

CSCC46H, Fall 2025 Lecture 11

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# Today

A3 due next Tuesday, Dec 2 @ I I:59pm, plus no-penalty extension to next Friday, Dec 5 @ I I:59pm (late days not needed)

# Today

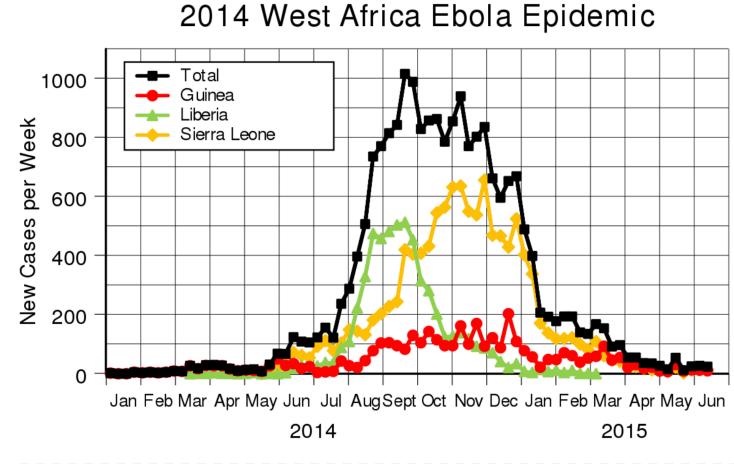
Epidemics and Contagion Voting

Why study epidemics in a computer science class?

Epidemics are diseases that travel socially

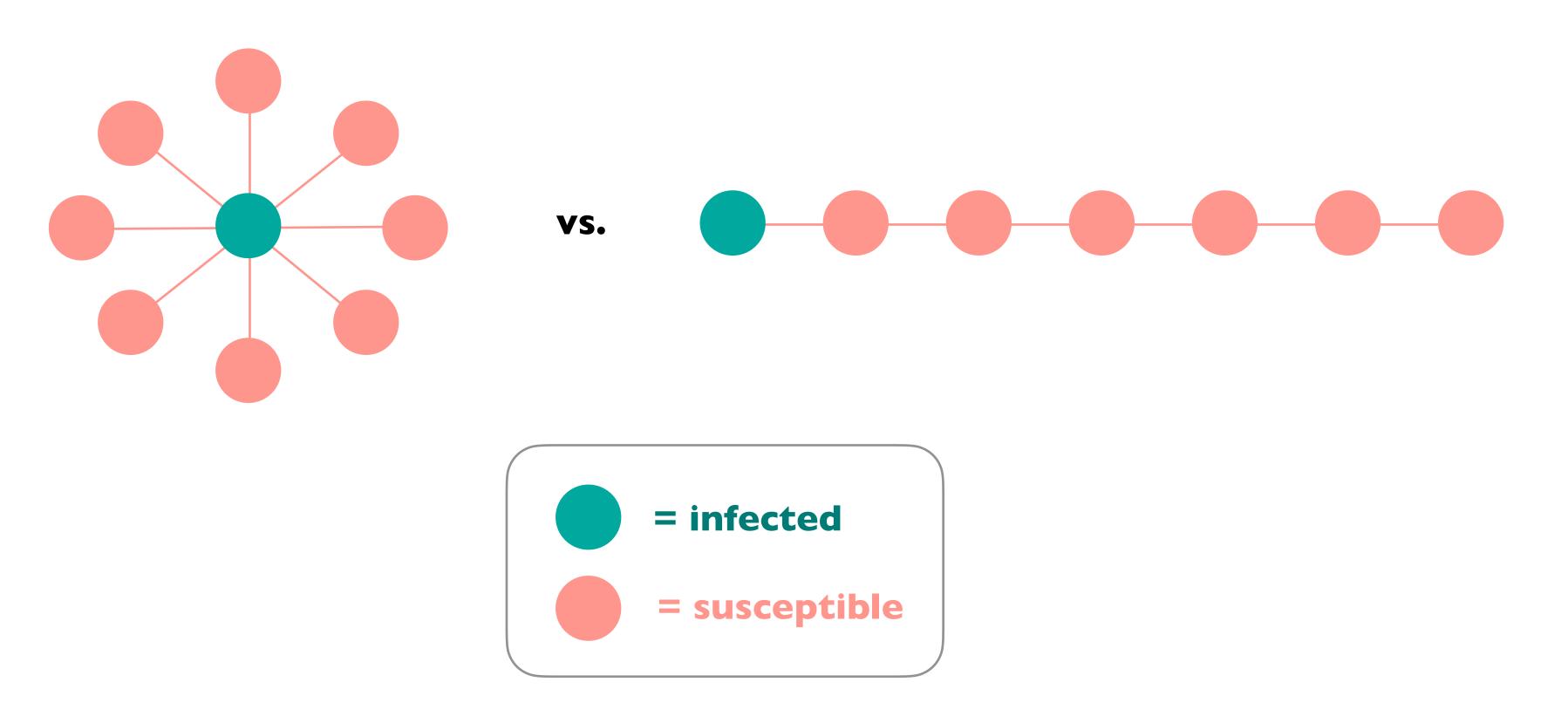
The structure of social interaction networks determine the spread of disease







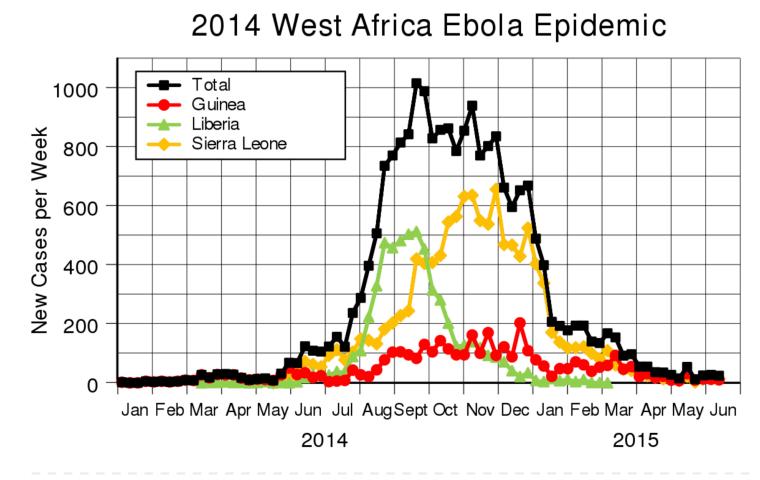
Which outbreak is more dangerous to the population?



#### Types of epidemic diffusion:

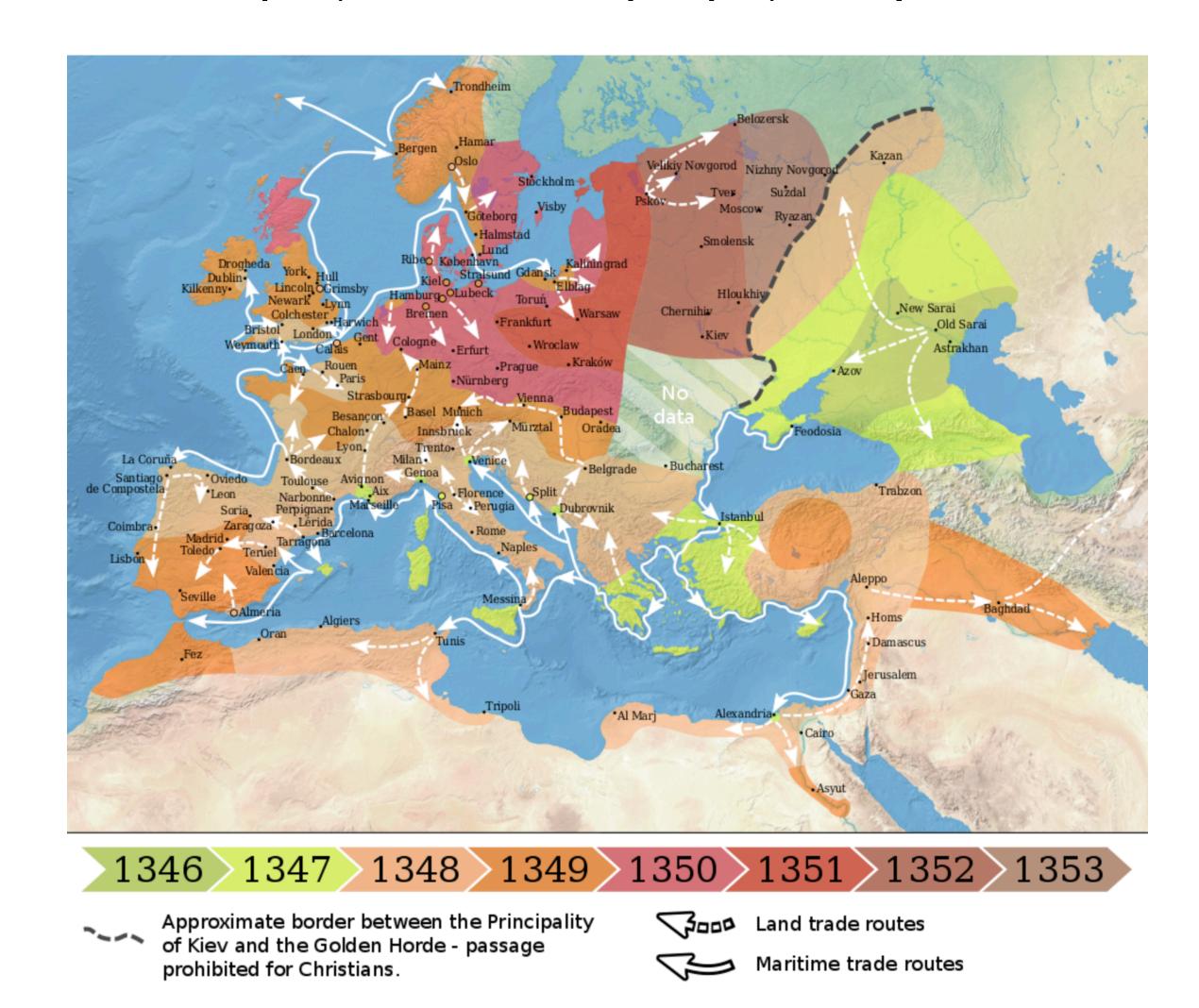
- Explosive spread through a population
- "Slow burn" persistence over long periods of time
- Wave-like cyclical patterns



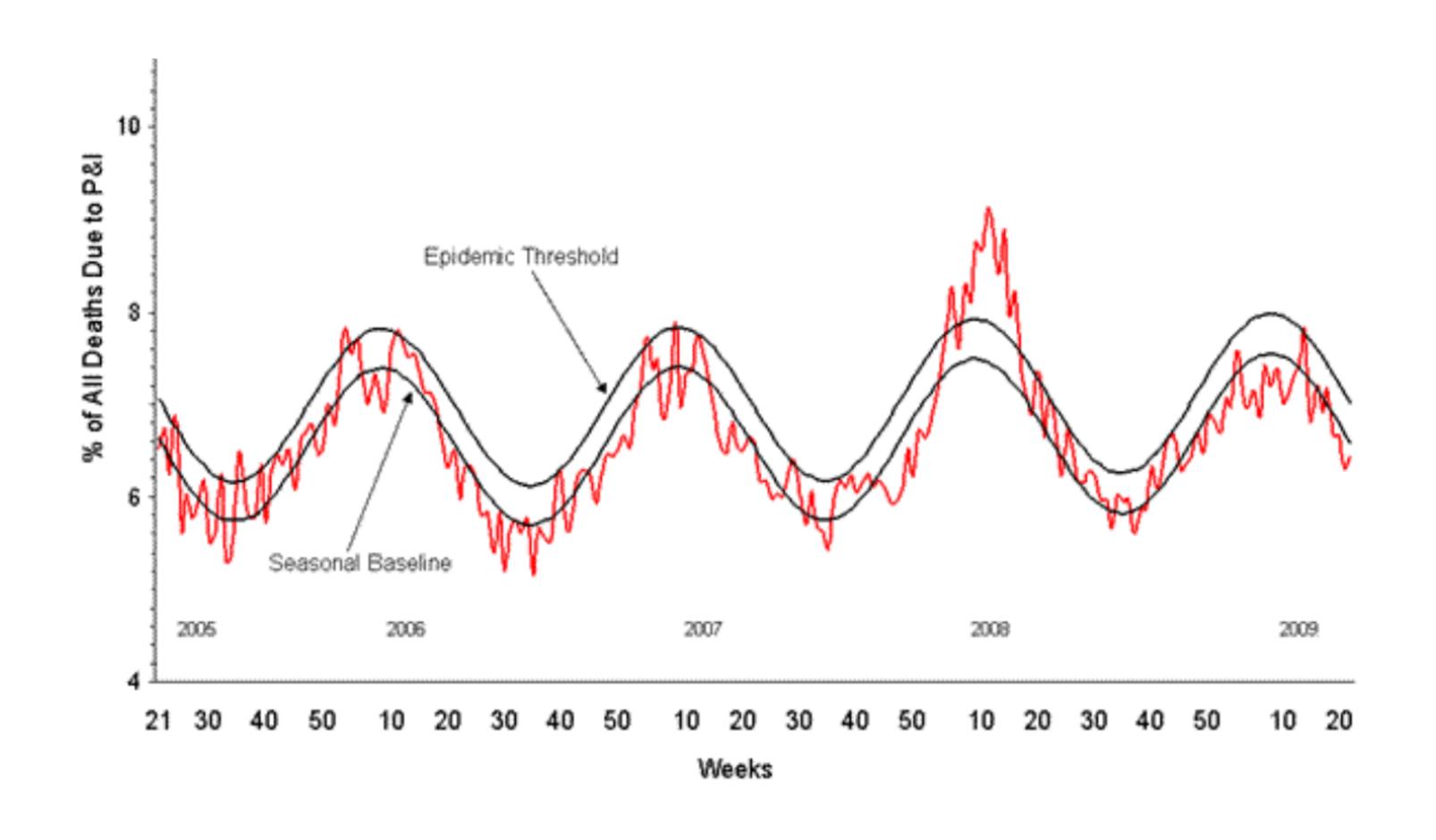




**Explosive spread**: Bubonic Plague (the "Black Death"): wiped out ~50% of the population in Europe (~150 million people) in 7 years



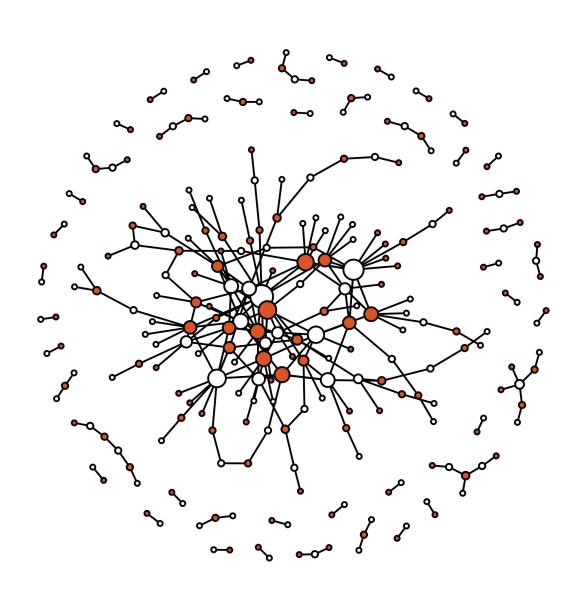
#### Other epidemics are cyclical



What determines how an epidemic might spread?

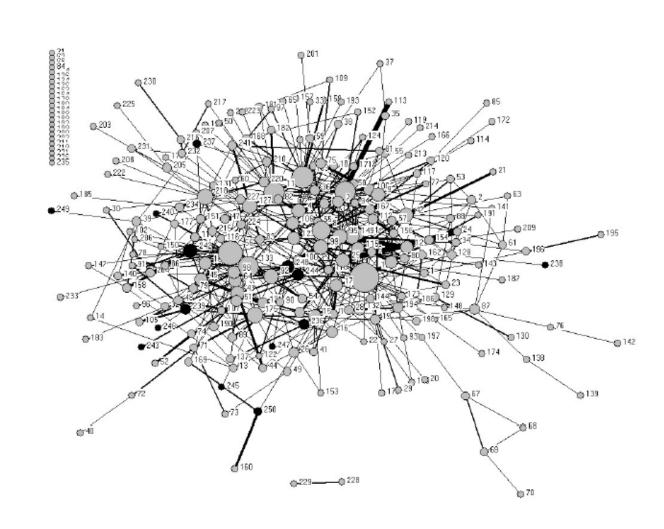
- Properties of the disease
- Structure of the network

What network?

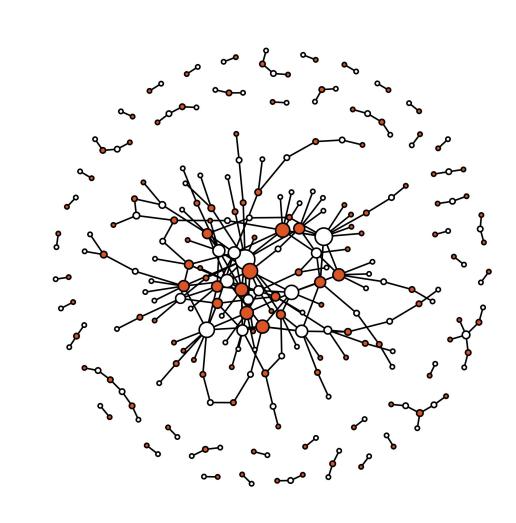


Node for each person

**Edge** if two people come into contact with each other in a way that makes it possible for a disease to spread



UK fish farm exchanges

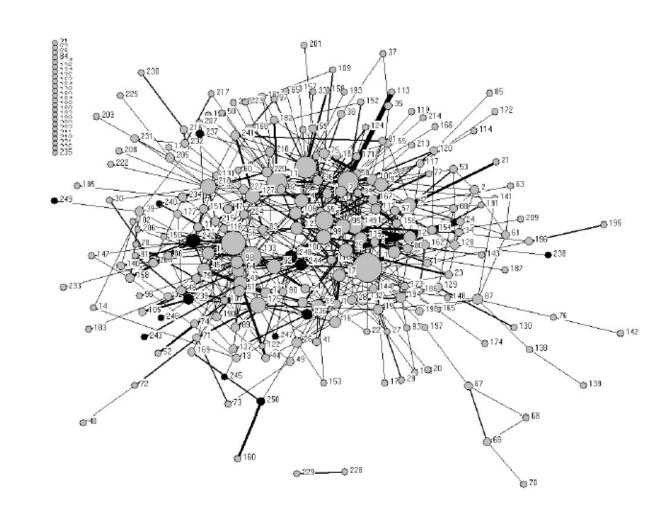


CH sexual contact network

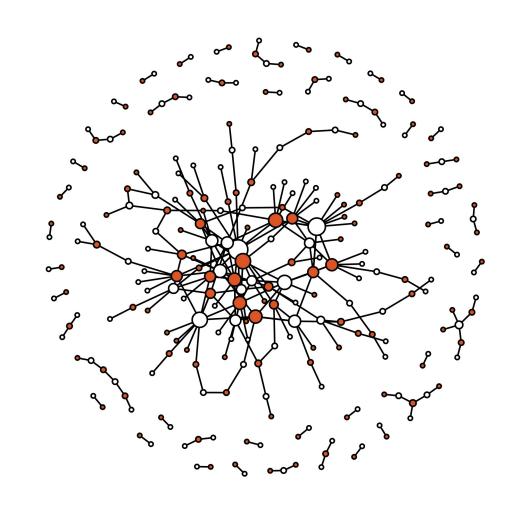
Once you've got through the laborious process of mapping out the contact network, can you use it to study any disease?

