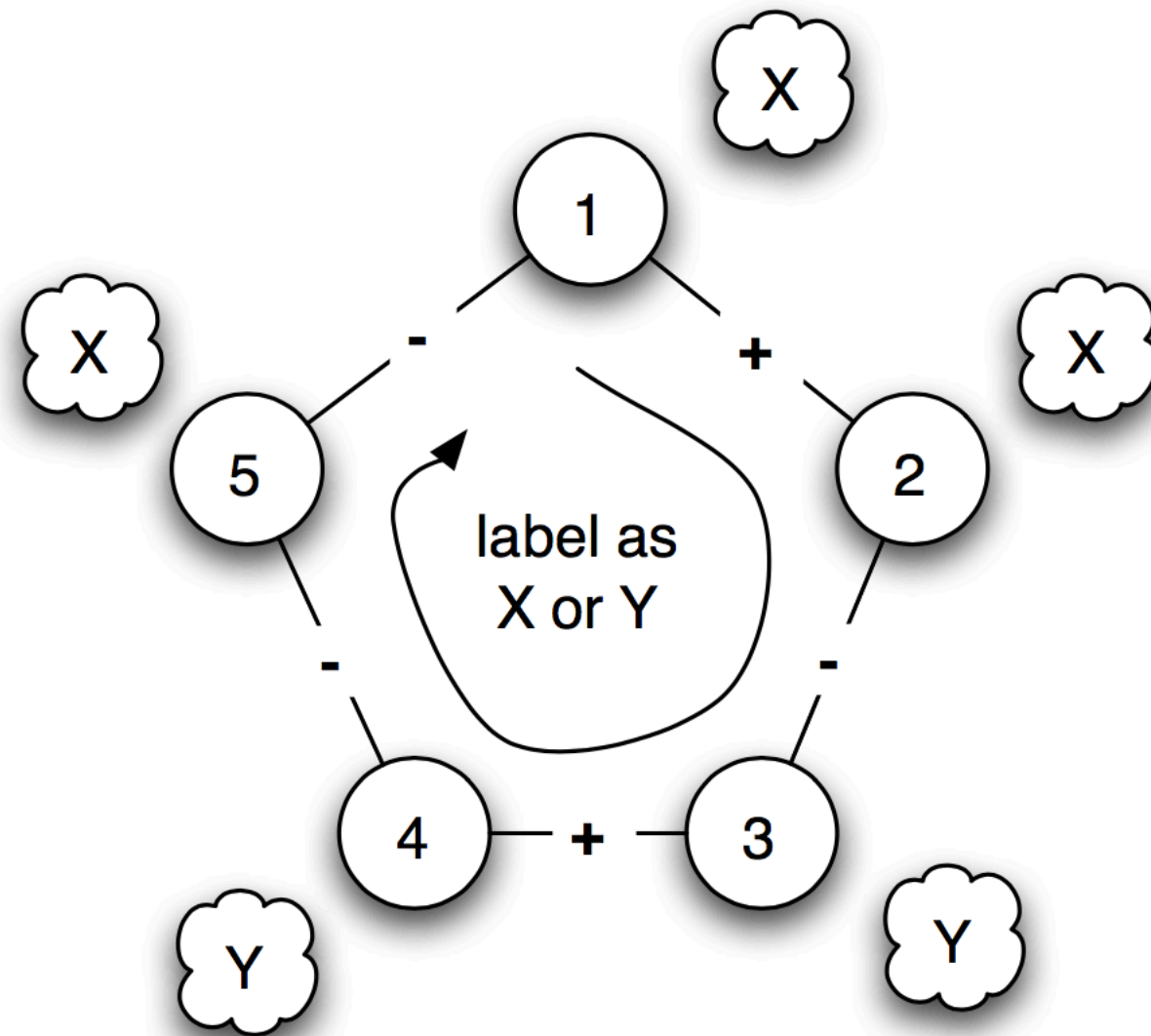
# Balance in General Networks

Why isn't this graph balanced?



# Balance in General Networks

Why isn't this graph balanced?

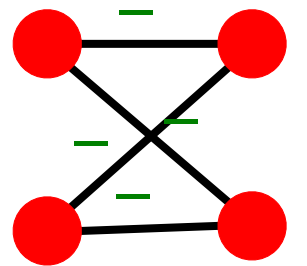


Walk around a cycle, every time we see a negative edge we have to switch coalitions

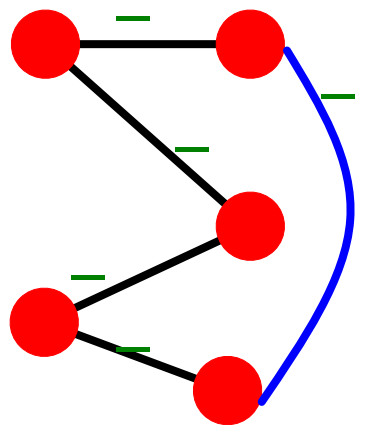
# Is a Signed Network Balanced?

Theorem: Graph is **balanced** if and only if it contains **no cycle with an odd number of negative** edges

[Harary 1953, 1956]



Even length  
cycle



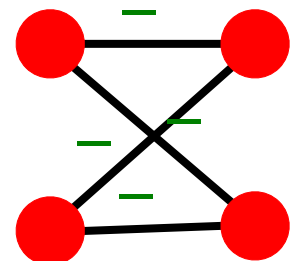
Odd length  
cycle

# Is a Signed Network Balanced?

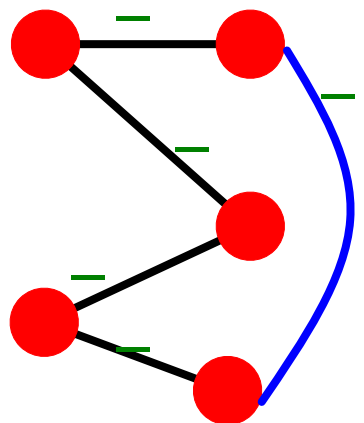
Theorem: Graph is **balanced** if and only if it contains **no cycle with an odd number of negative** edges

[Harary 1953, 1956]

This theorem is saying that the **only way** a graph can be unbalanced is if there is **a cycle with an odd number of negative cycles**. That's the only possible problem!



Even length cycle



Odd length cycle

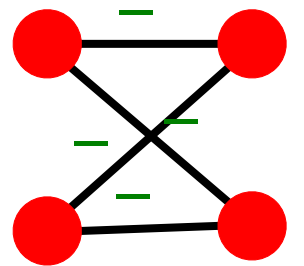


# Is a Signed Network Balanced?

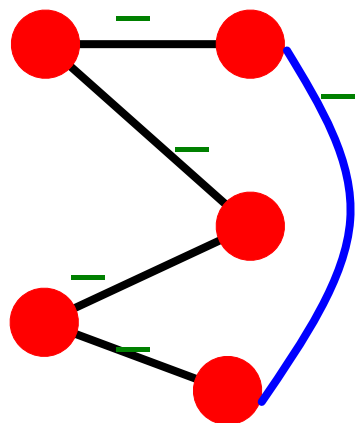
Theorem: Graph is **balanced** if and only if it contains **no cycle with an odd number of negative** edges

[Harary 1953, 1956]

Proof: We will show that every graph is either **balanced** or contains **a cycle with odd number of negative edges**



Even length cycle



Odd length cycle

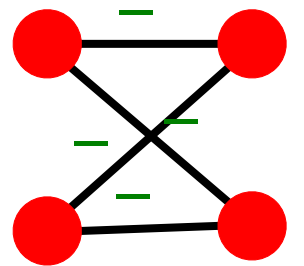


# Is a Signed Network Balanced?

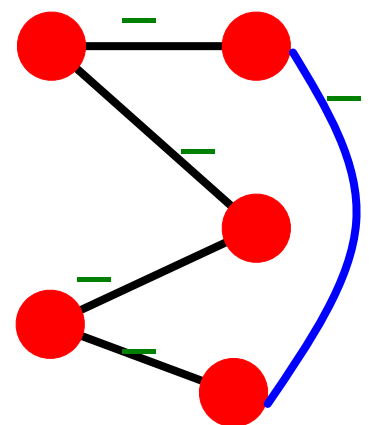
Theorem: Graph is **balanced** if and only if it contains **no cycle with an odd number of negative edges**  
[Harary 1953, 1956]

Proof by algorithm: We will do this by actually constructing an algorithm that either **outputs a division into coalitions** or a **cycle with odd number of negative edges**

Because these are the **only two outcomes**, this **proves the claim**



Even length cycle



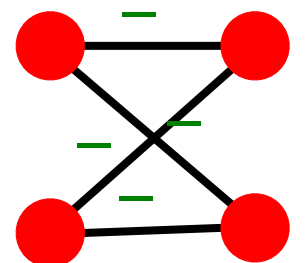
Odd length cycle

# Is a Signed Network Balanced?

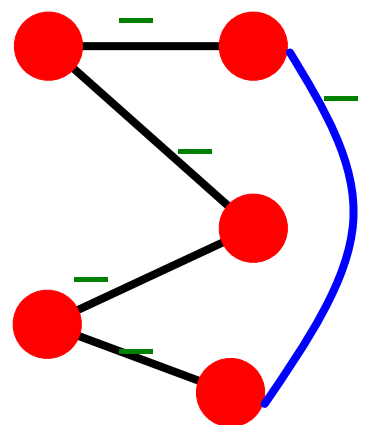
Theorem: Graph is **balanced** if and only if it contains **no cycle with an odd number of negative** edges

[Harary 1953, 1956]

Proof sketch: Our algorithm will try to assign nodes to coalitions such that the graph is balanced. We will reason that the **only way it can fail** is if there is a cycle with an odd number of negative edges.