No! Definition of "contact" depends on the disease

- Airborne transmission: edge between everyone who was in the same car, etc.)
  - many edges
- Close contact / sexual transmission: sparser graph



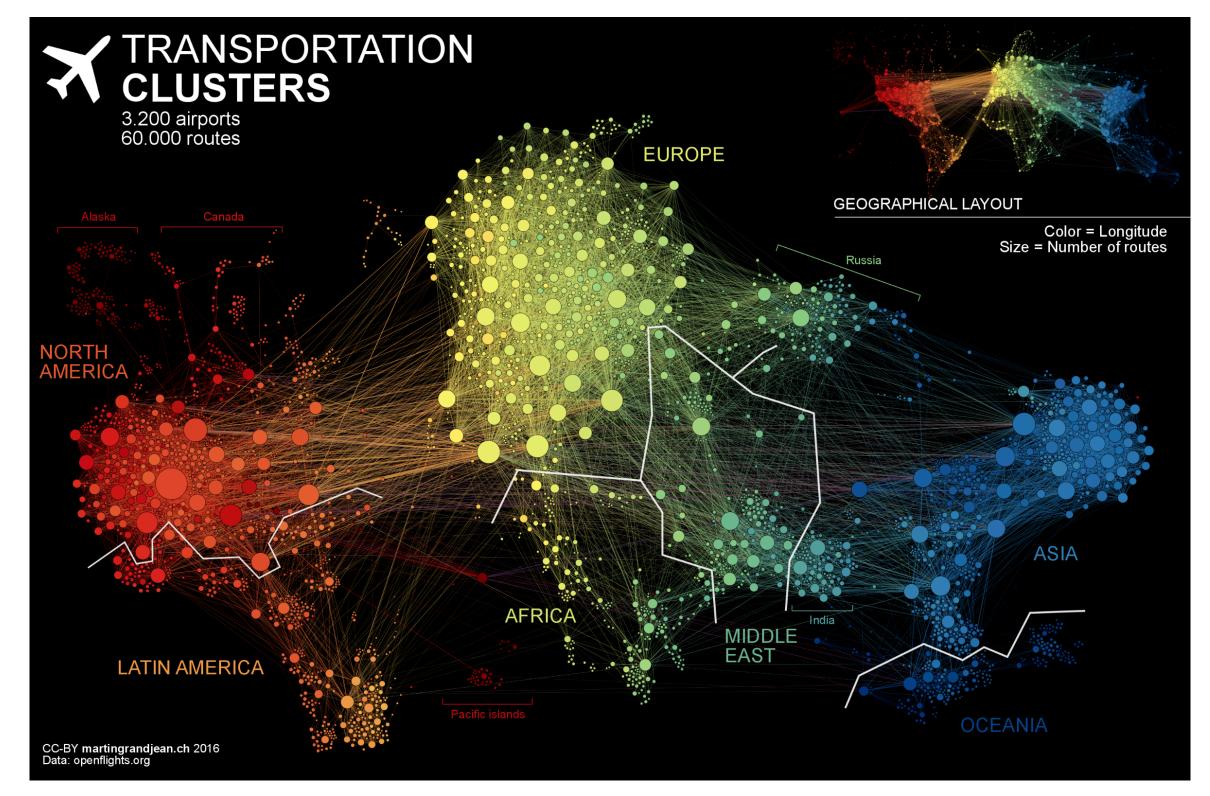
UK fish farm exchanges



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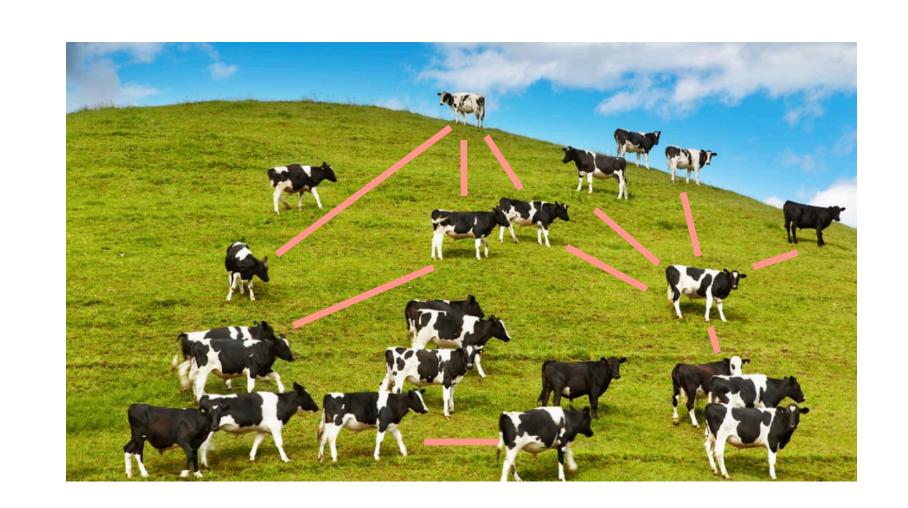
Big part of real-world epidemic research is constructing contact networks

Lots of work on travel patterns in cities, the worldwide airline network, etc. to understand how diseases can spread in today's world



Not just human contact networks

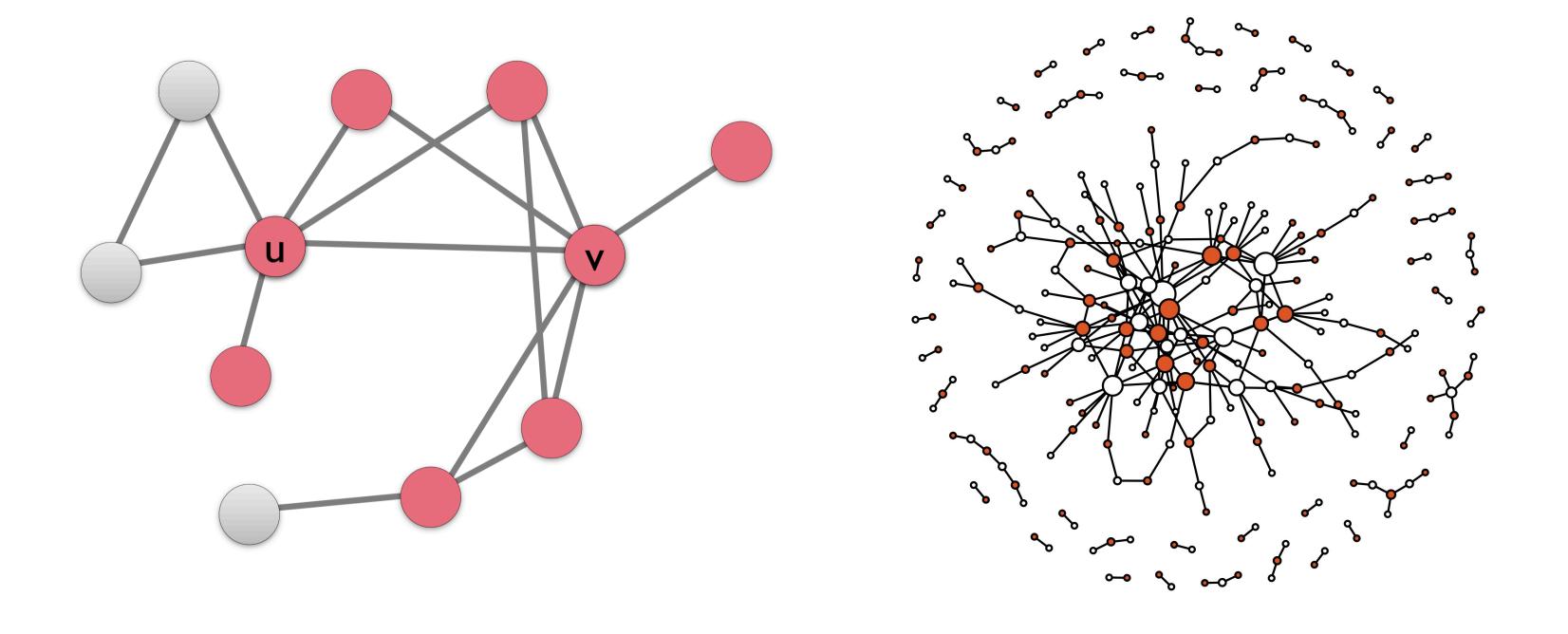
Animal/livestock networks and plant networks





### Behavioural vs. Biological Contagion

Biological/epidemic diffusion: no decision-making!



Decision cascade

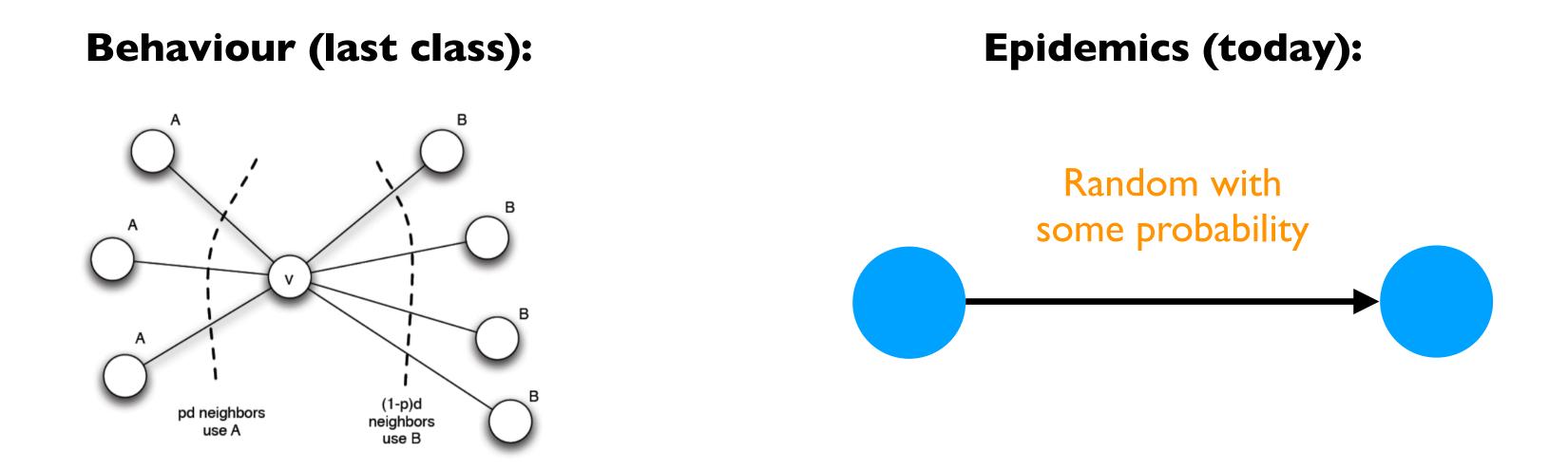
High school contact network

### Modeling Epidemic Diffusion

Biggest difference: model transmission as random

No decision-making, but also the processes by which diseases spread from one person to another are so complex and unobservable at the individual level that it's most useful to think of them as random

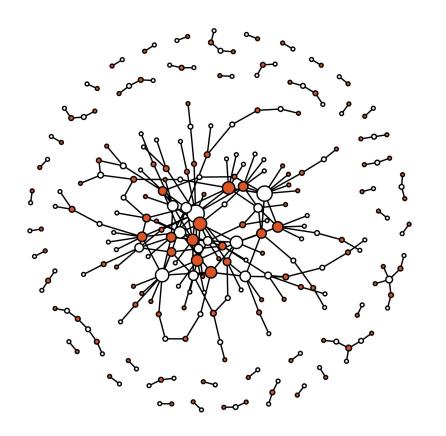
Use randomness to abstract away difficult biological questions about the mechanics of spread



### Modeling Epidemic Diffusion

#### Basic structure of epidemic diffusion:

- Someone gets infected
- Then they infect some number of people
- Those people infect others

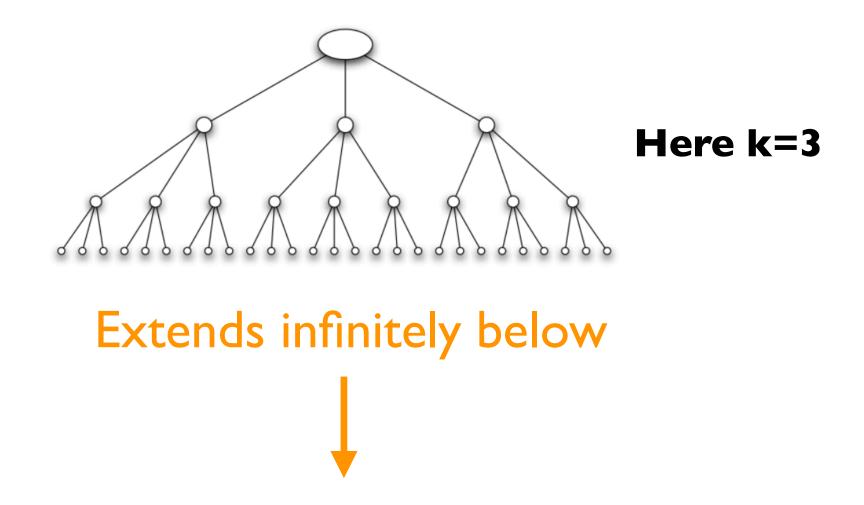


Model as a random process on a tree:

Wave I: First person infected, infects each of k neighbors with independent probability p

**Wave 2:** For each infected person, they infect each of k neighbors with independent probability p

Wave 3+: repeat for each infected person

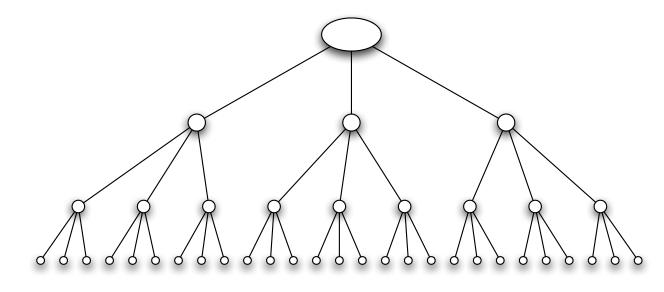


#### Model parameters:

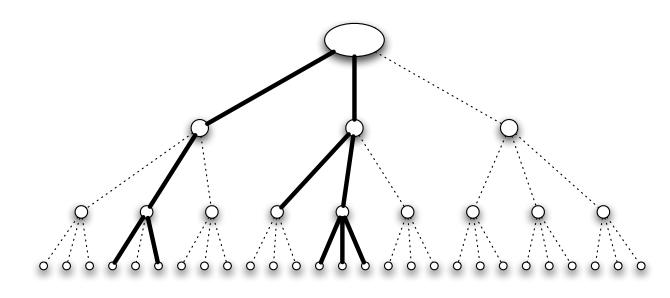
k: number of individuals each person can possibly infect:

Higher transmission probability p:

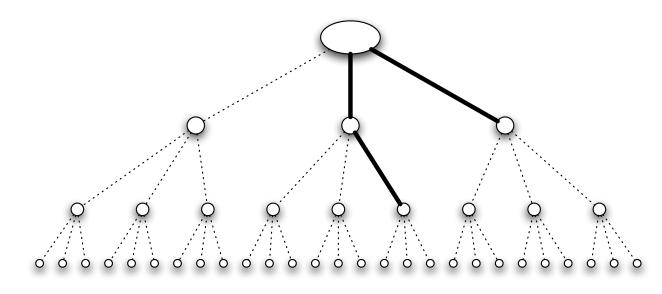
Lower transmission probability p:



(a) The contact network for a branching process



(b) With high contagion probability, the infection spreads widely



(c) With low contagion probability, the infection is likely to die out quickly

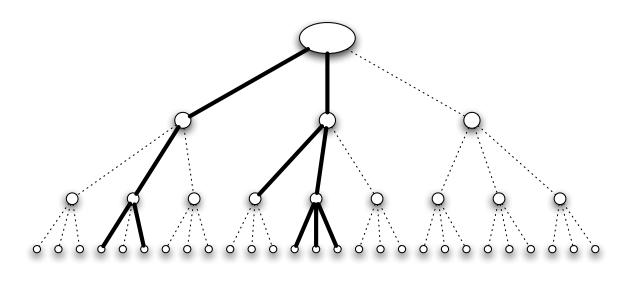
#### **Branching Process: Outcomes**

Only two possibilities in the long run: **blow up** or **die out** How does it die out?

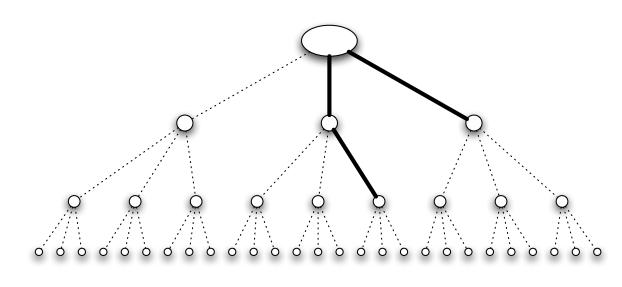
Dies out if and only if none of the nodes on a given level are infected

Disease might blow up:

Disease has already died out:



(b) With high contagion probability, the infection spreads widely

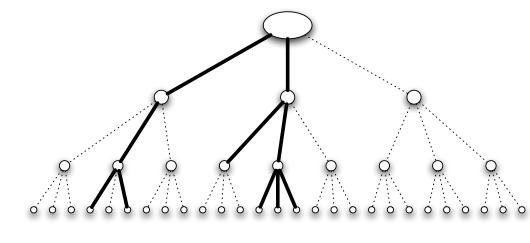


(c) With low contagion probability, the infection is likely to die out quickly

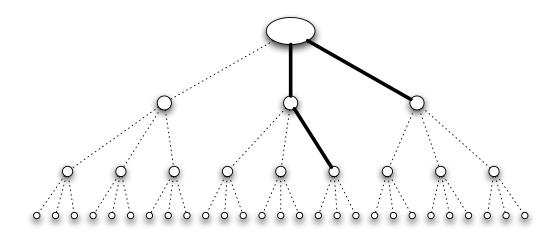
Only two possibilities in the long run: blow up or die out How does it die out?

Dies out if and only if none of the nodes on a given level are infected

Define Basic reproductive number R<sub>0</sub>: the number of expected new cases caused by an individual



(b) With high contagion probability, the infection spreads widely



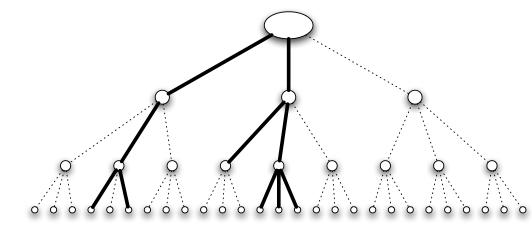
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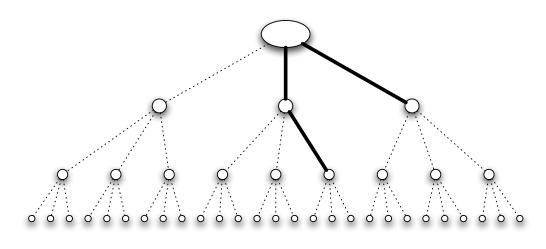
Dies out if and only if none of the nodes on a given level are infected

Define Basic reproductive number R<sub>0</sub>: the number of expected new cases caused by an individual

$$R_0 = pk$$



(b) With high contagion probability, the infection spreads widely



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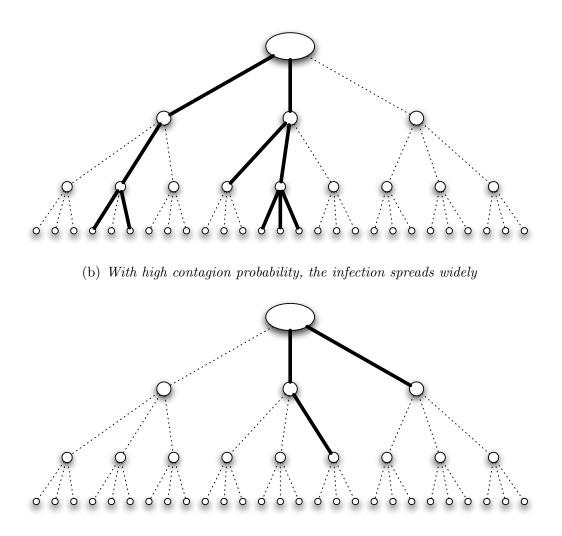
### Branching Process: Ro

<u>Claim</u>: Epidemic spread in the branching process model is entirely controlled by the reproductive number  $R_0$ :

- If  $\mathbf{R_0} < \mathbf{I}$  then with probability I the disease dies out after a finite number of steps.
- If  $\mathbf{R_0} > \mathbf{I}$  then with probability > 0 the disease persists by infecting at least one person in each wave.

"Go big or go home."

 $R_0 = pk$ 



(c) With low contagion probability, the infection is likely to die out quickly

### Branching Process: Ro

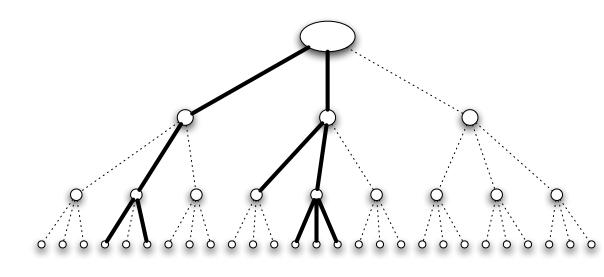
#### $R_0 = pk < 1$ :

With probability I the disease dies out after a finite number of steps Below replacement; disease isn't able to replenish itself. Even if it grows momentarily, it trends downward.

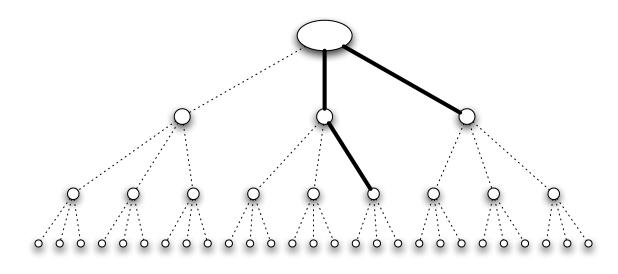
#### $R_0 = pk > 1$ :

with probability > 0 the disease persists by infecting at least one person in each wave

Always trending upward. Could still get "unlucky" and die out, but there's a non-zero chance it runs forever.



(b) With high contagion probability, the infection spreads widely



(c) With low contagion probability, the infection is likely to die out quickly

### Branching Process: Ro

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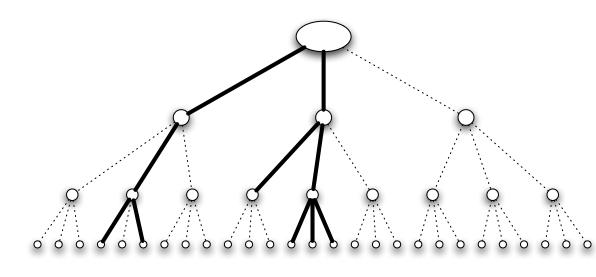
Even if it grows momentarily, it trends downward.

#### $R_0 = pk > 1$ :

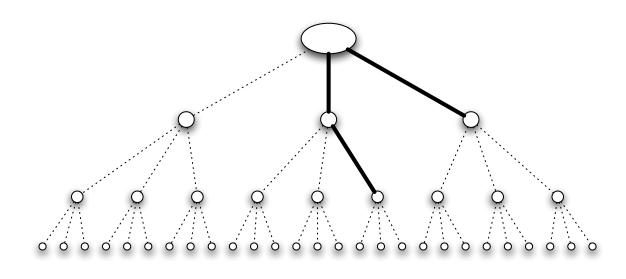
with probability > 0 the disease persists by infecting at least one person in each wave

Always trending upward. Could still get "unlucky" and die out, but there's a non-zero chance it runs forever.

What happens when  $\mathbf{p}$  or  $\mathbf{k}$  change near  $\mathbf{pk}=1$ ?



(b) With high contagion probability, the infection spreads widely

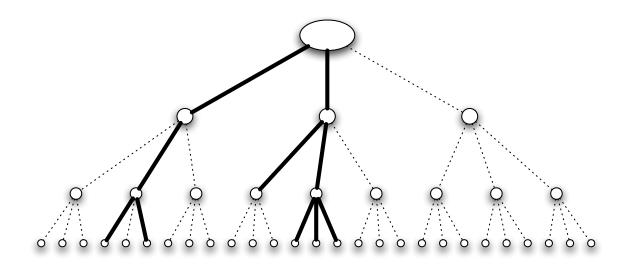


(c) With low contagion probability, the infection is likely to die out quickly

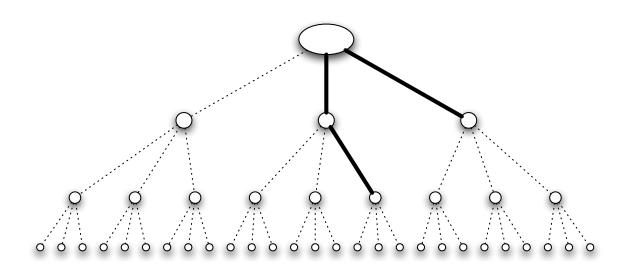
## Sensitivity of p and k

Because epidemics have a "critical threshold", it can be worth it to do a lot of work or expend resources to push **p** or **k** down a little bit.

Quarantine = reduce **k**Improved sanitation = reduce **p** 



(b) With high contagion probability, the infection spreads widely



(c) With low contagion probability, the infection is likely to die out quickly

Disease	Transmission	R <sub>0</sub>
Discase	Halisillission	110
Measles	Airborne	12–18
Diphtheria	Saliva	6-7
Smallpox	Airborne droplet	5–7
Polio	Fecal-oral route	5–7
Rubella	Airborne droplet	5–7
Mumps	Airborne droplet	4–7
HIV/AIDS	Sexual contact	2–5
Pertussis	Airborne droplet	5.5 <sup>[2]</sup>
SARS	Airborne droplet	2-5 <sup>[3]</sup>
Influenza	Airborne droplet	2–3 <sup>[4]</sup>
(1918 pandemic strain)		
Ebola	Bodily fluids	1.5-2.5 <sup>[5]</sup>
(2014 Ebola outbreak)		

COVID-19: ~2 COVID Delta: ~5 COVID Omicron: ~8

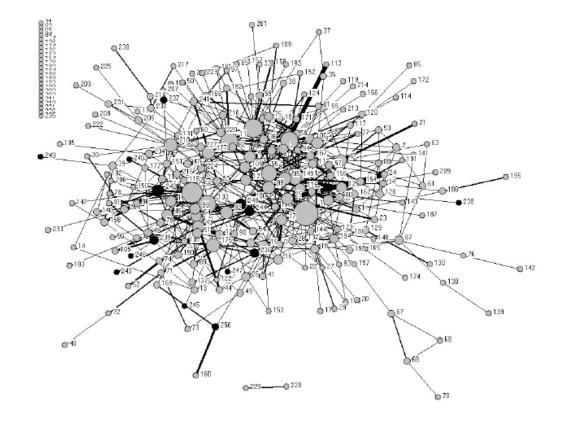
# General Models of Contagion

### **Epidemics on General Graphs**

We just studied epidemics as ideal trees

But of course real-life networks are more complicated than that

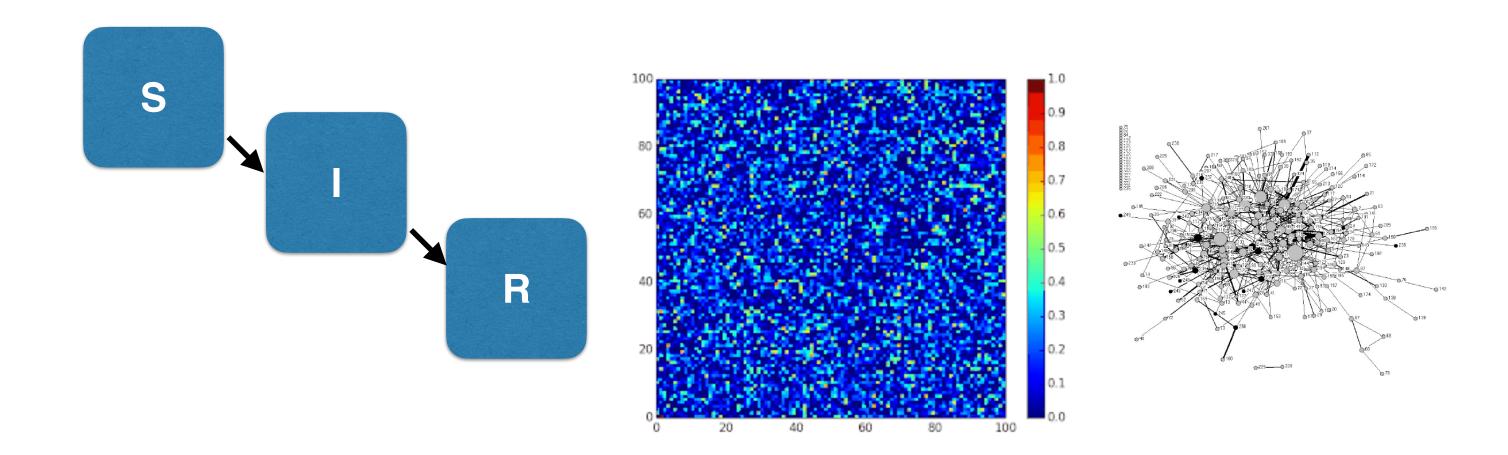
What does epidemic diffusion look like in general graphs?



### SIR Epidemic Models

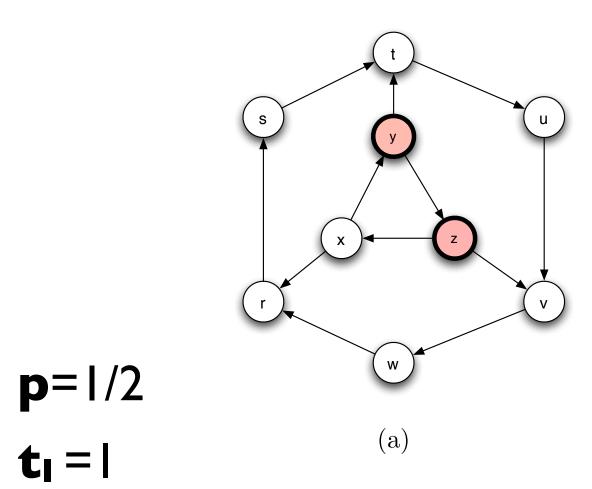
Simple lifecycle model with three stages:

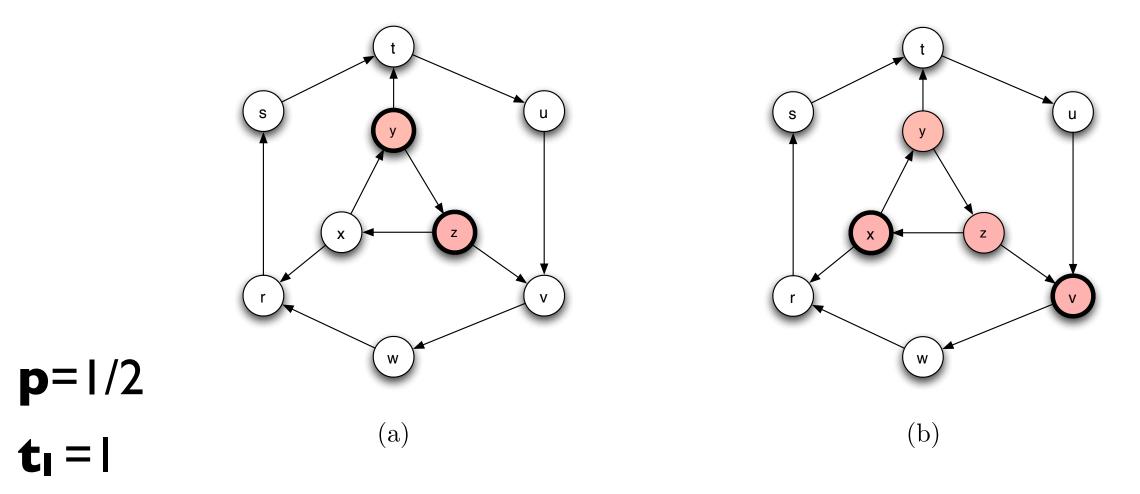
- S = Susceptible
- I = Infectious: node is infected and infects with prob p
- R = Removed: after t<sub>I</sub> time, no longer infected or infectious

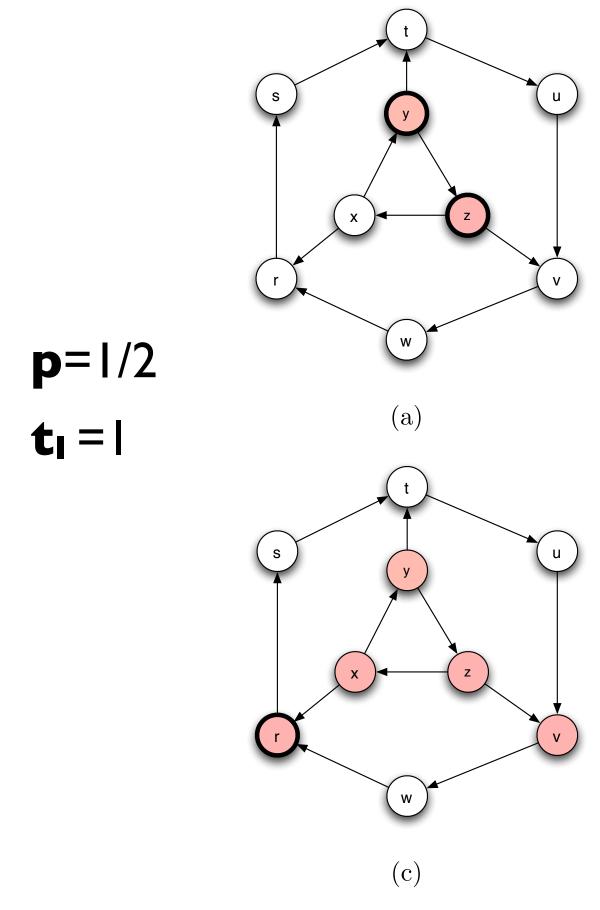


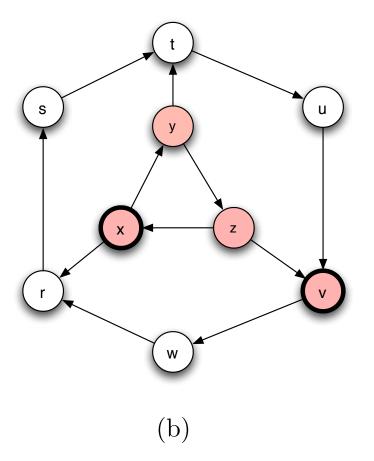
### SIR Epidemic Models

- **S** = Susceptible
- I = Infectious: node is infected and infects with prob **p**
- R = Removed: after t<sub>1</sub> time, no longer infected or infectious
- Initially some nodes in I state, rest in S state.
- Each node in I state remains infected for t<sub>I</sub> time steps
- During each step, each node has probability **p** of infecting each susceptible neighbour
- After t<sub>I</sub> time steps, no longer S nor I; removed to R

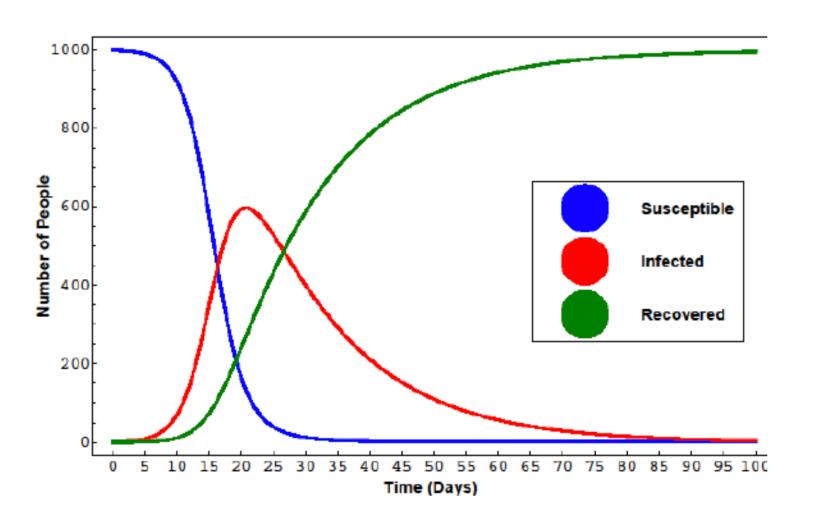








Typical run of SIR on a graph representing a contact network



**Big questions in epidemiology**: how many will an epidemic infect? How will the spread change with changes in parameters? Based on that, what are best defences?

#### SIR Epidemic Extensions

Many extensions to accommodate different parameters

Some contacts more likely than others:

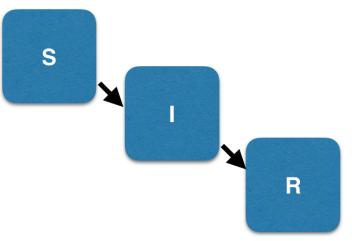
→ probability **p**<sub>uv</sub> that is pair-dependent

Disease goes through different stages (infectious incubation, then less infectious symptomatic transmission):

→ SEIR or S"III"R: either Exposed state or several different infectious states (with different p's or t's)

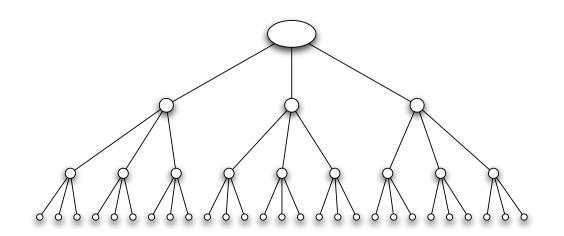
**SIS:** later in the lecture

Mutations (infectiousness, breaking immunity, etc)



#### From trees to networks

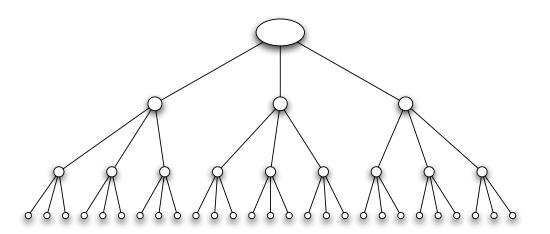
Recall that analysis of R<sub>0</sub> was for trees:



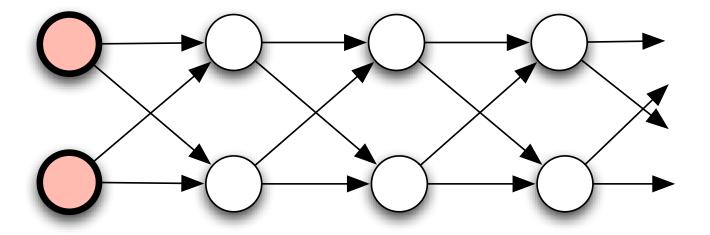
Do we have the same knife-edge R0 ~ I result in general graphs?

#### From trees to networks

Recall that analysis of R<sub>0</sub> was for trees:

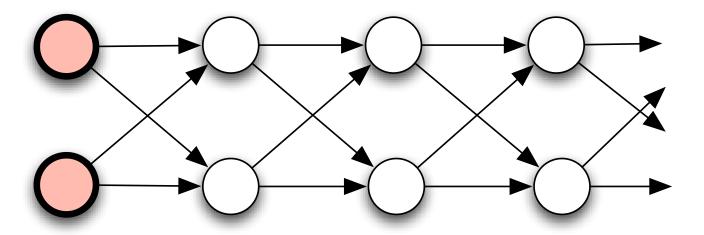


What happens on other networks? Consider p=2/3, k=2.



#### From trees to networks

What happens on other networks? Consider p=2/3, k=2.



Calculate R<sub>0</sub> as number of expected new cases per node

$$R0 = (2/3)*2 = 4/3 > 1$$

But this will almost certainly die out:  $(1/3)^4 = 1/81$  chance that all four edges fail even if both nodes are infected

Prob that this happens after finite number of steps converges to I

#### Now: SIS Epidemic Model

**S** = Susceptible

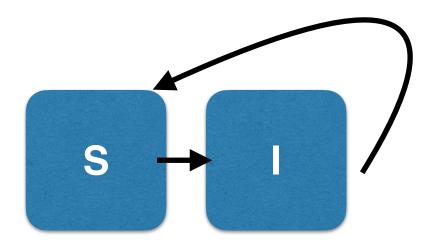
= Infectious: node is infected and infects with prob p

Initially some nodes in I state, rest in S state.