Even length cycle



Odd length cycle

# Is a Signed Network Balanced?

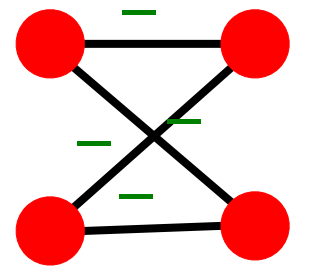
## Signed graph algorithm:

**Step 1:** Find connected components on + edges and for each component create a super-node

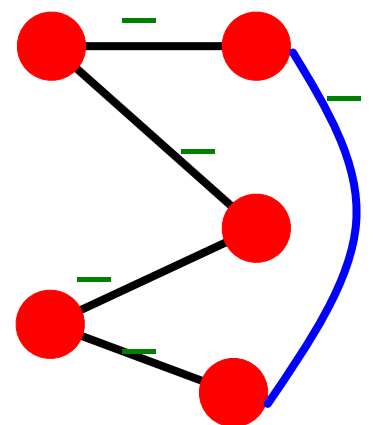
- Since nodes connected by a + edge must be in same coalition
- If any – edge in the super node, done (cycle with 1 negative edge)

**Step 2:** Connect components A and B if there is a negative edge between the members

- Note there are only negative edges pointing out of a super-node (otherwise should've connected the two super-nodes that have a positive edge)



Even length cycle

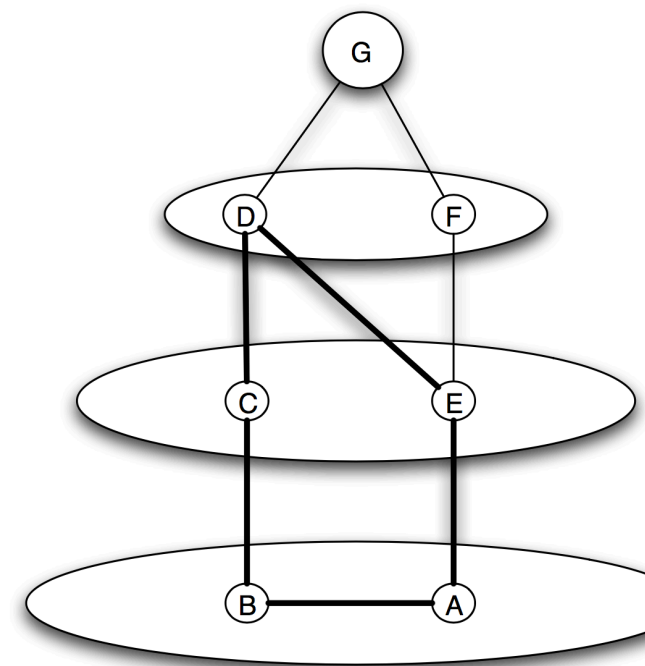


Odd length cycle

# Is a Signed Network Balanced?

## Signed graph algorithm

- Now we have a graph on super-nodes joined by negative edges
- Just need to consistently assign super-nodes to coalitions X and Y
- BFS starting at any node in the super-node graph (which only has – edges)
- Produces a set of layers of increasing distances from the root
- Call all even layers X and odd layers Y
- If edges are only between adjacent layers (not within-layer), then all – edges point between X and Y, **balanced**!
- Otherwise, within-layer edge A-B. Cycle G-A-B-G has length  $2k+1$ , therefore it's odd, therefore **unbalanced**!



# Is a Signed Network Balanced?

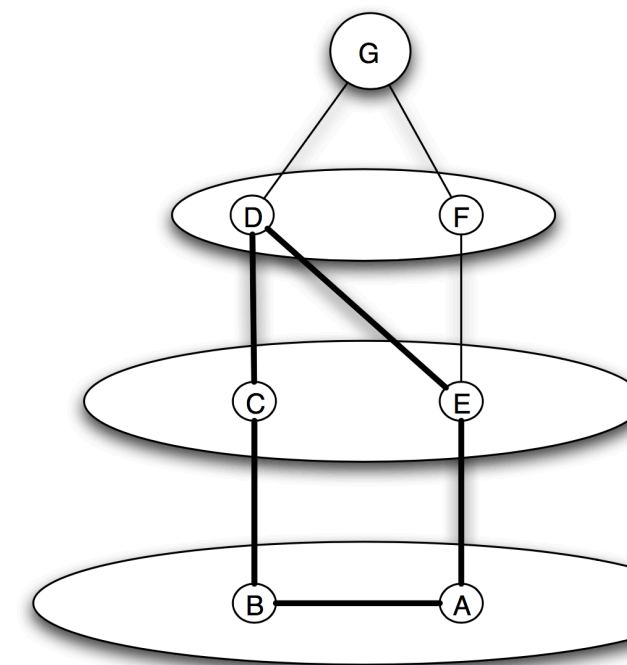
Two outcomes:

1) label each super-node as either X or Y, in such a way that every edge has endpoints with opposite labels.

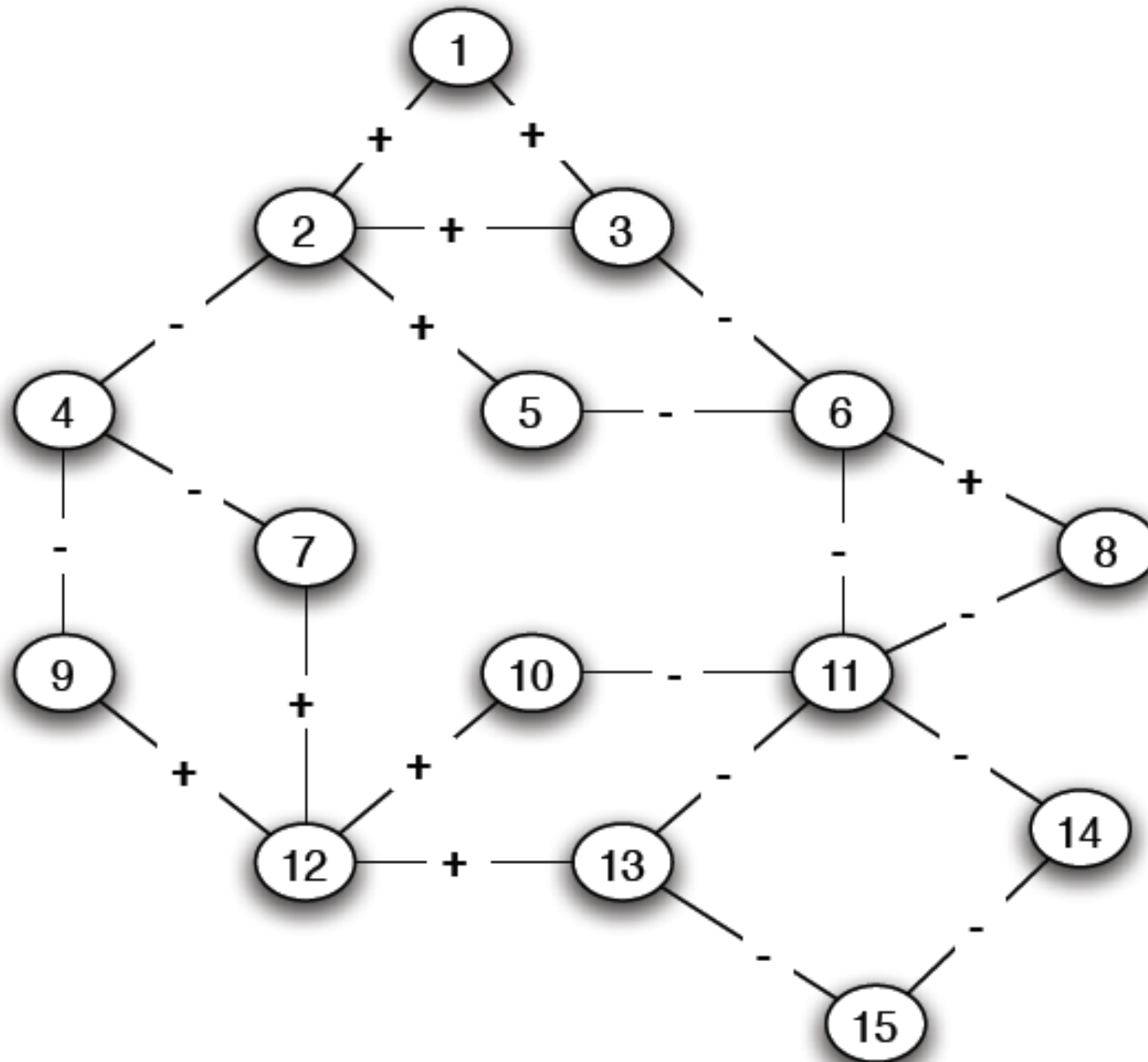
Then we can create a balanced division of the original graph, by labeling each node the way its supernode is labeled in the reduced graph.