Each node in I state remains infected for t<sub>I</sub> time steps

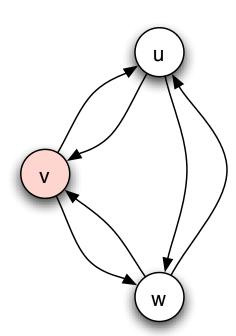
During each step, each node has probability **p** of infecting all neighbors

After t<sub>I</sub> time steps, node returns to S



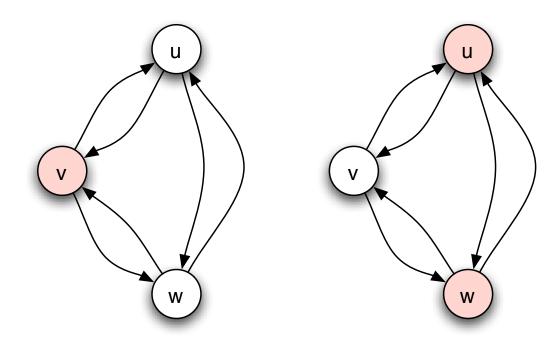
$$p=1/2$$

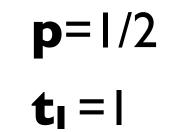
$$\mathbf{t_l} = \mathbf{I}$$

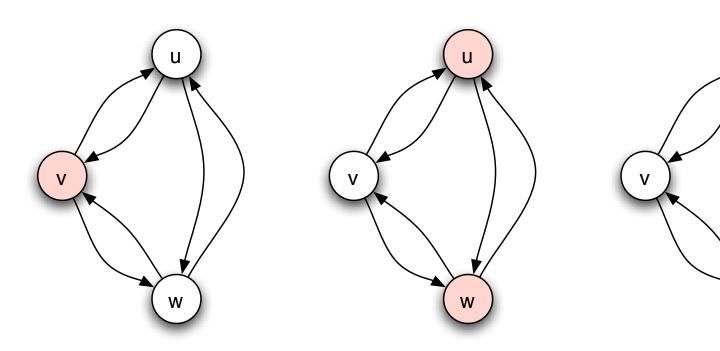


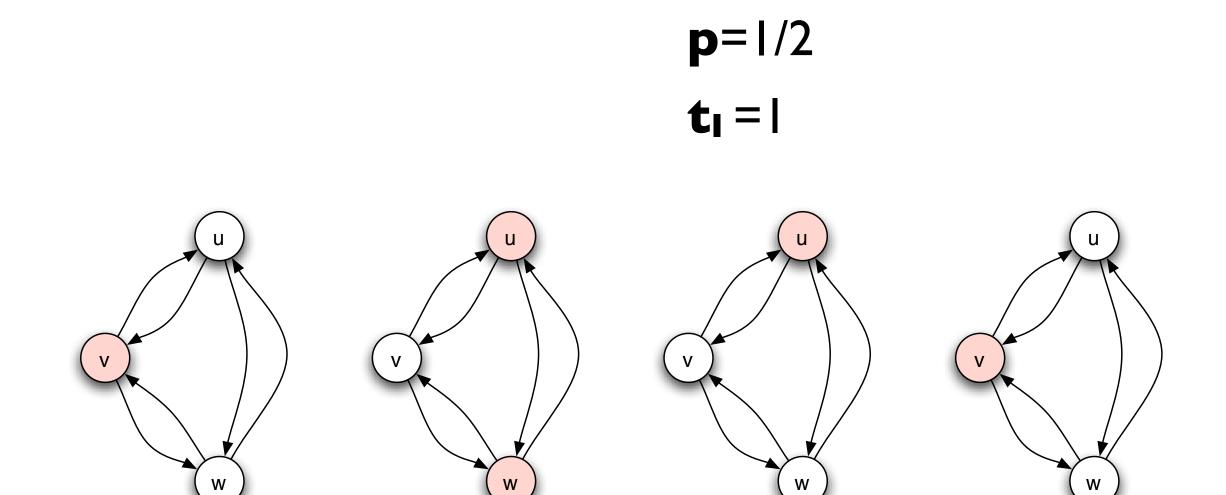
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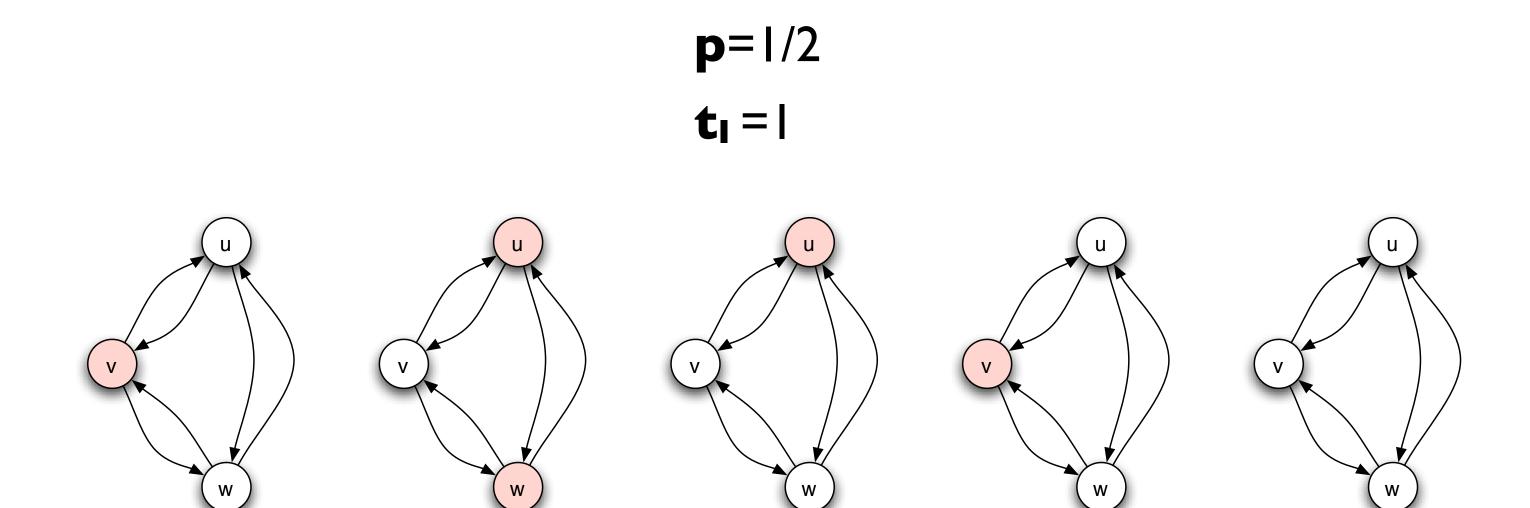
$$t_I = I$$









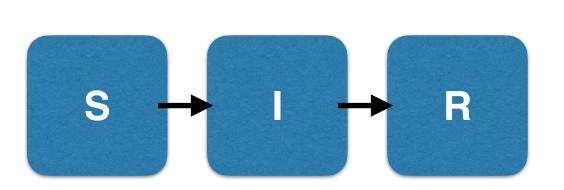


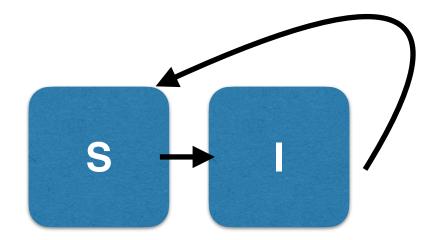
#### SIR vs. SIS

SIR: "burning through" a finite supply of susceptible

SIS: can run for a very long time, cycling through targets

SIS, like SIR, has a **critical threshold** ("knife-edge"/"tipping point"); trickier mathematical analysis. On non-trees both depend on more than just  $R_0$ .

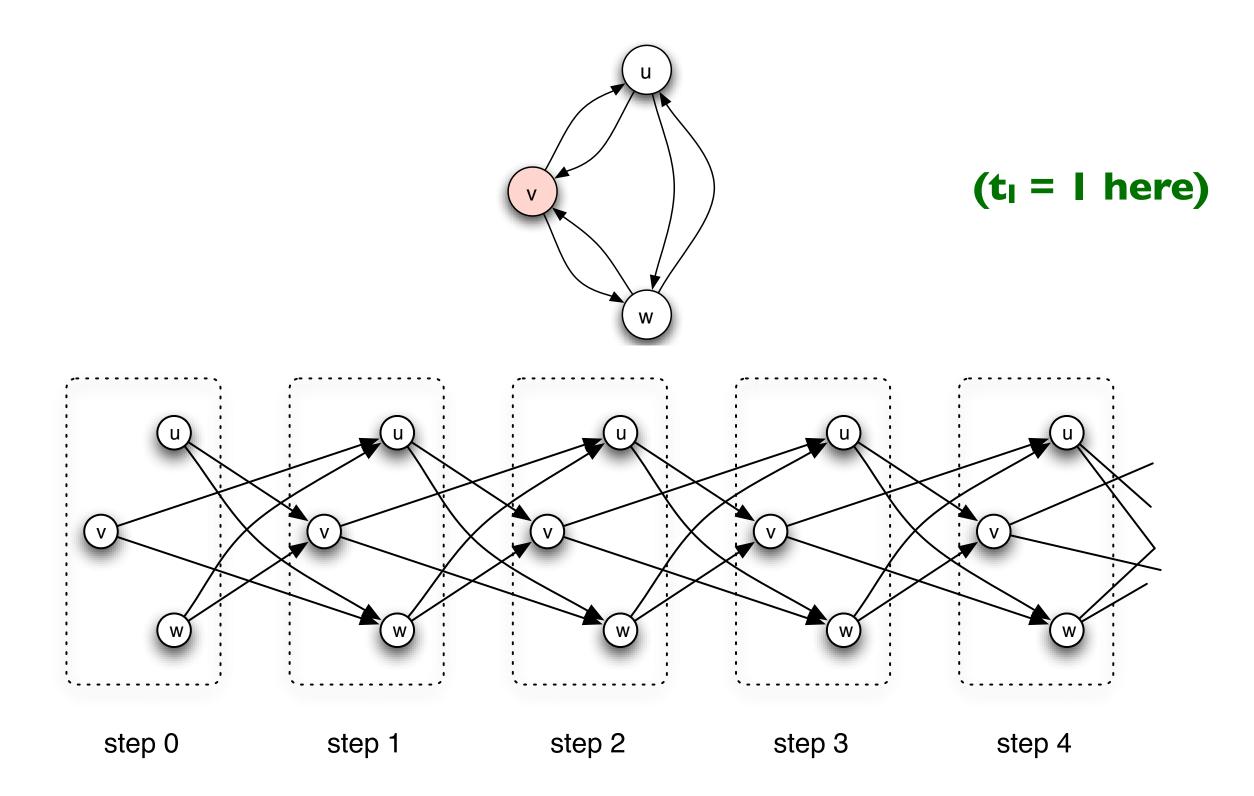




#### SIS as SIR on a bigger network

Consider time-expanded network: if u connects to v in network, have ut connect to vt+1

SIS is SIR on a time-expanded network.

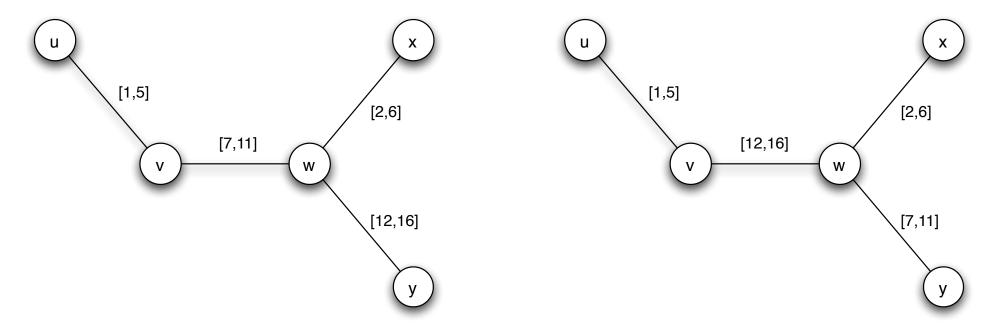


#### **Transient Contacts & Concurrency**

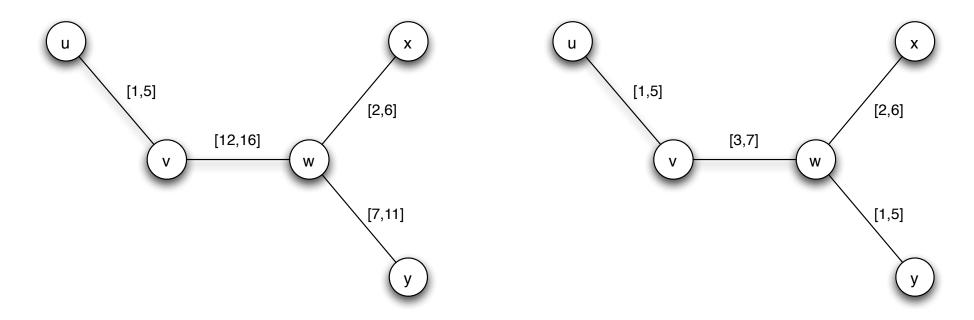
- So far, we've been analyzing **static** networks
- This is reasonable when the rate of transmission is typically much faster than edge creation/ deletion
- But some epidemic diseases last for years (HIV)
- When edges are active becomes very important

#### **Transient Contacts & Concurrency**

A less random model: it matters in what order contact is made in the contact network.

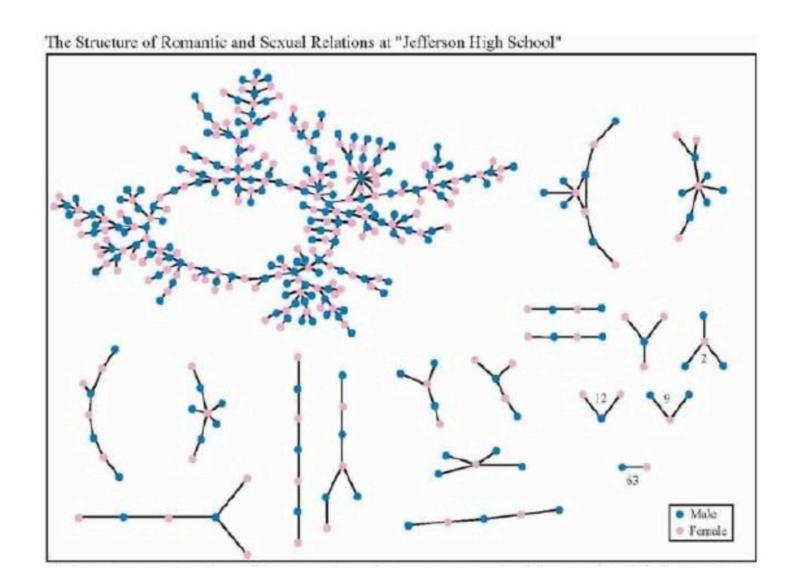


Concurrency: having two or more contacts at once.



#### **Transient Contacts & Concurrency**

Small changes in times can produce large differences in global epidemic spread
There are rich classes of network models incorporating transience and concurrency
It's not enough to just know the structure



#### Oscillations

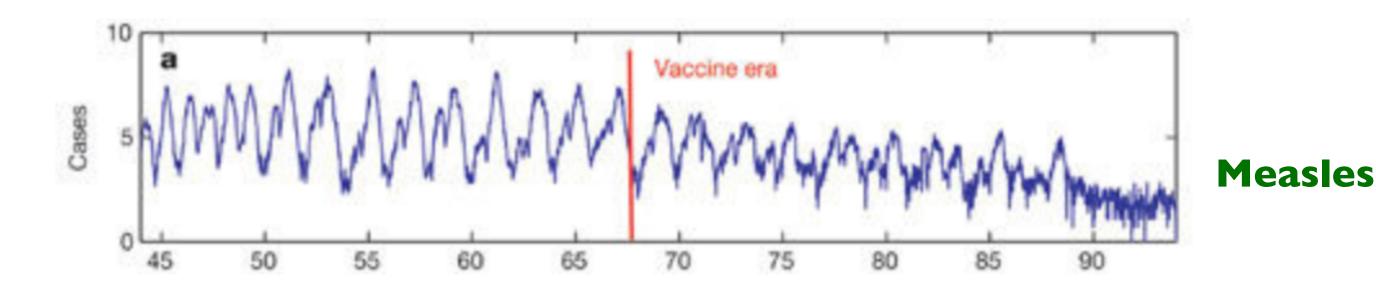
Diseases can be cyclical / have oscillations (like measles and syphilis)

To model this, vary the model so nodes have temporary immunity

SIRS: Susceptible, Infected for I steps, Recovered for R steps, then Susceptible again

This can produce oscillations in very localized parts of the network

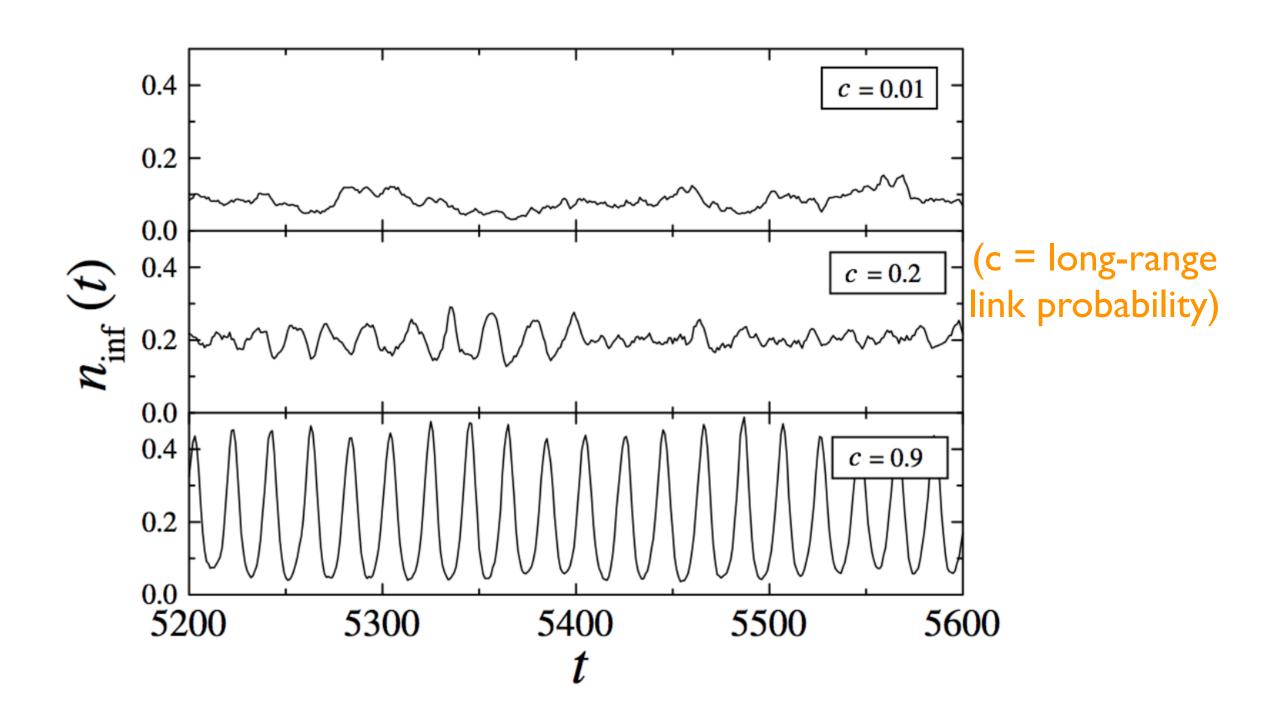
But for large fluctuations at the global network level, need **small-world structure** (random long-range contacts)



#### Oscillations

Diseases can be cyclical / have oscillations (like the flu)

But for large fluctuations at the global network level, need **small-world structure** (random long-range contacts)



In epidemic models, nodes get infected from one particular other node

To model information spread, people often use epidemic models ("viral diffusion")



But many social phenomena (behaviours, beliefs, practices, etc.) are complex: costly, risky, uncertain, etc.

When a behaviour is risky, costly, or uncertain, you may not do it just because one of your friends is (but this is what epidemic diffusion looks like)

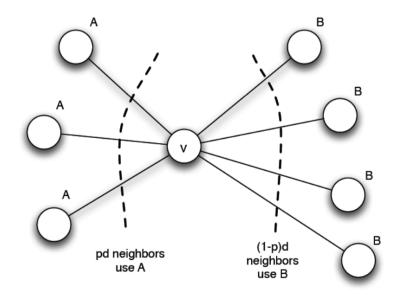
Social movements, health technologies, political activism, etc.

E.g.: PrEP medication is the best latest in HIV prevention  $\rightarrow$  one pill a day gives 90% prevention.

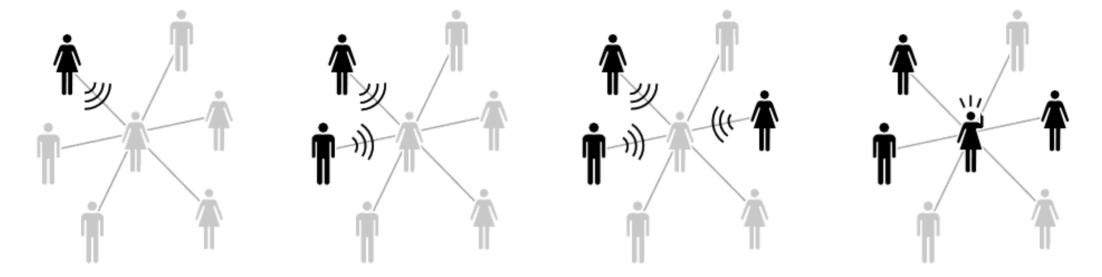
But in two trials in sub-Saharan Africa, it didn't work ... because no one was taking it! (fears of discrimination, etc.)

How do you get behaviour to diffuse?

Previously we saw a model of behaviour diffusion based on utility



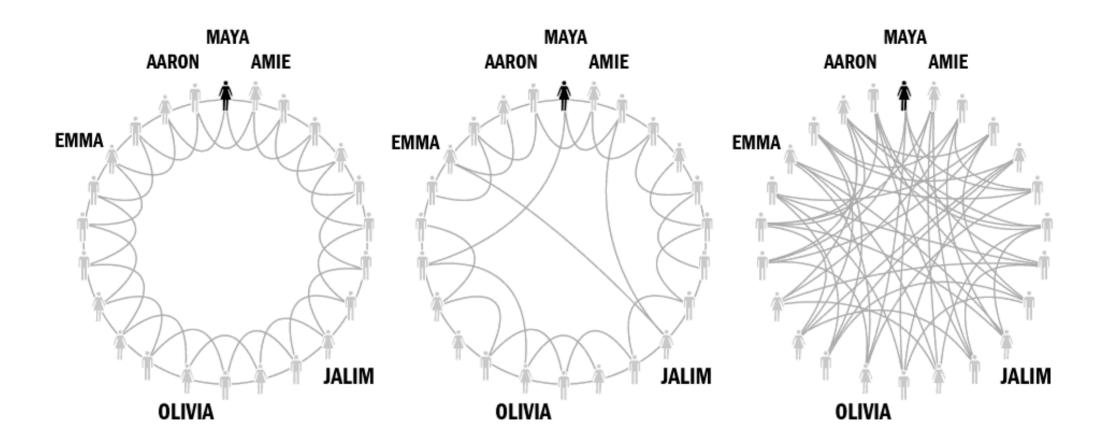
This is an example of *complex* diffusion: in general, need more than one neighbour to adopt before you adopt a behaviour.



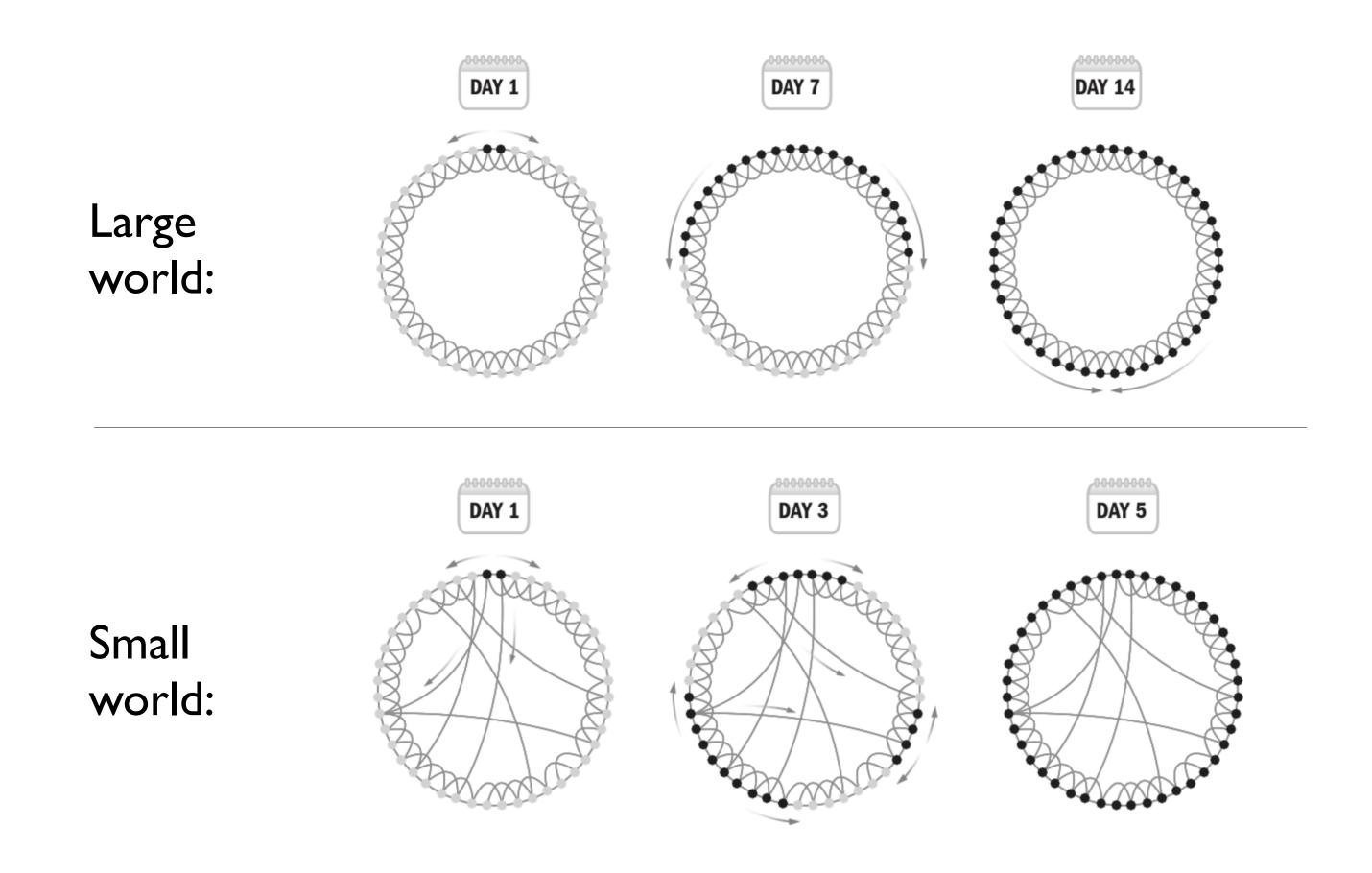
Simple vs. complex diffusion Epidemics vs. behaviour

What's the difference?

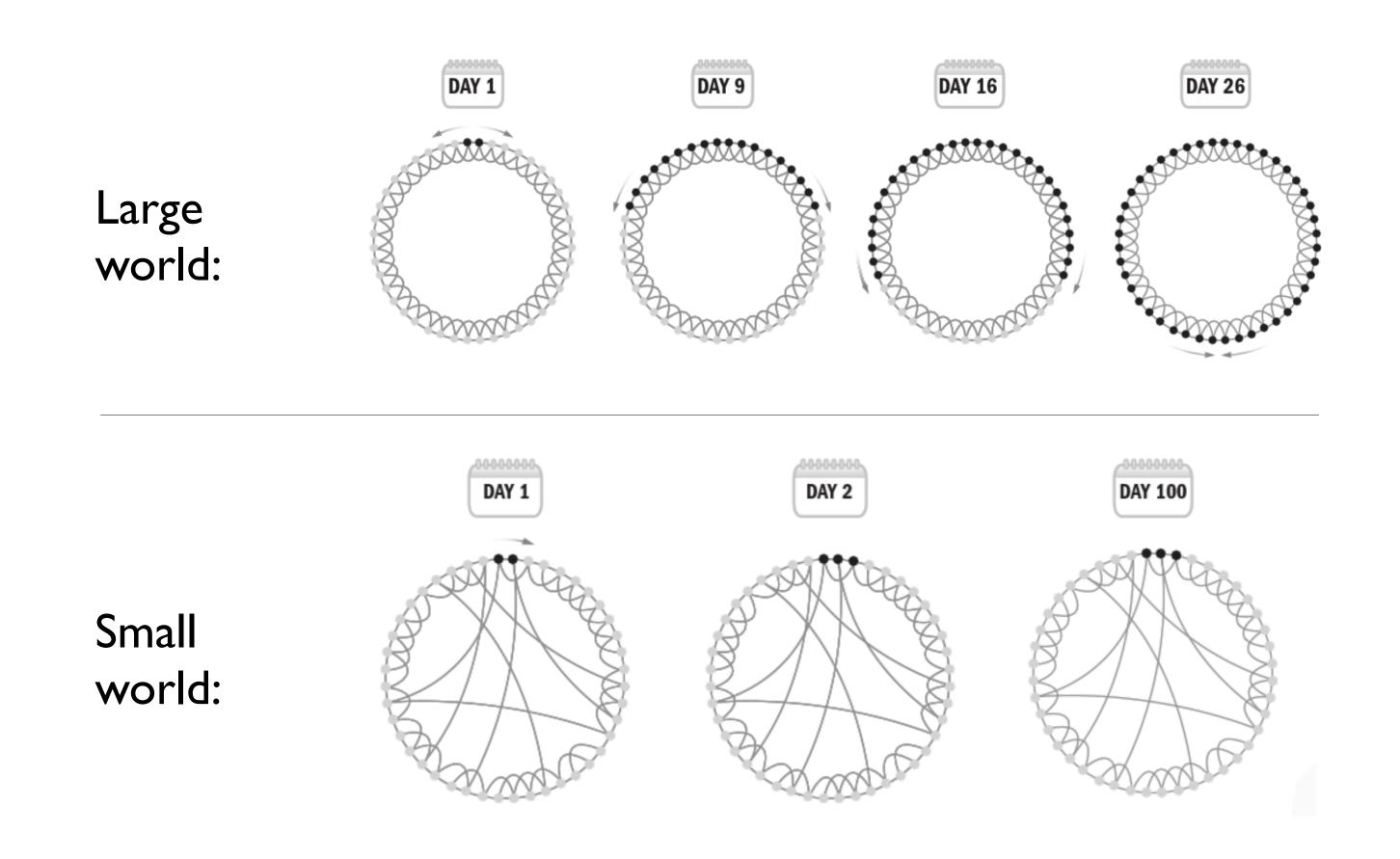
#### Recall the small-world model



## Simple Diffusion

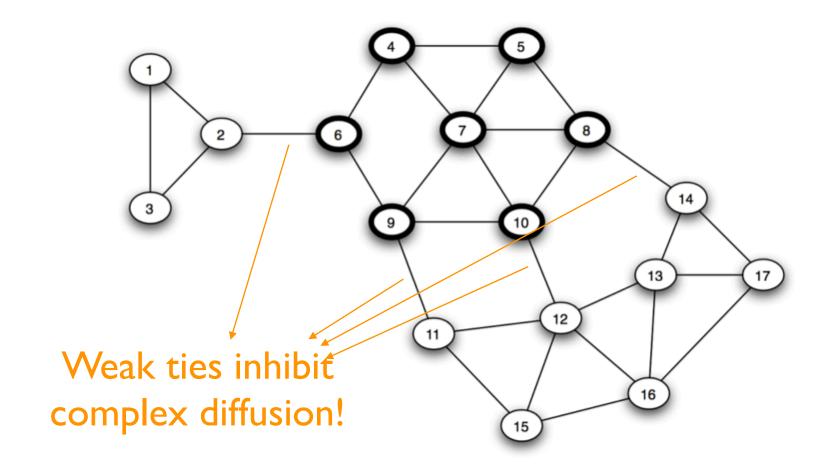


## Complex Diffusion



#### Simple vs. Complex Diffusion

Weak ties are extremely useful for simple diffusion and contagion, but they inhibit complex diffusion!



# Voting

## Voting

Why have voting?

#### Synthesize the preferences of a group

Aggregate information, preferences, beliefs, decisions

Voting on:

Candidates

Laws

Verdicts for trials

**Awards** 





## Simple example

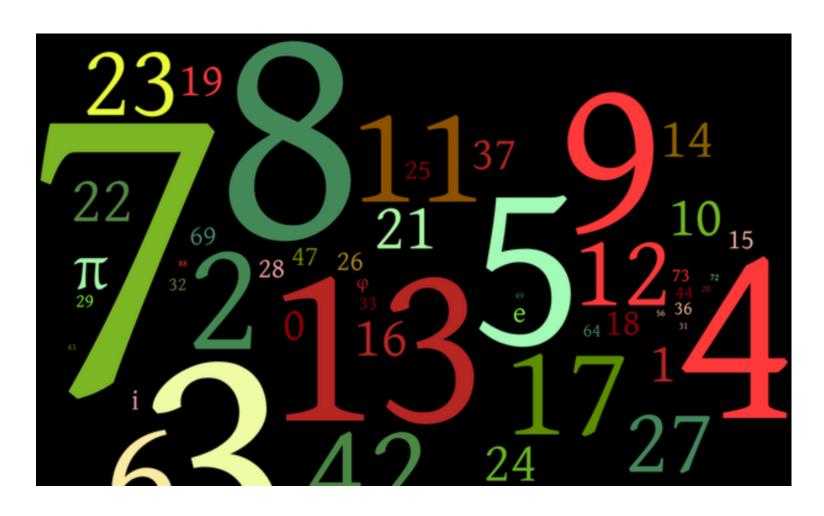
Say you want to pick the fairest outcome for the group

Everyone has their preferred number (e.g. price)

What should you do?

Easy...take the average

Why fair? Minimizes the squared loss



#### Why voting is hard

But in many situations there is no natural "average"!

Voting on:

Candidates

Laws

Verdicts for trials

Awards

#### Averaging fails here...



## Why voting is hard

Often need to pick a single winner that becomes binding for the group

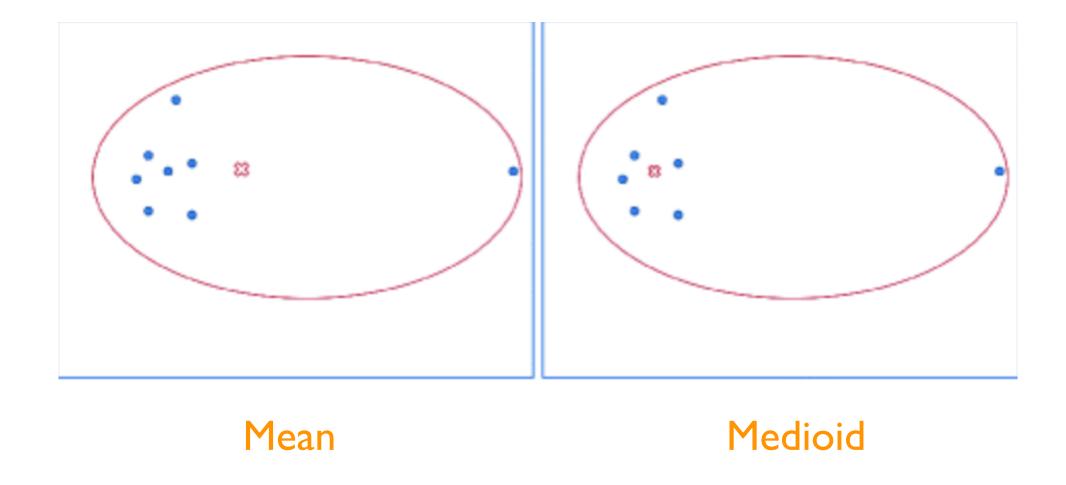
President

Award-winner

Policy decision

Voting as group decision making

Parallels to clustering: finding the centre vs finding the "medioid"—the best representative element



We want to aggregate many individuals' preferences

What are individual preferences?

Setup: a group of k people are evaluating a finite set of possible alternatives



The people want to produce a single **group ranking** that orders the alternatives from best to worst

The ranking should reflect the collective opinion of the group

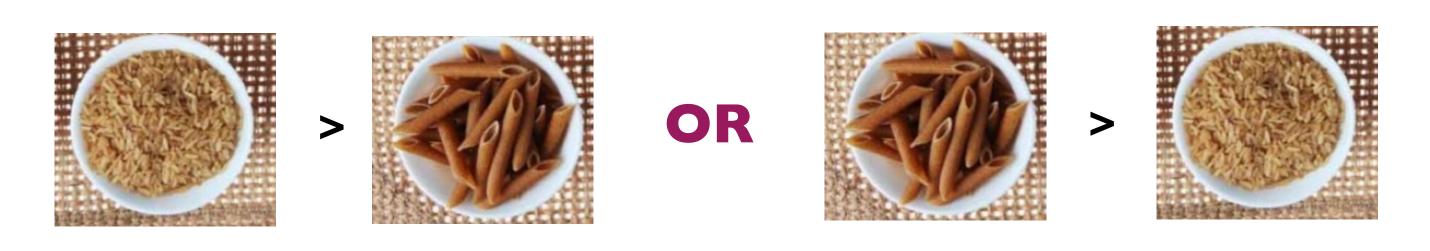
The challenge: how do we define what it means to reflect multiple, potentially contradictory opinions?



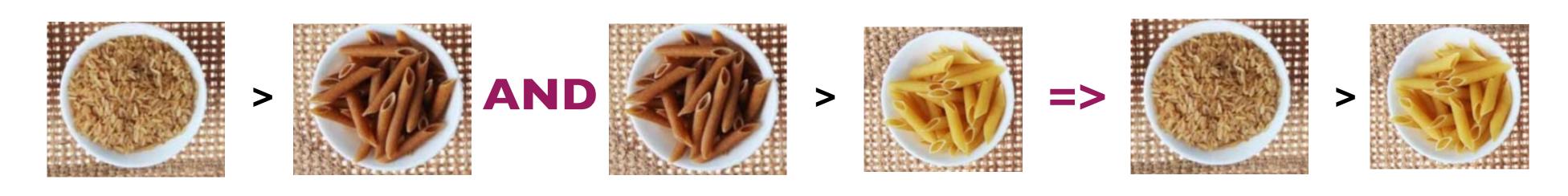
Every person has a preference relation over the alternatives, denoted  $>_i$  for player i

Must satisfy two properties:

**Complete**: all pairs of distinct alternatives X and Y, either  $X >_i Y$  or  $Y >_i X$ 



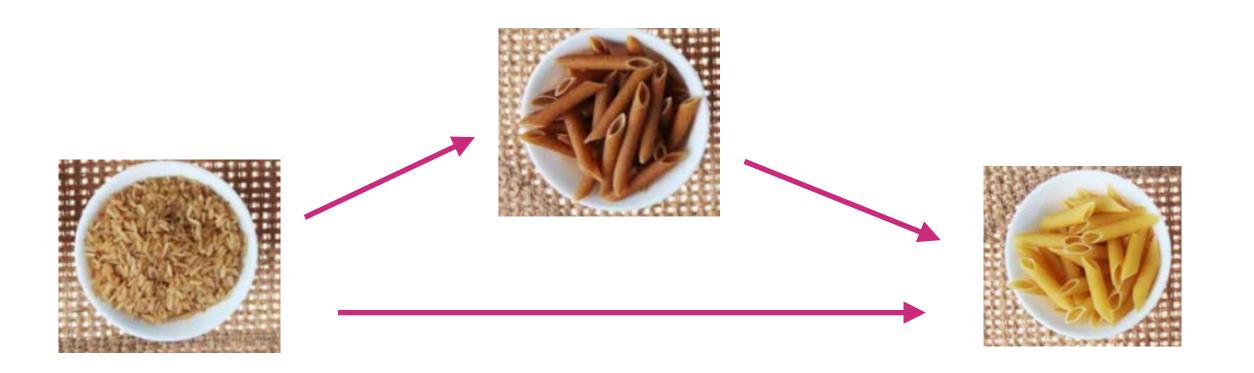
**Transitive**: if  $X >_i Y$  and  $Y >_i Z$  then  $X >_i Z$ 



A way to think about preference relations: as a graph

**Nodes:** alternatives

Directed edges:  $Y \longrightarrow X$  if  $X >_i Y$ 



(complete and transitive example)

Another way of expressing preferences: ranked list

#### For example:



Ranked list → preference relation

Obviously complete and transitive

Preference relation → ranked list

Less obvious but still true

Claim: Ranked list → Preference relation

#### Proof:

A ranked list is complete, since for any pair of alternatives X and Y, either X>Y or Y>X

A ranked list is transitive, since if X is higher than Y and Y is higher than Z, then X is also higher than Z.