2) find a cycle in the original graph that has an odd number of negative edges

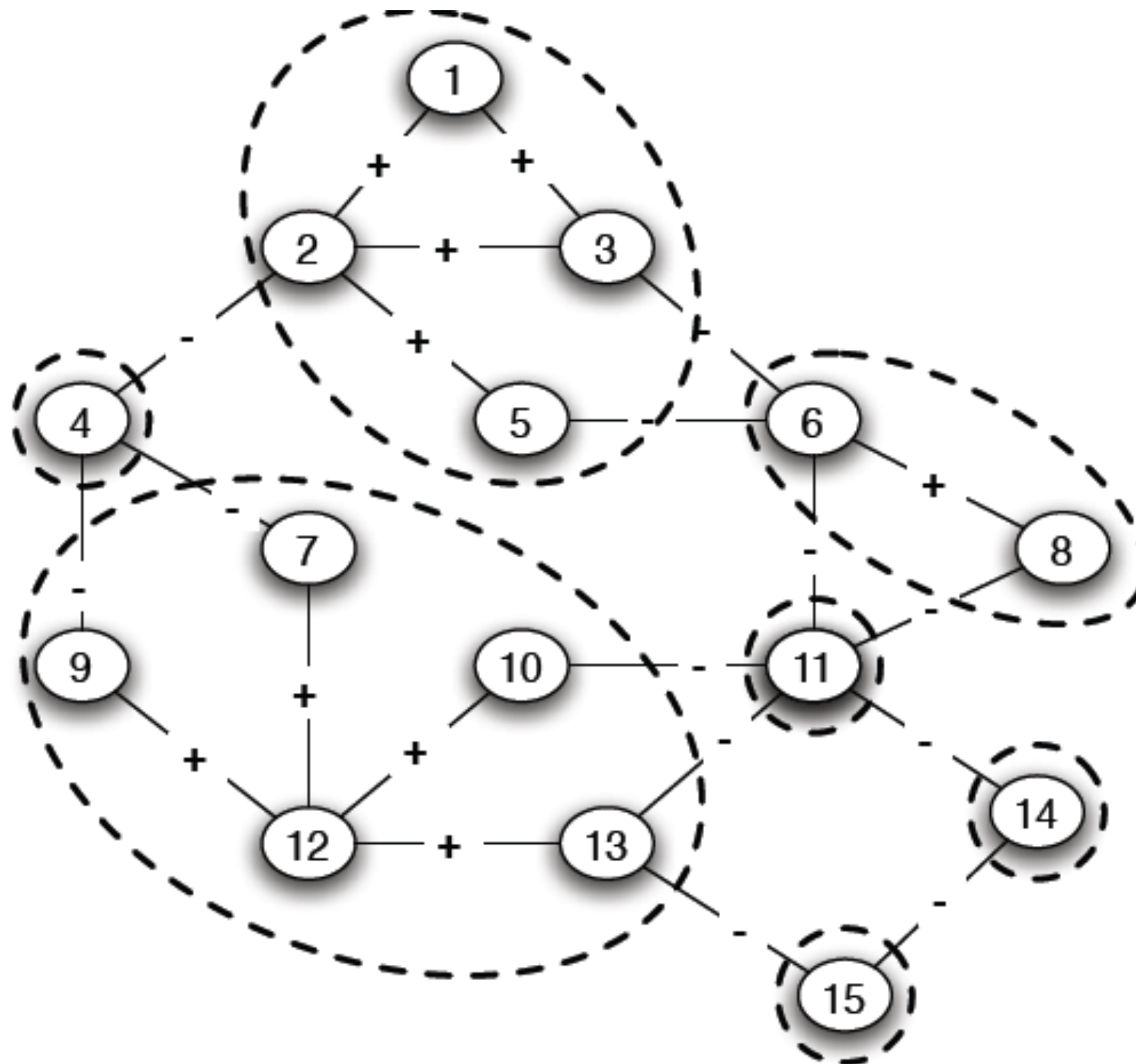
Simply “stitch together” these negative edges using paths consisting entirely of positive edges that go through the insides of the supernodes



# Signed Graph: Is it Balanced?

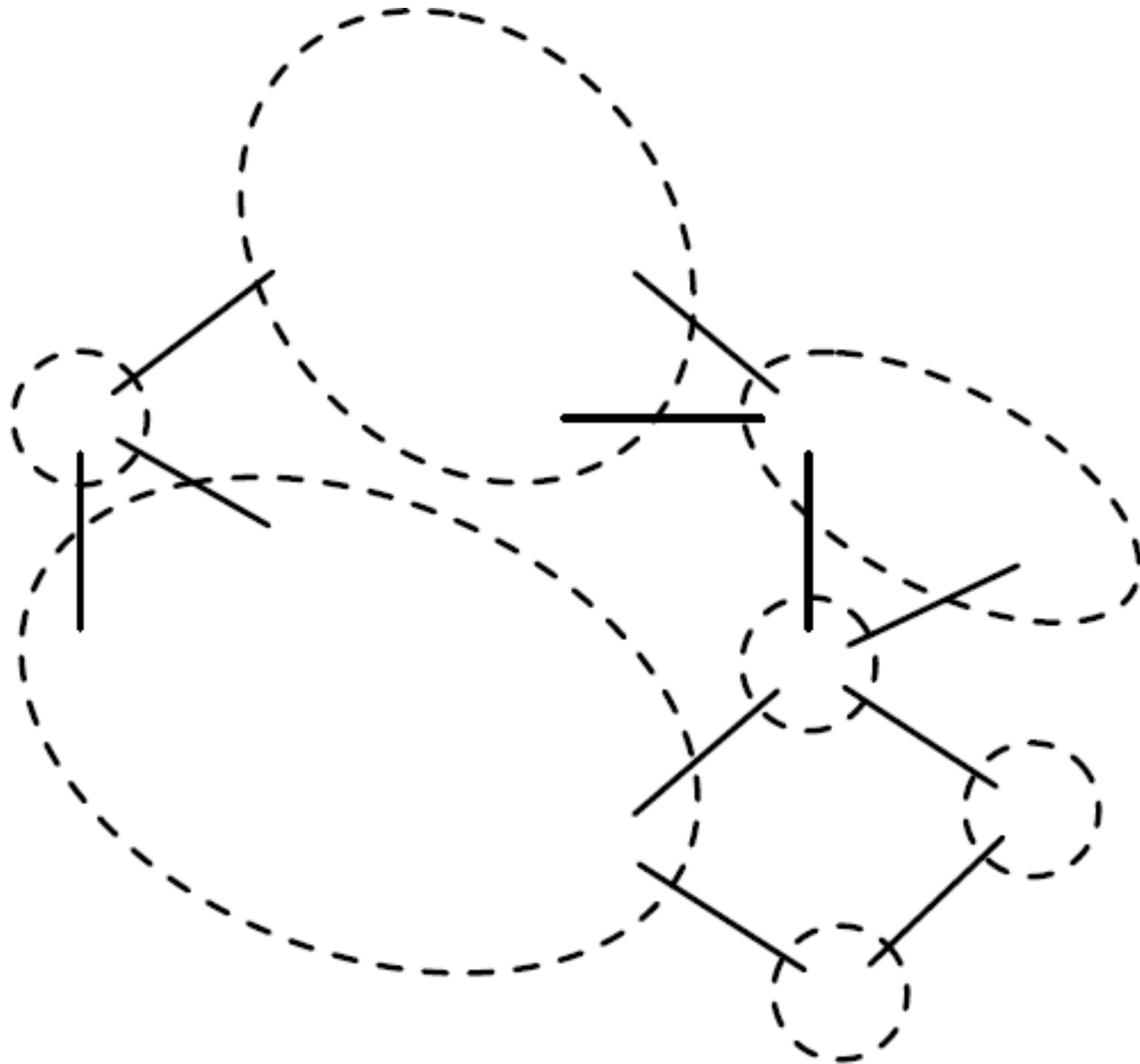


# Positive Connected Components





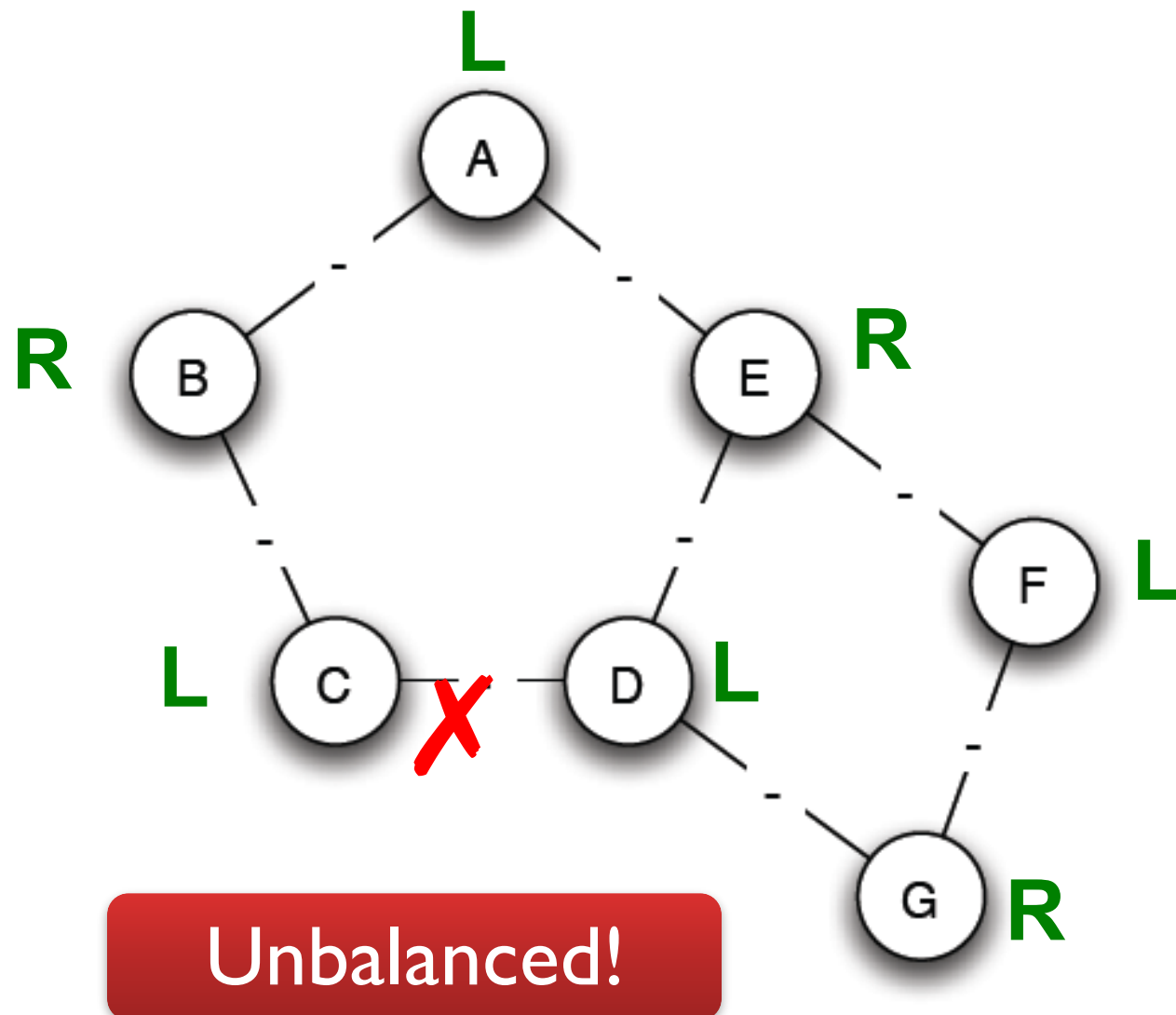
# Reduced Graph on Super-Nodes



# BFS on Reduced Graph

Using BFS assign each node a **side**

Graph is **unbalanced** if any two connected super-nodes are assigned the **same side**



# Where Do Signed Edges Come From?

**In many online applications users express positive and negative attitudes/opinions:**

- Through **actions**:

- Rating a product/person
- Pressing a “like” button

- Through **text**:

- Writing a comment, a review

- **Success of these online applications is built on people expressing opinions**

- Recommender systems
- Wisdom of the Crowds
- Sharing economy

amazon.com.