Claim: Preference relation → ranked list

#### Proof:

Identify the alternative X that wins the most pairwise comparisons

Claim: X actually beats every other alternative

Why? Suppose  $Y >_i X$ . Then Y would beat everything X beats (by transitivity), and also X. Therefore beats more than X. Contradiction!

Put X at the top of the list, remove it from the set of alternatives, and recurse

Relation is still complete and transitive over remaining alternatives

Construct a list by repeatedly finding the alternative that beats everyone else

## Individual preferences

### Summary:

Preference relation → Ranked list

Ranked list → Preference relation

Therefore preference relations and ranked lists are equivalent!

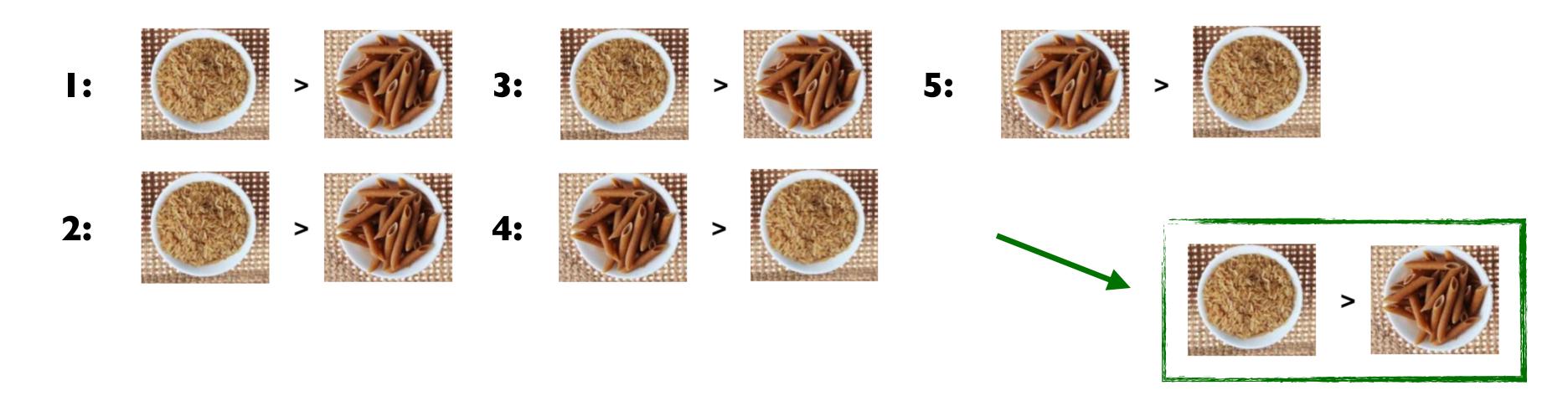
# Voting Systems

<u>Voting system</u>: a method that takes a set of complete and transitive individual preference relations (or ranked lists) and **outputs a group ranking** 

When there's only two alternatives, what should we do?

Majority Rule: whoever is preferred by a majority of the voters wins, other one is second

#### (let k be odd to avoid ties)



Easy enough, what about majority rule with more than two alternatives?

### What's a natural way to extend it?

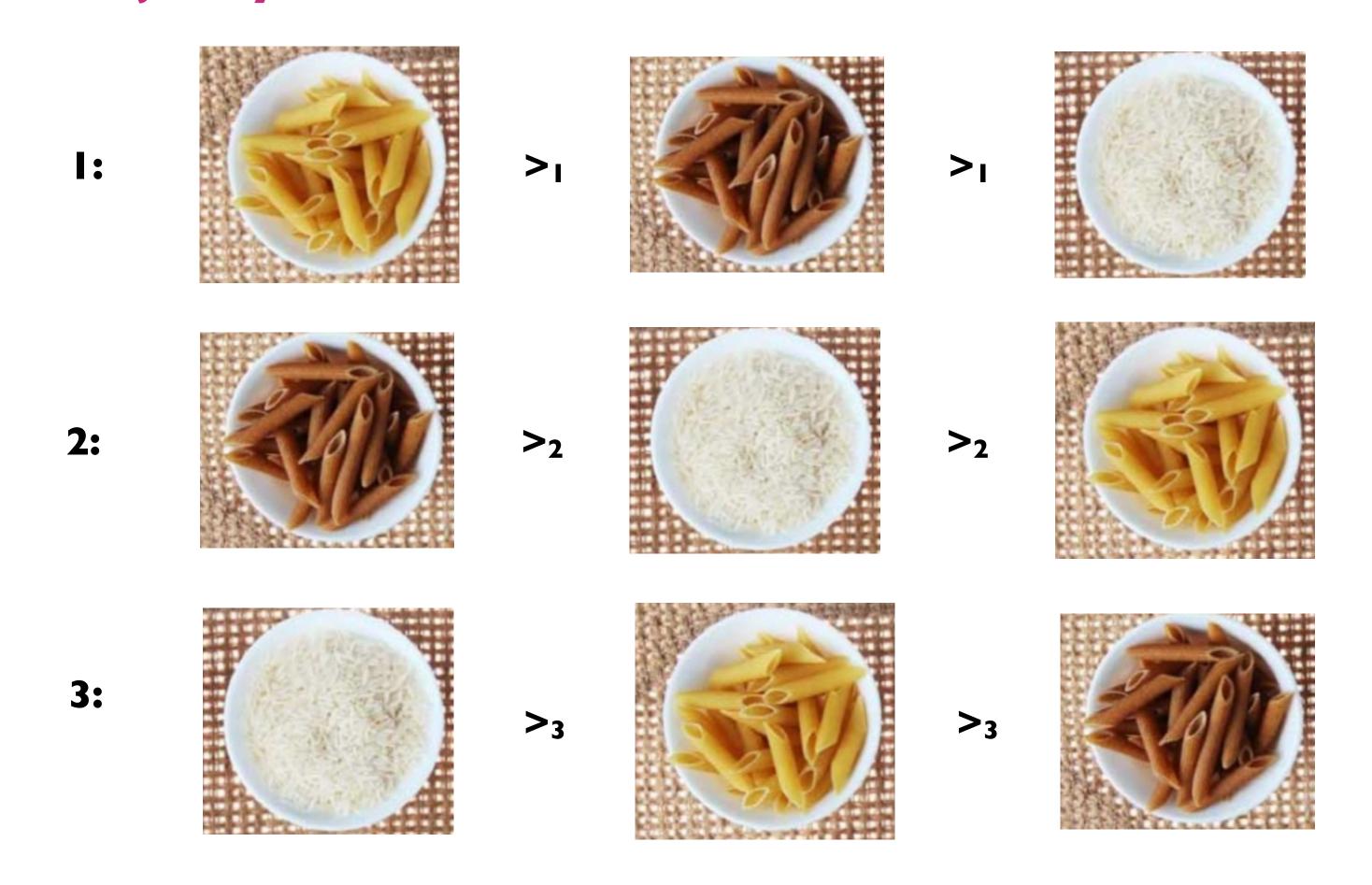
Majority rule on every pair of alternatives: X > Y if a majority of voters have  $X >_i Y$ 

### Is this complete?

Everyone has a preference for every pair, and there's always a majority (assume k is odd). So this is **complete** 

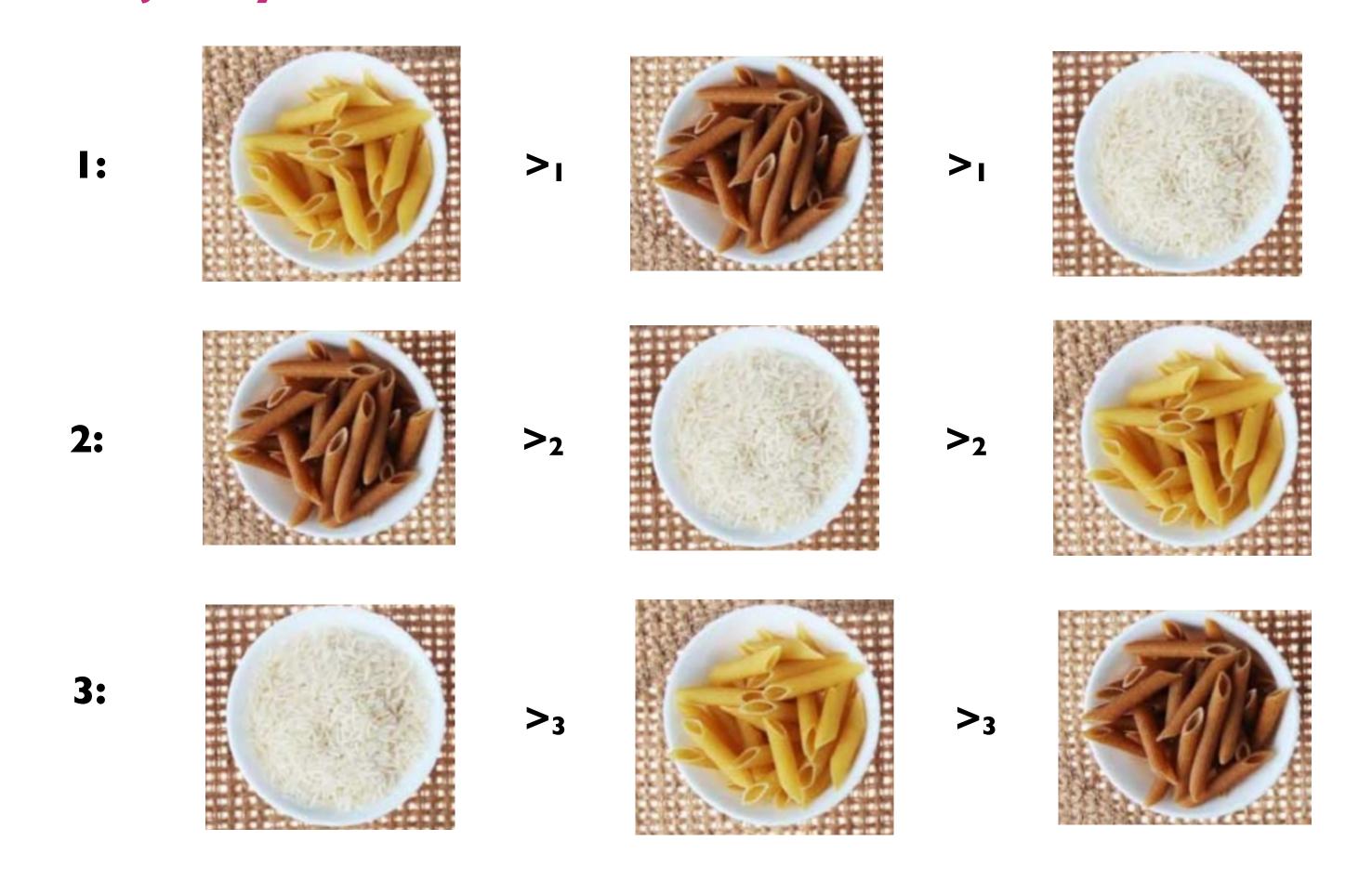
#### Is this transitive?

Is majority rule on at least 3 alternatives transitive?



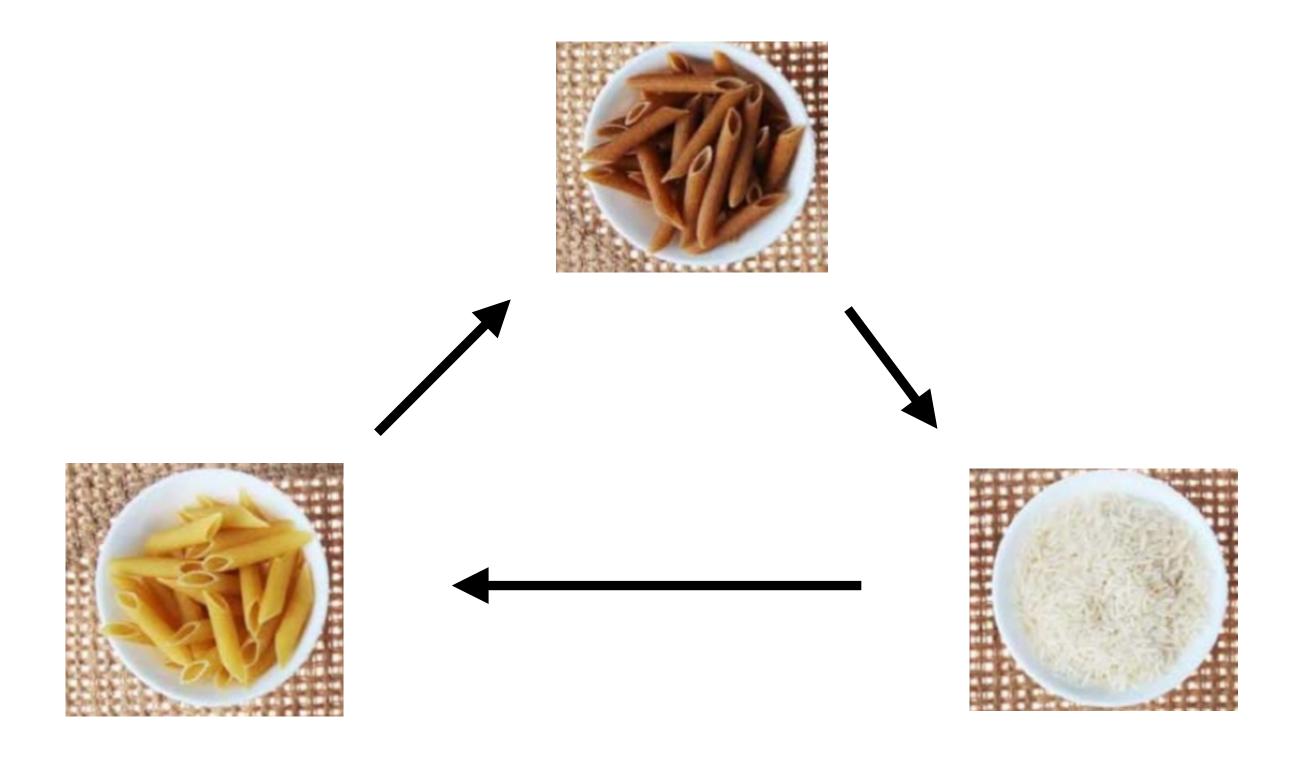
What does majority rule do here?

Is majority rule on at least 3 alternatives transitive?



Y pasta > B pasta, B pasta > rice, rice > Y pasta!

Majority rule with at least three alternatives can produce a non-transitive group ranking



Cycle on preferences => non transitive => bad!

## **Condorcet Paradox**

Majority rule with at least three alternatives can produce a *non-transitive* group ranking

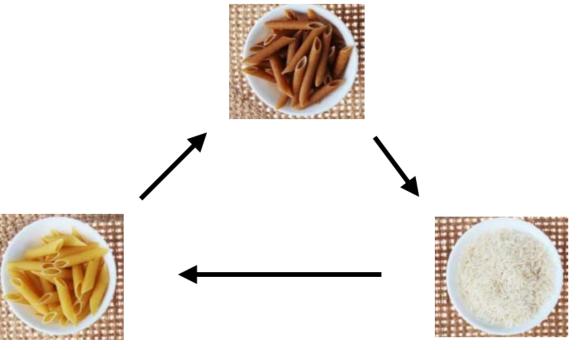
Called the "Condorcet Paradox"

### Really bad news!

Everyone had perfectly plausible preferences

But they behave incoherently as a group, can't even decide on a favourite





### **Condorcet Paradox**

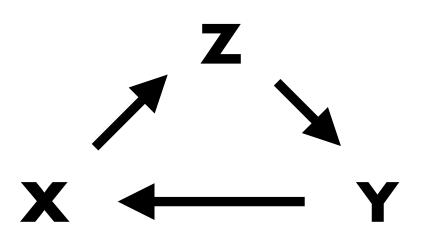
# Condorcet Paradox can even happen within a single individual person

Consider a student deciding which college to attend

Wants a highly-ranked college, a small average class size, and maximum scholarship money

Plans to decide between pairs by favouring the one does better on the most criteria

College	National Ranking	Average Class Size	Scholarship Money Offered
X	4	40	\$3000
Y	8	18	\$1000
Z	12	24	\$8000





### What about using majority rule another way?

Iterative approach: find a winner, remove from the list, and recurse

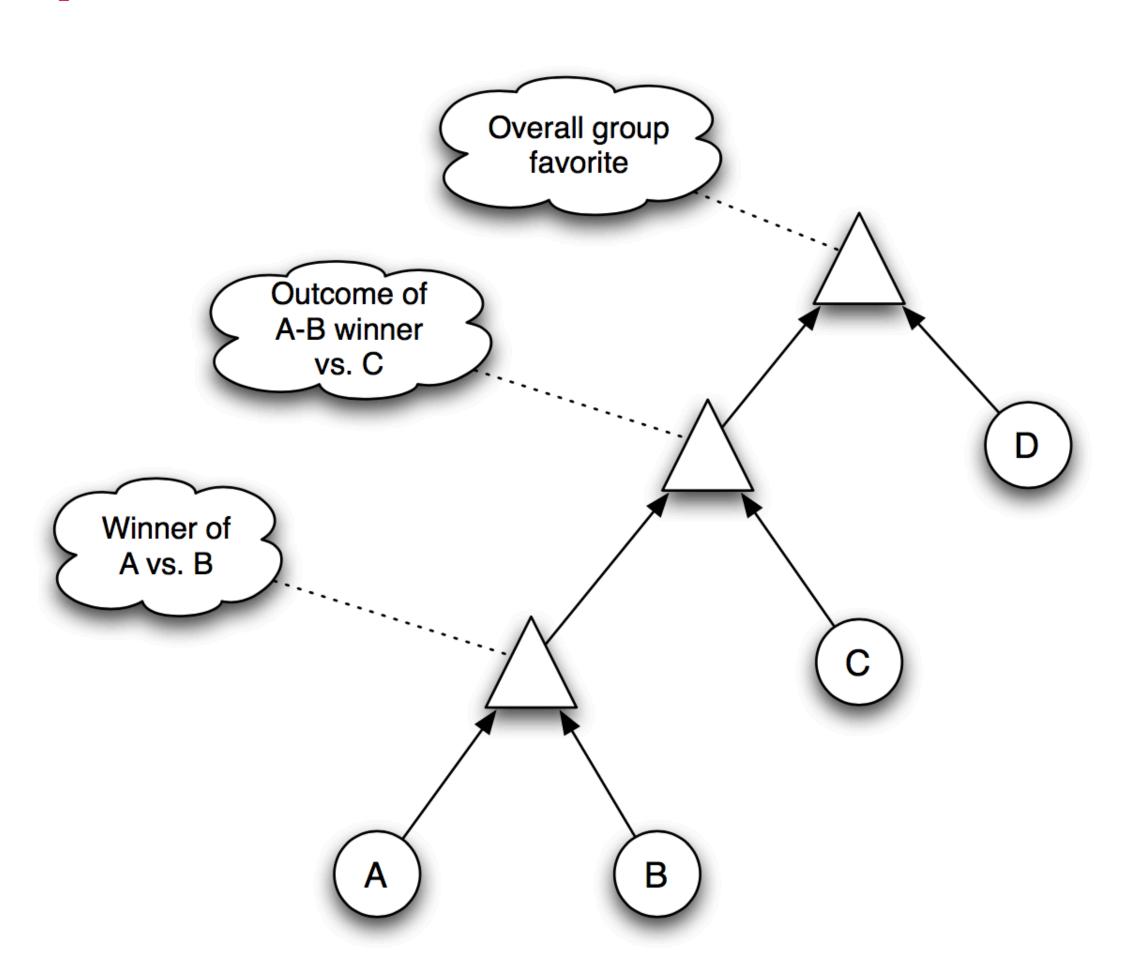
One idea: arrange alternatives in some order, then compare by majority vote, compare the winner to the third alternative, and so on.