WIKIPEDIA  
The Free Encyclopedia

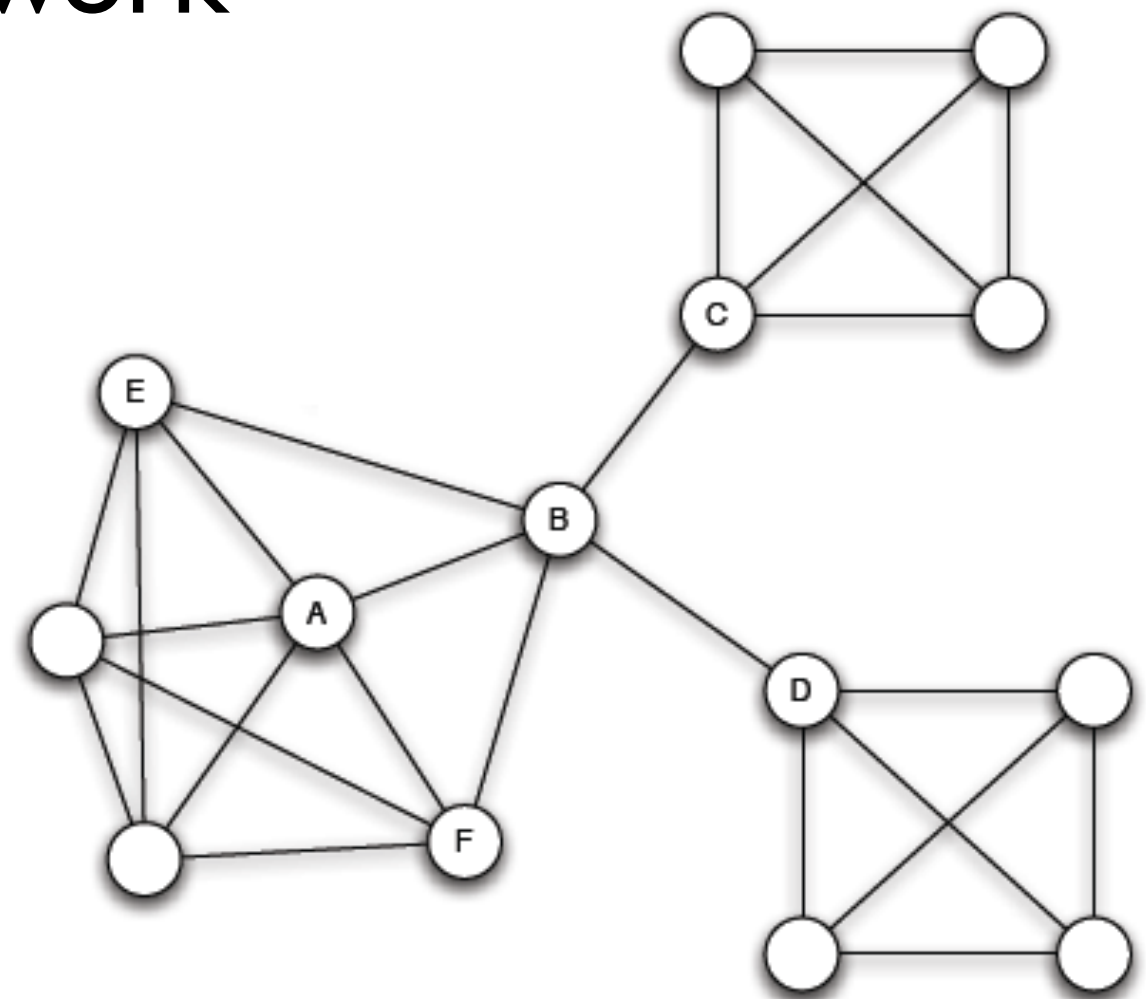


# Global Structure of Signed Nets

Intuitive picture of social network  
in terms of  
densely linked clusters

**How does structure  
interact with links?**

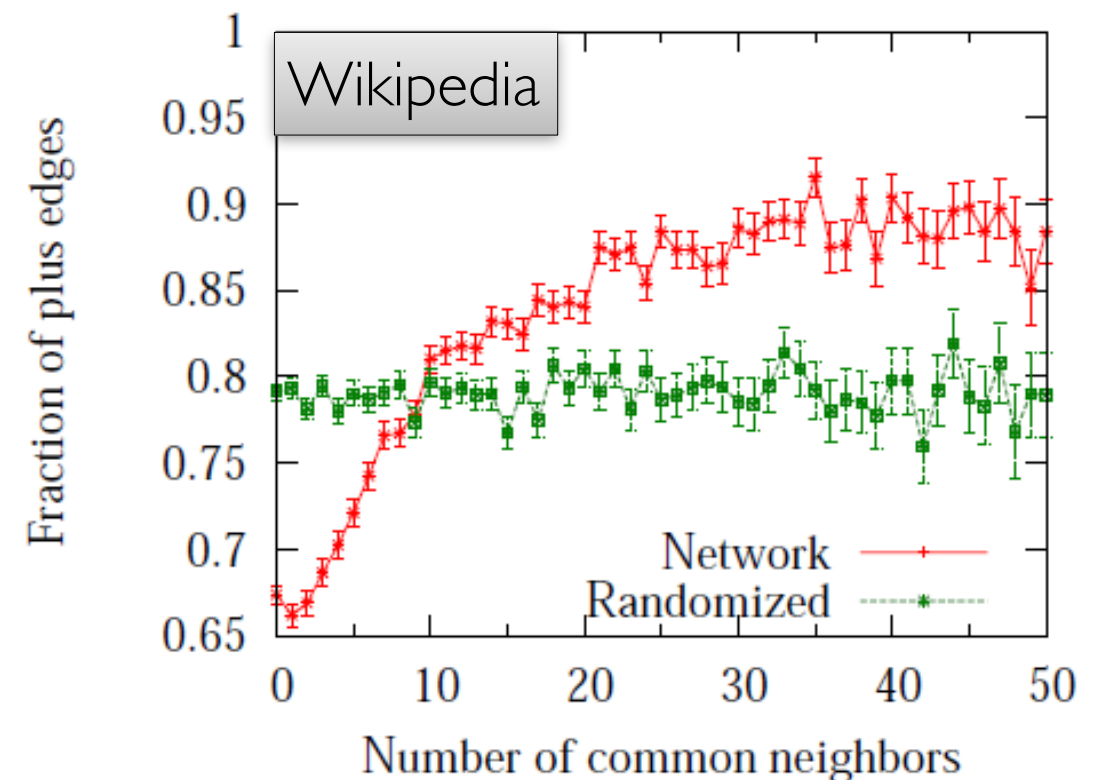
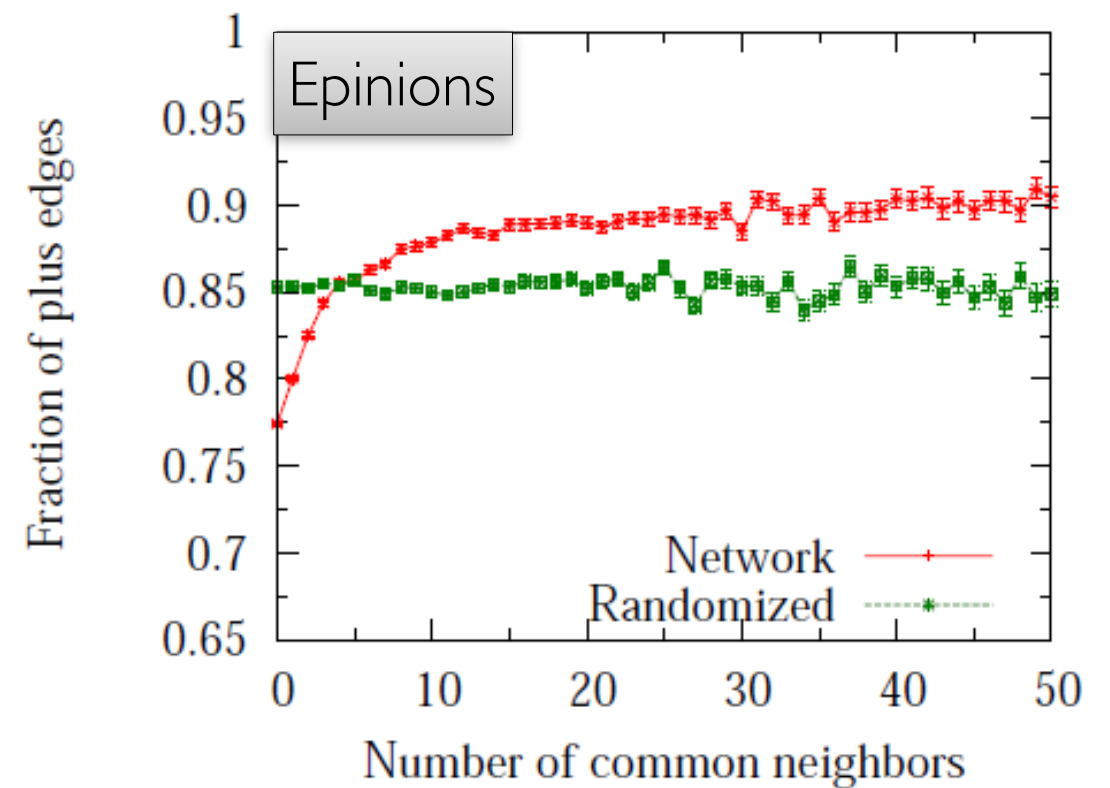
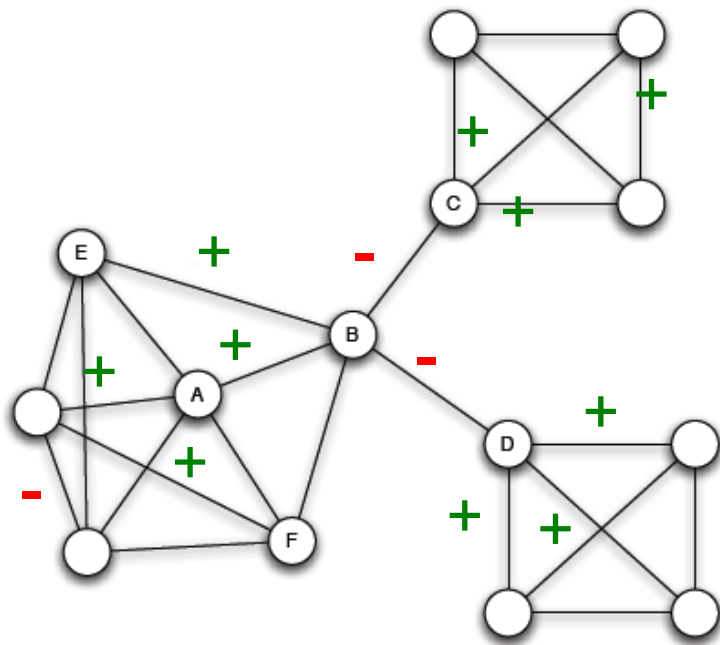
**Embeddedness of  
link (A,B):** Number of  
shared neighbors



# Global Factions: Embeddedness

## Embeddedness of ties:

Positive ties tend to be **more** embedded



# Real Large Signed Networks

Each link **A-B** is **explicitly** tagged with a sign:

**Epinions:** Trust/Distrust

Does A trust B's product reviews?  
(only positive links are visible to users)

**Wikipedia:** Support/Oppose

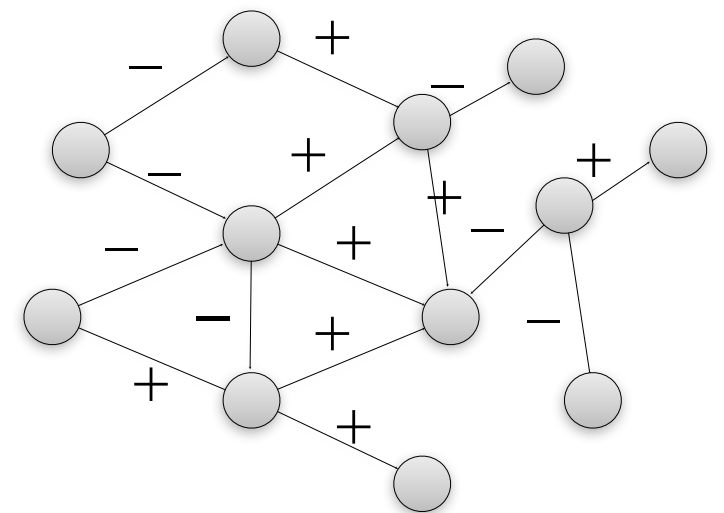
Does A support B to become  
Wikipedia administrator?

**Slashdot:** Friend/Foe

Does A like B's comments?

**Other examples:**

Online multiplayer games

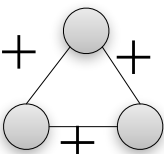
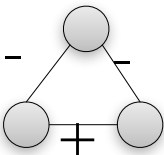
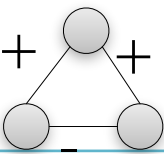
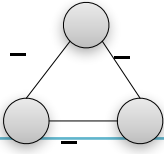


	Epinions	Slashdot	Wikipedia
Nodes	119,217	82,144	7,118
Edges	841,200	549,202	103,747
+ edges	85.0%	77.4%	78.7%
- edges	15.0%	22.6%	21.2%

# Balance in Our Network Data

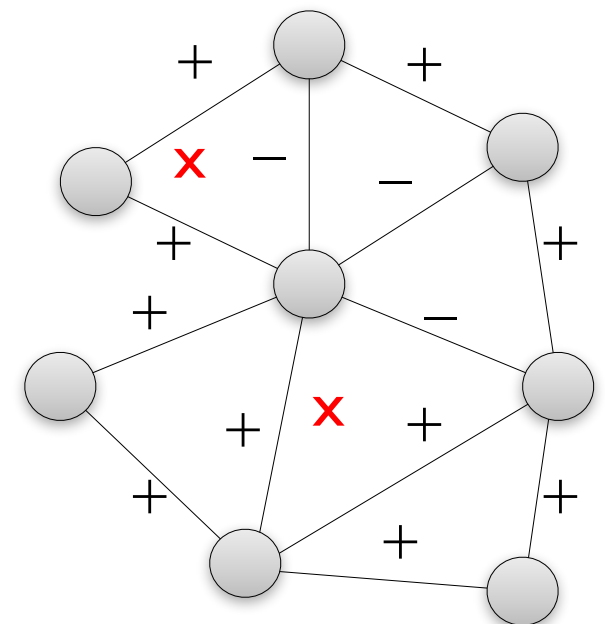
## Does structural balance hold?

Compare frequencies of signed triads in real and “shuffled” signs

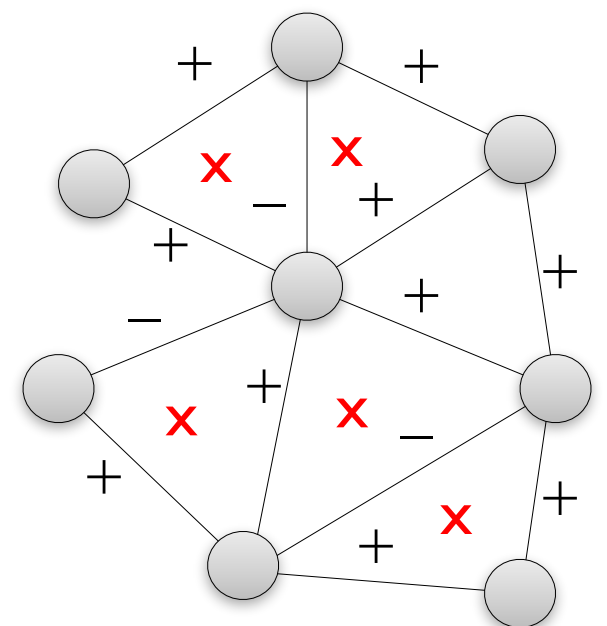
	Triad	Epinions		Wikipedia		Consistent with Balance?
		P(T)	P <sub>0</sub> (T)	P(T)	P <sub>0</sub> (T)	
Balanced		0.87	0.62	0.70	0.49	✓
		0.07	0.05	0.21	0.10	✓
Unbalanced		0.05	0.32	0.08	0.49	✓
		0.007	0.003	0.011	0.010	✗