Winner of the final comparison is the group favourite

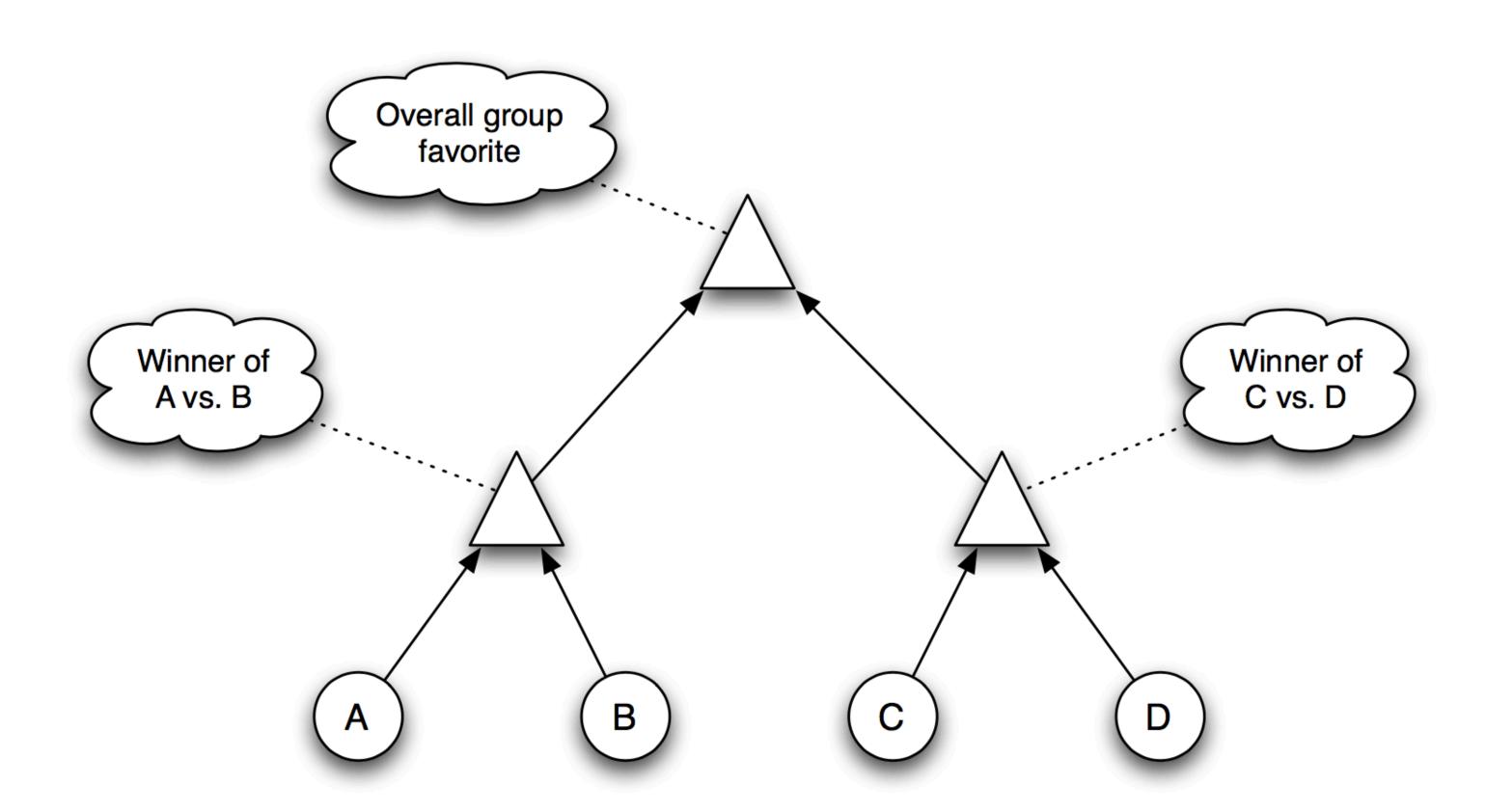
More generally, we can schedule any kind of elimination tournament to determine the favourite

→ Then recurse!

### **Graphically:**



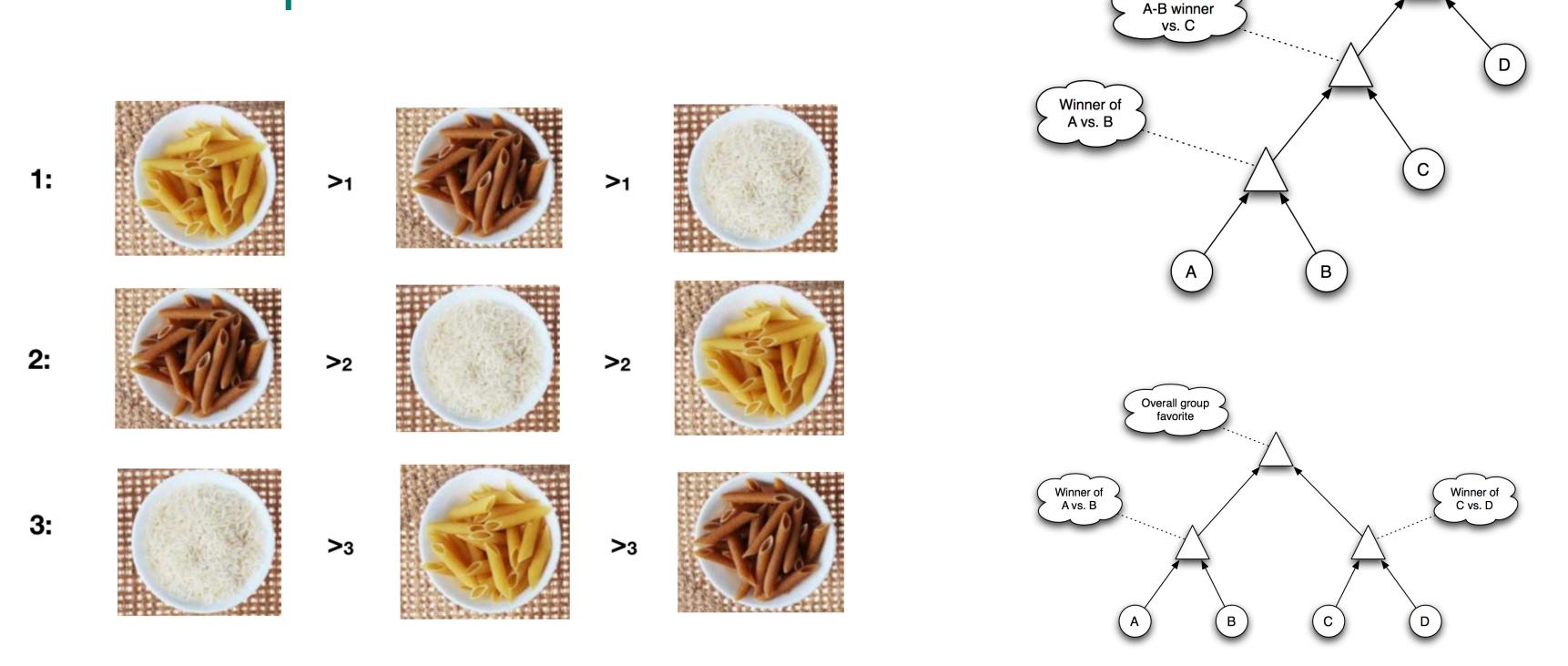
#### Other kind of elimination tournament:



What's wrong with this?

Strategic agenda setting: order matters!

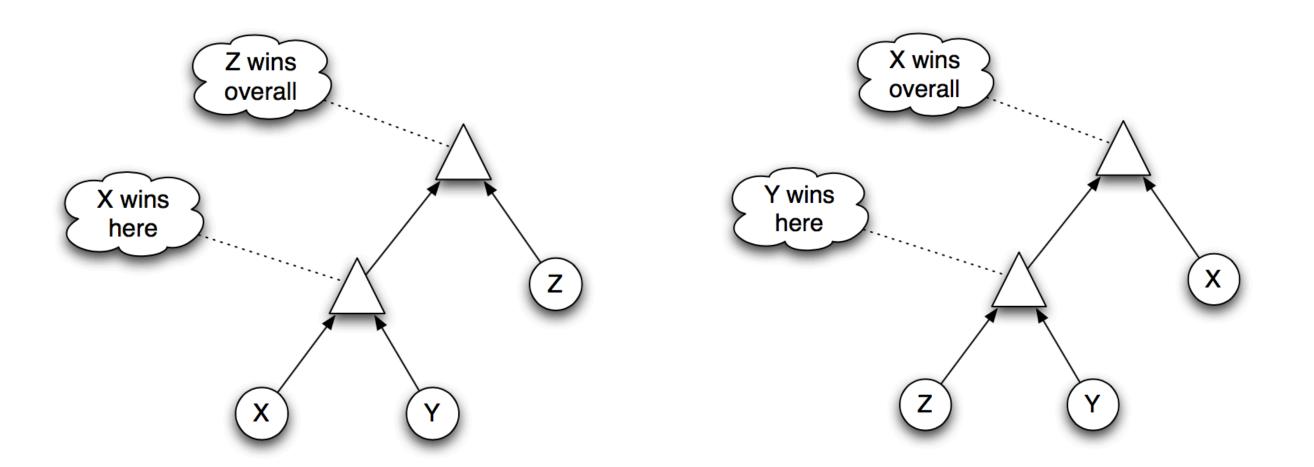
Consider example from before:



Overall group

In what order do we evaluate the alternatives?

In what order do we evaluate the alternatives?



Entire ranking is entirely determined by the order in which we evaluate!





















## Other systems?

Majority rule led to some bad outcomes

What about other strategies?

**Positional voting:** produce a group ranking directly from the individual rankings

Forget pairwise comparisons

Each alternative receives a certain **weight** based on its positions in all the individual rankings

Heisman trophy in college football (and NBA MVP, etc.) all use the following method: get weight 0 for being picked last, I for being second last, ..., k-I for being picked first

Repeat for each voter, tally up the scores, and rank

### Example: two voters, four alternatives

Voter I:  $A >_I B >_I C >_I D$ 

Voter 2:  $B >_2 C >_2 A >_2 D$ 

A: 3 + 1 = 4

B: 2 + 3 = 5

C: 1 + 2 = 3

D: 0 + 0 = 0

Group ranking: B > A > C > D

Called the "Borda Count"





You can create your own variants (and many have) by changing the number of points per position

Example: if only top 3 matter, you could assign 3 for first place, 2 for second place, I for third place, and 0 otherwise

Any such system is a "positional voting system"

Ignoring ties, Borda Count always produces a complete, transitive ranking!



But the Borda Count has its own problems

Magazine tries to rank greatest movie of all time, asks five film critics to rank Citizen Kane and The Godfather

Three prefer CK, two prefer TG => CK>TG => all good!

At the last second, they want to inject some modernity into the discussion, so they include Frozen

First three only like old movies, so they vote:

$$CK >_i TG >_i F$$

Critics 4 and 5 only like past 40 years, so:

$$TG >_i F >_i CK$$

What is the Borda Count now?



First three only like old movies, so they vote:

$$CK >_i TG >_i F$$

Critics 4 and 5 only like past 40 years, so:

$$TG >_i F >_i CK$$

Borda:

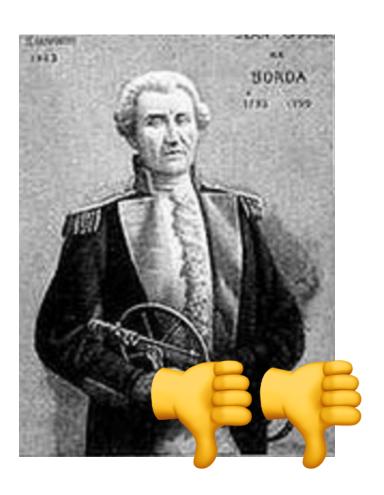
$$CK: 6, TG: 7, F: 2 => TG > CK > F$$

But before Frozen was introduced it was CK > TG!

TG and CK flip because of Frozen??

#### Both TG and CK beat Frozen head-to-head

Yet still Frozen influenced CK > TG



Borda Count is susceptible to "irrelevant alternatives"

What voters think of Frozen should be irrelevant to how they feel about relative ranking of TG and CK

#### But it isn't

This gives rise to another problem: voters can strategically misreport their preferences

For example, say voters 4 and 5 actually had the true ranking

$$1,2,3: CK >_i TG >_i F$$

$$4,5:TG >_i CK >_i F$$

Borda: 
$$CK >_i TG >_i F$$

By lying and reporting TG  $>_i$  F  $>_i$  CK, they get TG to win



## Irrelevant Alternatives in Politics

These problems with "irrelevant alternatives" and strategic misreporting have happened in elections around the world

Most vote with **plurality voting:** the candidate ranked at the top by most voters wins

Q: is this a positional voting system?

A: Yes: I for winner, 0 otherwise

"Third-party effects"/"spoiler effects": if very few people favour some candidate, this can swing outcome of two leading contenders

In response, some people strategically misreport their preferences

### What's The Deal?

Voting is one society's most important institutions

On its face, seems like a relatively simple problem

But we can't find a system that doesn't have horrible pathologies!

Is there any system that is free of pathologies?

### What's The Deal?

### Is there any system that is free of pathologies?

Let's define "Free of pathologies"

- Criterion I "Unanimity": if there is a pair X and Y for which  $X >_i Y$  for every i, then  $X >_i Y$
- Criterion 2 "Independence of Irrelevant Alternatives" (IIA): the ordering of X and Y should only depend on the relative positions X and Y in individual rankings

If we have a bunch of rankings that produces a group ranking with X > YThen we move some Z around in the individual rankings

#### It should still be the case that X > Y

• Criterion 3 "Non-Dictatorship": the group ranking should not just always be what one particular voter thinks

## Independence of Irrelevant Alternatives







# Good Voting Systems

### What satisfies Unanimity and IIA and non-dictatorship?

With two alternatives, majority rule clearly satisfies all

Arrow's Theorem [Arrow 1953]: With at least three alternatives, no voting system satisfies Unanimity, IIA, and Non-dictatorship

# In general, there is no good voting system!

In practice, this means that there will always be inherent tradeoffs we have to choose from



### What Do We Do Now?

# How do we vote, how do we decide on things in the presence of Condorcet's Paradox and Arrow's Theorem?

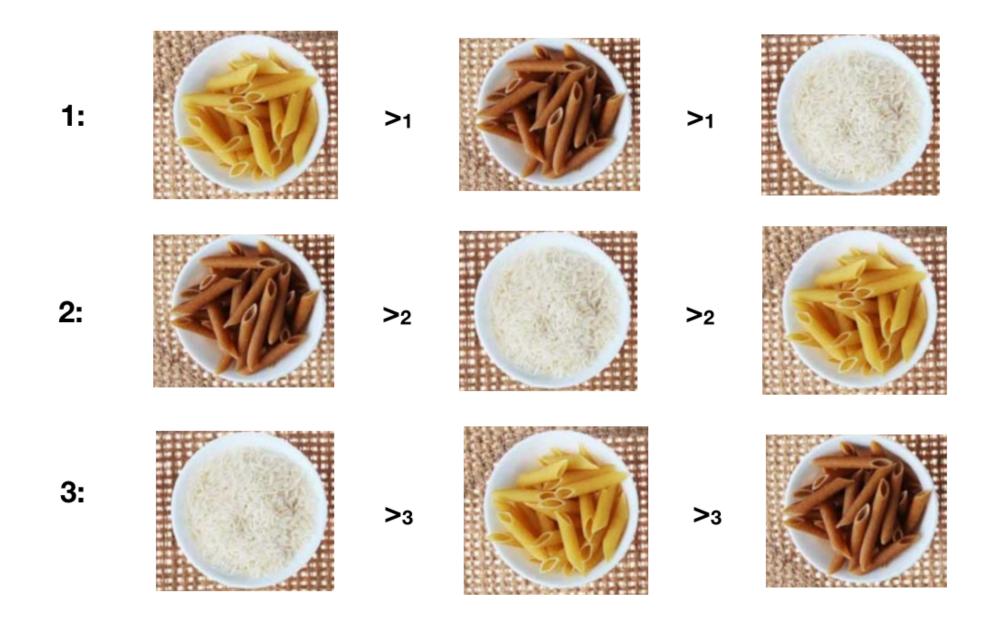
If you're faced with an impossibility result, you don't just give up

One common technique is to look for important special cases

Arrow's Theorem is a **general result**, so it doesn't necessarily apply if we make some additional assumptions

### What Do We Do Now?

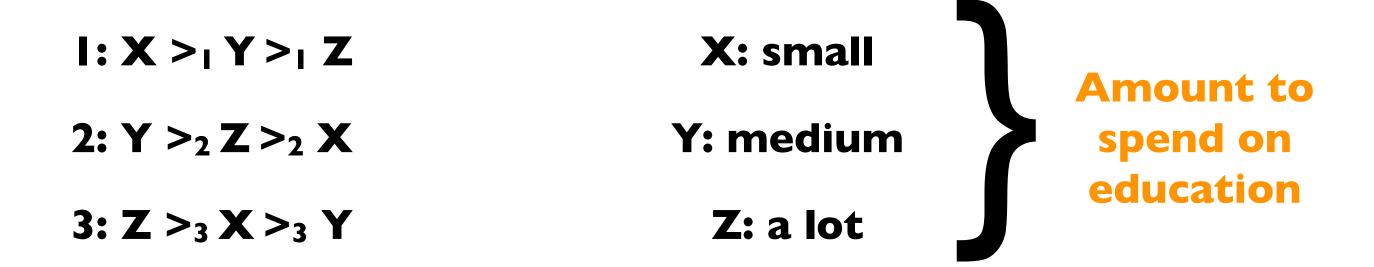
Go back to original Condorcet problem



Replace food with choices about how much money to spend on education

### What Do We Do Now?

Go back to original Condorcet problem with money now:



Voter I's preferences "make sense"

Voter 2's preferences do too: prefer between Y and Z, so say Y then Z then X

Voter 3's preferences are harder to justify

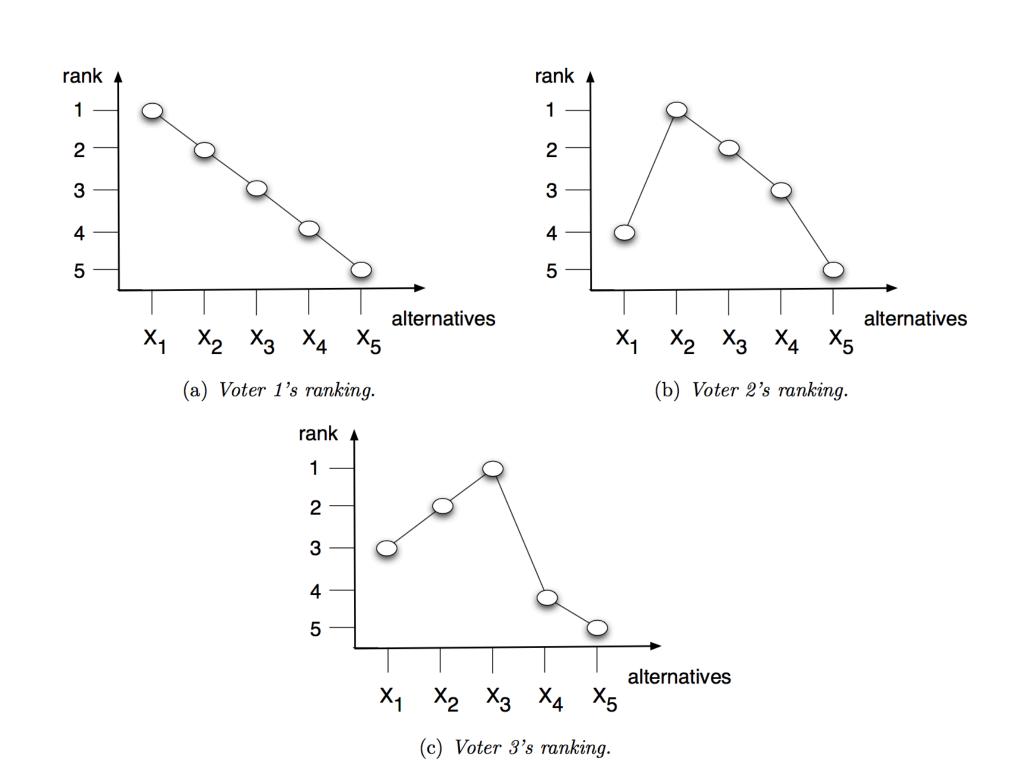
Not impossible, but they're more unusual