P(T) ... fraction of a triads

P<sub>0</sub>(T)... triad fraction if the signs would appear at random



Real data

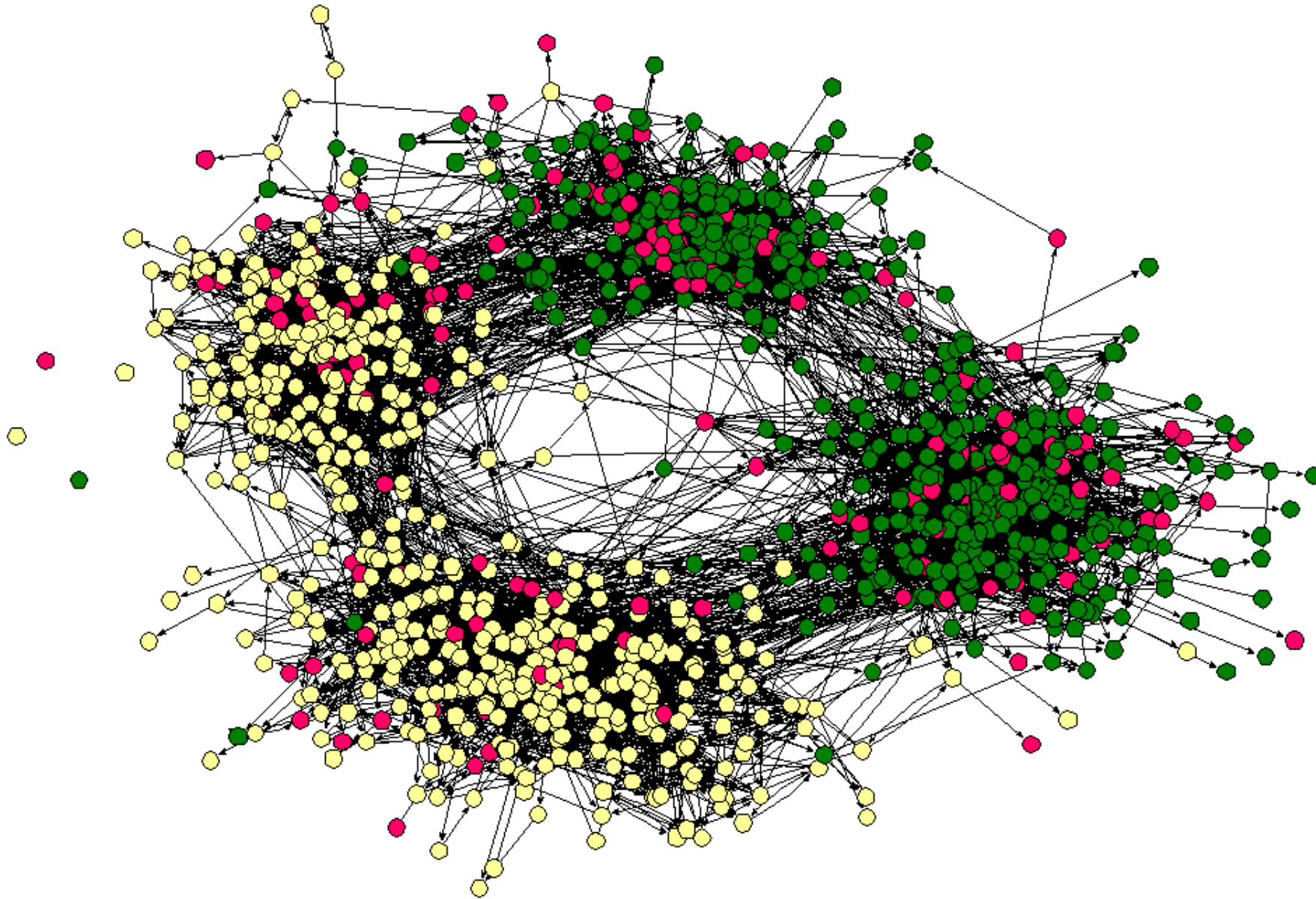


Shuffled data



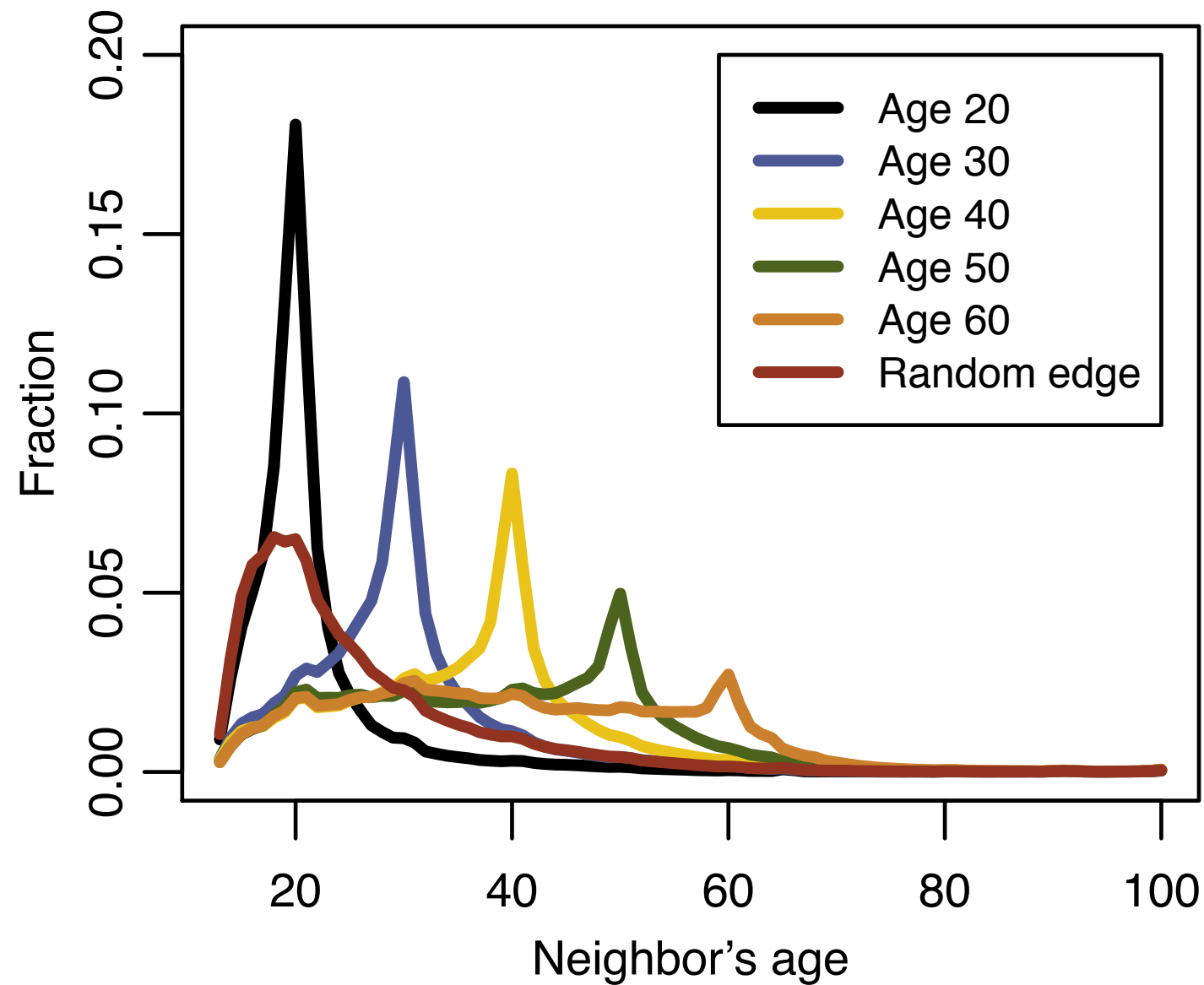
**Homophily**  
**“Birds of a Feather Flock Together”**

# Homophily



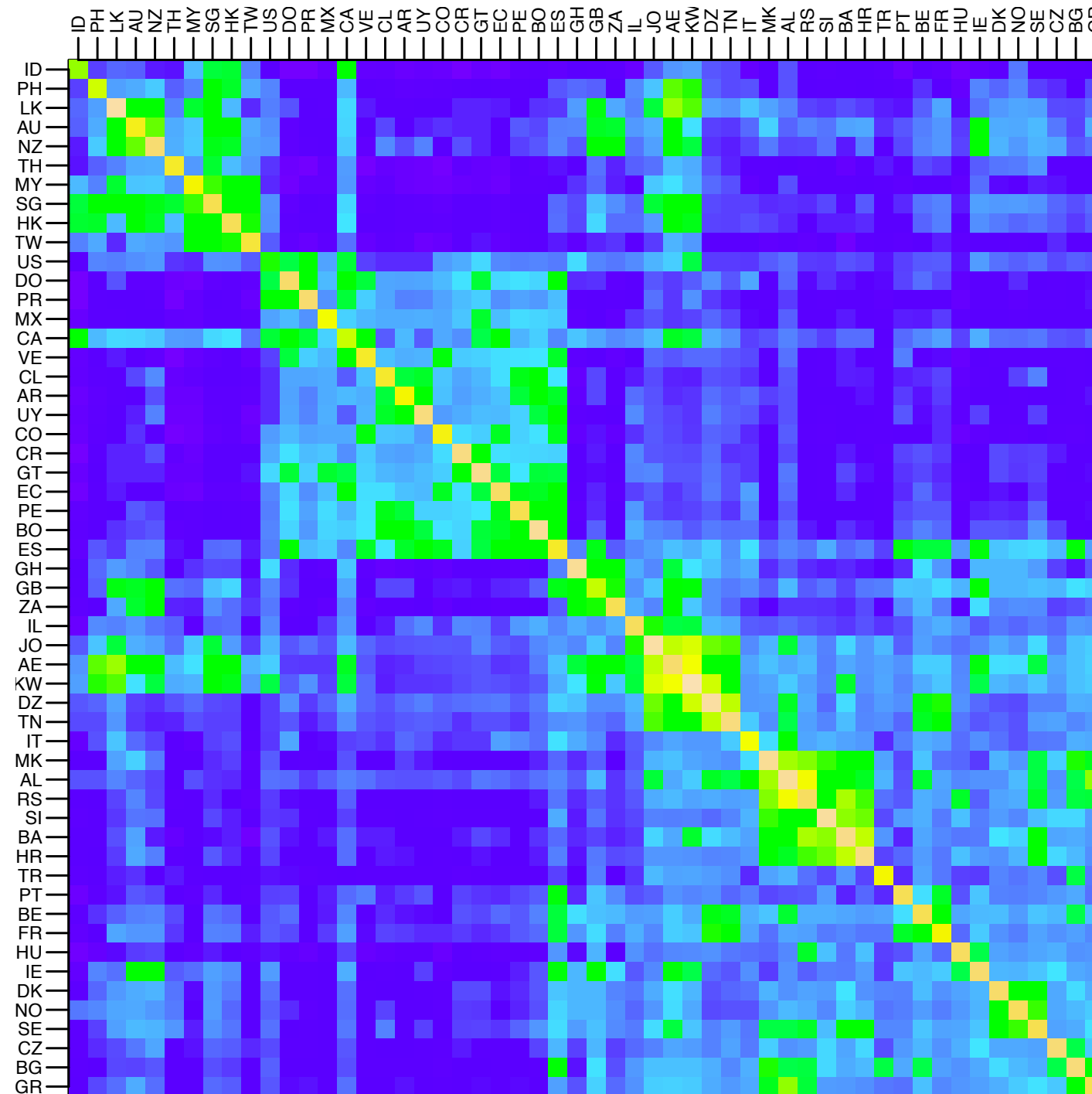
- US middle school + high school
- node color = self-identified race

# Homophily: Age



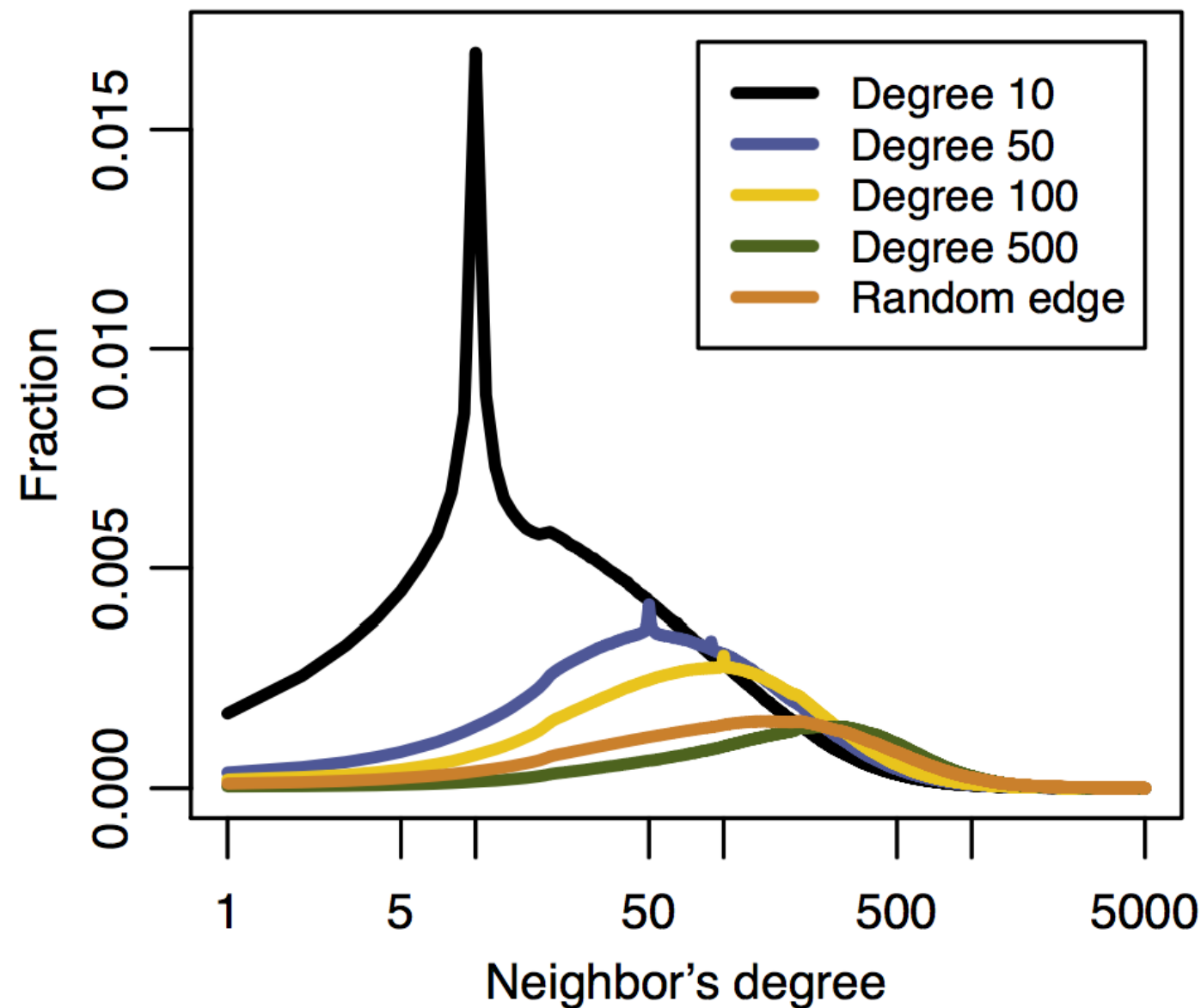
- Facebook friendship network, 2011

# Homophily: Nationality



- Facebook friendship network, 2011

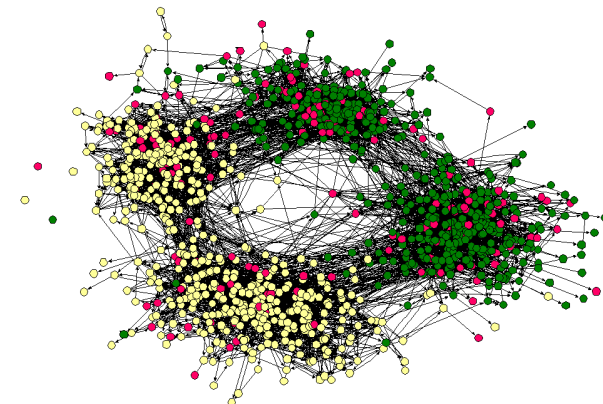
# Homophily: Friend count



- Facebook friendship network, 2011

# Homophily

- Connections don't form uniformly at random
- **Null model**: what if they were forming at random?
- **Measuring homophily**: are there fewer connections between nodes across traits than you'd expect at random?
- **Homophily test**: If the fraction of cross-gender edges is significantly less than at random, then there is evidence of homophily.



# Homophily

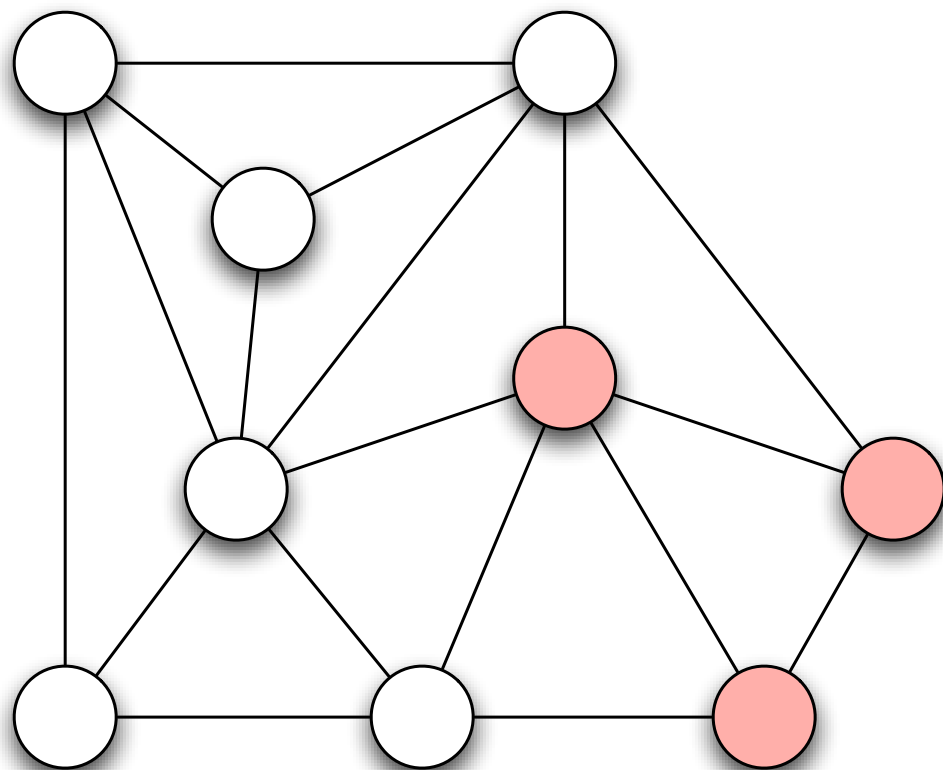
$p$  = Probability that a node is white

$q$  = Probability that a node is red

Prob an edge is between two white nodes?

Prob an edge is between two red nodes?

Prob an edge is between 1 red, 1 white?



Homophily test:



# Homophily

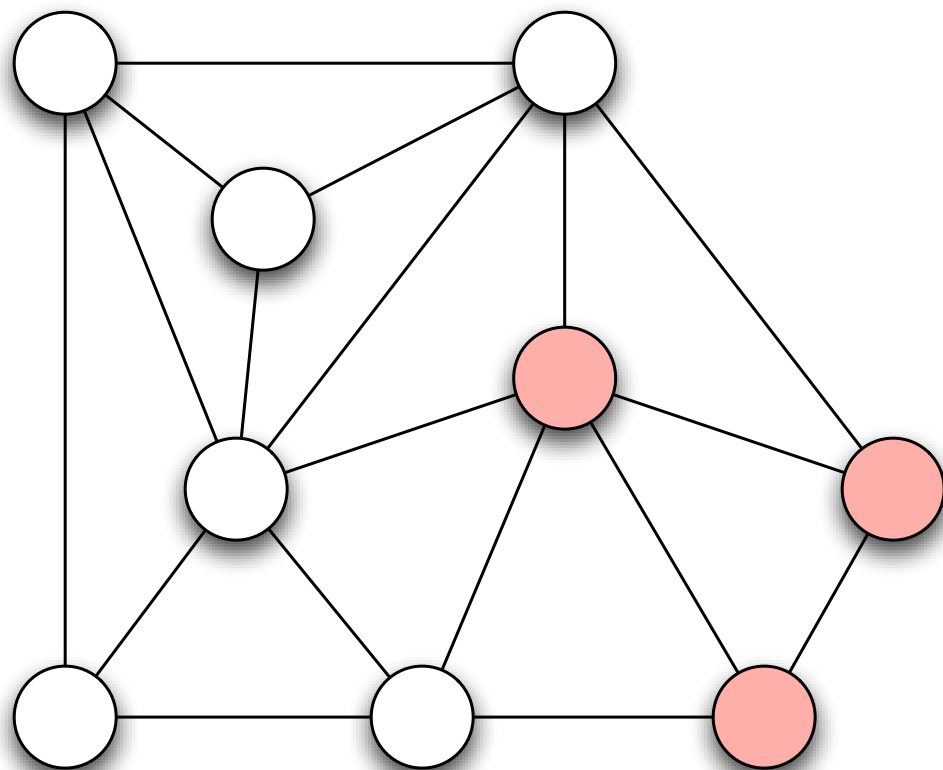
$p$  = Probability that a node is white  $6/9=2/3$

$q$  = Probability that a node is red  $3/9=1/3$

Prob an edge is between two white nodes?  $p^2$

Prob an edge is between two red nodes?  $q^2$

Prob an edge is between 1 red, 1 white?  $2pq$



Homophily test:  $2pq = 4/9 = 8/18$

Observed: 5/18