## Ideal Points

Assume the preferences lie on a one-dimensional spectrum, and each voter has an "ideal point" on the spectrum

They evaluate alternatives by proximity to this ideal point

Actually we can assume something weaker: each voter's preferences "fall away" consistently on both sides of their favourite alternative

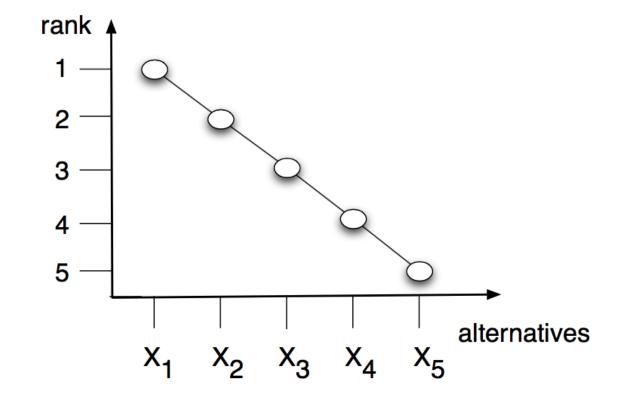


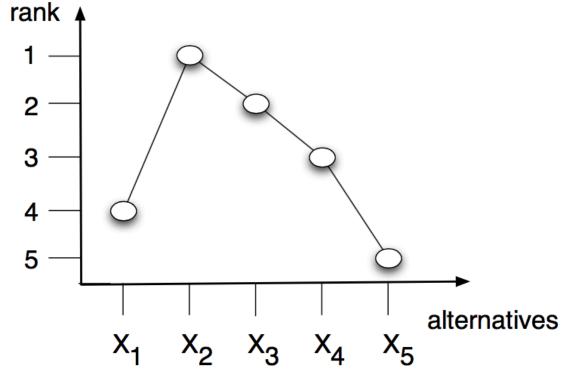
## Single-Peaked Preferences

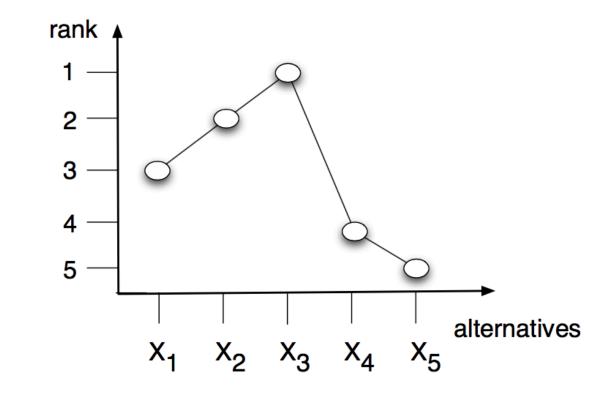
**Definition:** a voter has "single-peaked preferences" if there is no alternative  $X_s$  for which both neighbouring alternatives  $X_{s-1}$  and  $X_{s+1}$  are ranked above  $X_s$ 

**Equivalent to**: every voter i has a top-ranked option  $X_t$ , and her preferences fall off on both sides of t:

$$X_t \succ_i X_{t+1} \succ_i X_{t+2} \succ_i \cdots$$
 and  $X_t \succ_i X_{t-1} \succ_i X_{t-2} \succ_i \cdots$ 







## Single-Peaked Preferences

Majority rule with single-peaked preferences

Recall majority rule: compare every pair of alternatives X and Y, and decide X > Y or Y > X by the majority of voters

**Claim:** If all individual rankings are single-peaked, then majority rule applied to all pairs of alternatives produces a group preference relation that is **complete** and **transitive**.

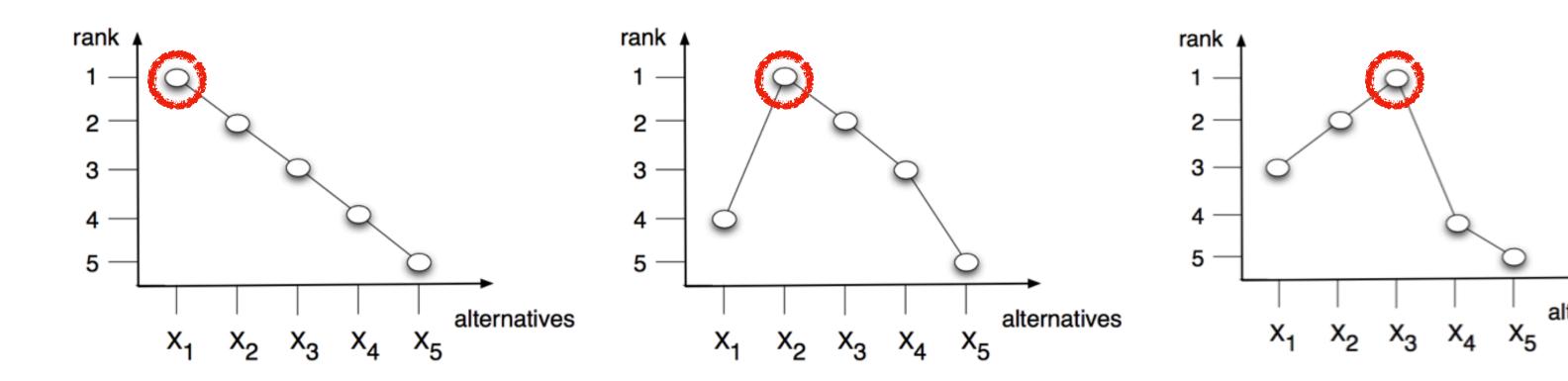
In other words, majority rule works!

Start off by trying to find a group favourite, then proceed by recursion on the rest of the alternatives

Consider every voter's top-ranked alternative — their peak — and sort this set of favourites from left to right along the spectrum

A popular alternative can show up many times

Now consider the **median** of these favourites



Favourites: X<sub>1</sub>, X<sub>2</sub>, X<sub>3</sub> Median: X<sub>2</sub>

The median individual favourite is a natural candidate for potential group favourite

Strikes a compromise between more extreme favourites on either side

**Median Voter Theorem:** With single-peaked preferences, the median individual favourite defeats every other alternative in a pairwise majority vote.

Let  $X_m$  be the median individual favourite, and  $X_t$  some random

Assume t > m, and sort the voters in order of their favourites

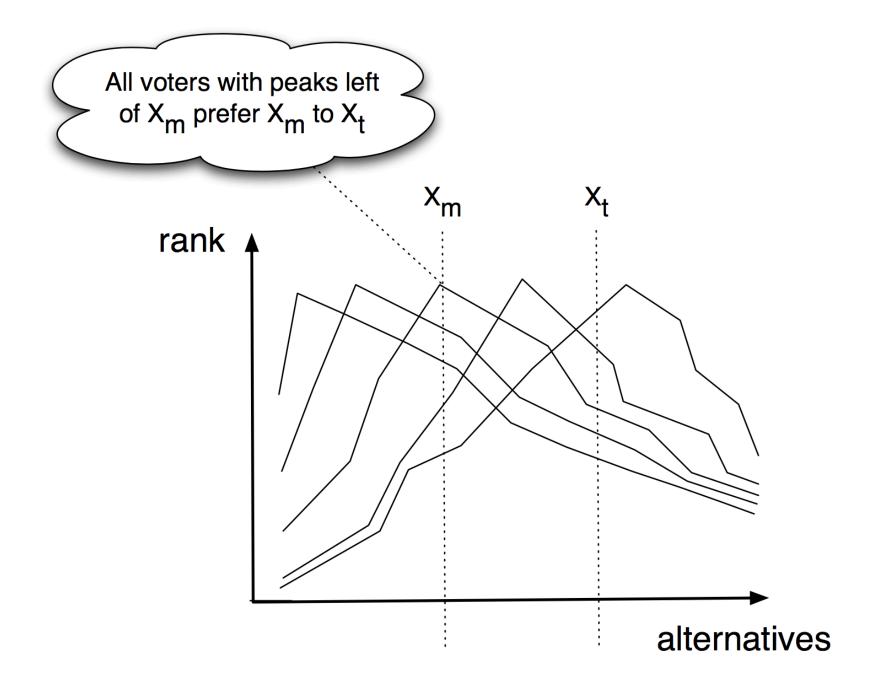
[remember k is odd]

By definition  $X_m$  is in position (k+1)/2 of sorted favourites

In the first (k+1)/2 positions, everyone either prefers  $X_m$  or their favourite is to the left of  $X_m$ 

For everyone in the latter group, both  $X_m$  and  $X_t$  on the right-hand "down-slope" of preferences, but  $X_m$  is higher than  $X_t$ 

So (k+1)/2 people prefer  $X_m$  to  $X_t$ —which is a strict majority!



So  $X_m$  can gather a majority of support against any other alternative  $X_t$  (case to the left of  $X_m$  is completely symmetric)

Build up group favourites one at a time: start by finding global favourite, then remove from individual rankings and recurse

Notice that removing it from the rankings still maintains the single-peaked property of everyone's preferences!

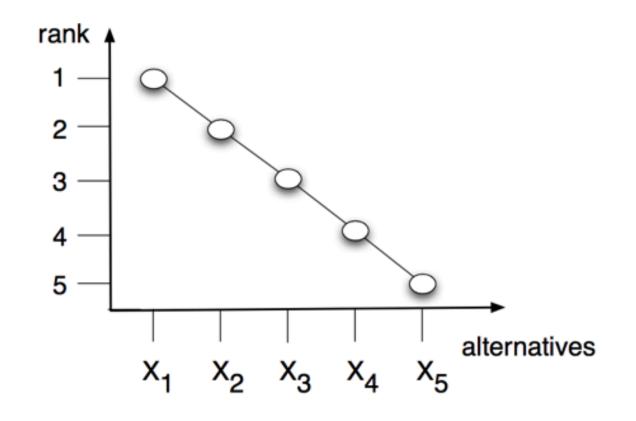
It just chops the peak off

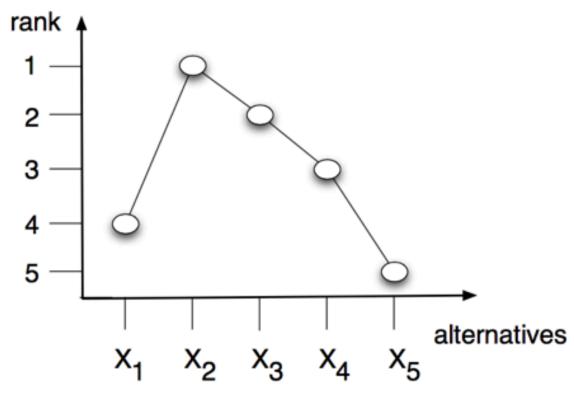
# Example

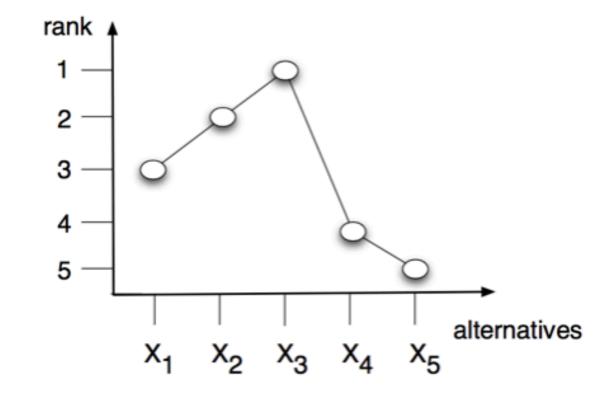
X<sub>2</sub> is global median favourite

Then favourites are  $X_1, X_3, X_3 => X_3$  median favourite

Eventually we get  $X_2 > X_3 > X_1 > X_4 > X_5$ 







## Voting as Information Aggregation

So far, trying to come up with methods for people who have different preferences

Sometimes there is a "true" underlying ranking and people with different information are trying to uncover it

Examples:

Jury deliberation

Board of advisors to a company

#### Simple Case: Simultaneous, Sincere Voting

Simple setting, two alternatives X and Y

One is genuinely the best choice, each voter casts vote on what she thinks the right choice is

#### Assume everyone votes sincerely

#### Model: similar to information cascades

Prior probability that X is best is 1/2

Each voter gets a private independent signal on which is best, prob of getting right signal is q (> 1/2)

With probability q, voter should vote for what her signal says

Condorcet Jury Theorem: as the number of voters increases, probability of the majority choosing correct decision goes to I

Oldest "wisdom of crowds" argument

#### Simple Case: Simultaneous, Sincere Voting

```
Formal Bayes argument
```

```
Recall Bayes Rule: P[A|B] = P[B|A]P[B]/P[A]
```

We want to compute P[X is best | X-signal]

```
Given: P[X \text{ is best}] = 1/2 \text{ and } P[X \text{-signal} \mid X \text{ is best}] = q
```

Voter's strategy: evaluate P[X is best | X -signal] then vote X if this probability > 1/2

P[X is best | X -signal] = P[X -signal | X is best]P[X is best]/P[X -signal]

X-signal can be observed if X is best or if Y is best:

P[X-signal] = P[X is best] \* P[x-signal observed | X is best] + P[Y is best] \* P[X-signal observed | Y is best] = 1/2q + 1/2(1-q) = 1/2

So overall: P[X is best | X -signal] = (1/2)q / (1/2) = q

Voter favours the alternative that is reinforced by her signal

# Insincere Voting

We just assumed sincere voting

But there are very natural situations where a voter should actually lie, even though her goal is to maximize the probability that the group chooses the right alternative!

Example, information cascades-style:

Experimenter has two urns, 10 marbles each

One urn has 10 white marbles ("pure") and the other has 9 green and one white ("mixed")

Three people privately draw one marble and guess what urn it is, and all win money if the majority of them are right

# Insincere Voting

Suppose you draw a white marble

→ Way more likely that urn is pure than mixed

If you draw a green marble

→ Know for sure it's mixed

But what should you guess?

#### First, when will your guess actually matter?

If the two others agree, then your guess doesn't change anything!

Only case where it matters is if they're split

If they're split, someone said mixed, so they know it's mixed!

Then you should guess mixed to break the tie the right way!

Assuming others vote sincerely, you have an incentive to vote insincerely => everyone voting sincerely is **not** a Nash equilibrium

#### Insincere Voting

This is very naturally thought of as a game

Voters are players, guesses are strategies, and they result in certain payoffs

This is highly stylized setting so we can see what's going on

But it happens in the real world too

Consider a jury deliberating on a verdict: guilty or innocent

There is a "best" answer — whether the defendant is actually guilty or innocent

Compare with Condorcet Jury Theorem setup:

- I. Juries require a unanimous vote. Guilty only if everyone says guilty
- 2. In Condorcet, evaluate alternatives just by picking most likely one (if > 1/2 sure, pick it). Here, only pick guilty if sure beyond a reasonable doubt:

 $\Pr\left[defendant \ is \ guilty \mid \ all \ available \ information \right] > z \ \ \text{for some large z}$ 

```
Each juror gets an independent private signal: guilty signal (G-signal) or innocent signal (I-signal)
```

```
They usually get the right signal: P[G-signal \mid defendant guilty] = P[I-signal \mid defendant innocent] = q, q > 1/2
```

Assume prior probability of guilt of 1/2, but doesn't matter

What should a juror do?

- What should a juror do?
- Say you receive an I-signal
  - At first it seems obvious that you should vote to acquit
  - But: conviction criterion is Pr[defendant is guilty | available information] > zso if all the other jurors received G-signals you might still be above that threshold
  - Second, ask yourself key question from before: when does my vote actually matter?
  - Like before, your vote only changes the outcome if everyone except you is voting guilty!
    - If you vote guilty, defendant is found guilty
    - If you vote to acquit, defendant is found innocent

If everyone but you is voting guilty, what is the probability of defendant being guilty?

 $\Pr\left[defendant \ is \ guilty \mid you \ have \ the \ only \ I\text{-}signal}\right] \\ = \frac{\Pr\left[defendant \ is \ guilty\right] \cdot \Pr\left[you \ have \ the \ only \ I\text{-}signal} \mid \ defendant \ is \ guilty\right]}{\Pr\left[you \ have \ the \ only \ I\text{-}signal}\right]}$ 

 $\Pr[you\ have\ the\ only\ I\text{-}signal]$ 

 $= \Pr\left[defendant \ is \ guilty\right] \cdot \Pr\left[you \ have \ the \ only \ I\text{-}signal \ | \ defendant \ is \ guilty\right] + \\ \Pr\left[defendant \ is \ innocent\right] \cdot \Pr\left[you \ have \ the \ only \ I\text{-}signal \ | \ defendant \ is \ innocent\right] \\ = \frac{1}{2} \cdot q^{k-1}(1-q) + \frac{1}{2}(1-q)^{k-1}q.$ 

If everyone but you is voting guilty, what is the probability of defendant being guilty?

 $\Pr\left[defendant \ is \ guilty \mid you \ have \ the \ only \ I\text{-}signal}\right] \\ = \frac{\Pr\left[defendant \ is \ guilty\right] \cdot \Pr\left[you \ have \ the \ only \ I\text{-}signal \mid \ defendant \ is \ guilty}\right]}{\Pr\left[you \ have \ the \ only \ I\text{-}signal}\right]}$ 

$$\Pr\left[defendant \ is \ guilty \ | \ you \ have \ the \ only \ I\text{-}signal\right] \ = \ \frac{\frac{1}{2}q^{k-1}(1-q)}{\frac{1}{2}q^{k-1}(1-q)+\frac{1}{2}(1-q)^{k-1}q} \\ = \ \frac{q^{k-2}}{q^{k-2}+(1-q)^{k-2}},$$

- Since q>1/2,  $(1-q)^{k-2}$  is super small, so the probability goes to 1
- In only case where your vote to acquit matters, you should vote guilty despite your I-signal!

- Intuitively: because of the unanimity rule, you only affect the outcome when everyone else holds the opposite opinion
- Assuming everyone else is as informed as you, and **assuming** independence (remember information cascades!), then the conclusion is that they're probably collectively right
- The result is: assuming everyone else votes sincerely, you have an incentive to vote insincerely
  - All-sincere voting is not an equilibrium
- What is the equilibrium?
  - There are several
  - Most interesting is a mixed equilibrium (randomly disregard I-signal some fraction of the time to correct for possibility that it's wrong)
  - In this equilibrium, probability of convicting an innocent defendant does not go to zero as #jurors goes to infinity!

#### Jury Decisions

- Why do we get such a bad outcome?
- Unanimity is a very harsh constraint.
  - If we relax to only requiring a certain fraction f saying guilty, then the probability that we convict an innocent defendant goes to 0

## Summary

- Voting: synthesizing the preferences of many people into a single group preference
- Many fundamental issues:
  - Condorcet paradox: most natural method (majority rule) can turn a set of reasonable preference relations into an unreasonable one
  - Arrow's Theorem: no general voting system simultaneously satisfies unanimity, IIA, and non-dictatorship.
- Special case: single-peaked preferences
  - Median Voter Theorem says we can get good outcomes
- Jury deliberations: insincere voting can be incentivized



#### Arrow's Theorem

- "Voting profile" = set of k individual rankings
- Criterion I "Unanimity": if there is a pair X and Y for which  $X >_i Y$  for every i, then  $X >_i Y$
- Criterion 2 "Independence of Irrelevant Alternatives" (IIA): the ordering of X and Y should only depend on the relative positions X and Y in individual rankings
  - If we have a bunch of rankings that produces a group ranking with X > Y
  - Then we move some Z around in the individual rankings
  - It should still be the case that X > Y
  - Another way to say it: voting profile restricted to X and Y deletes everything except X and Y. Voting system should order X and Y same way on profile as on profile restricted to X and Y
- Criterion 3 "Non-Dictatorship": the group ranking should not just always be what one particular voter thinks