

Principles of Computer Networks

Tutorial 9

Problem 1 Solution:

a)

| | |
|----------|----------------------------|
| Router z | Informs w, $D_z(x)=\infty$ |
| | Informs y, $D_z(x)=6$ |
| Router w | Informs y, $D_w(x)=\infty$ |
| | Informs z, $D_w(x)=5$ |
| Router y | Informs w, $D_y(x)=4$ |
| | Informs z, $D_y(x)=4$ |

b) Yes, there will be a count-to-infinity problem. The following table shows the routing converging process. Assume that at time t_0 , link cost change happens. At time t_1 , y updates its distance vector and informs neighbors w and z. In the following table, “ \rightarrow ” stands for “informs”.

| time | t_0 | t_1 | t_2 | t_3 | t_4 |
|------|---|---|--|--|--|
| Z | \rightarrow w, $D_z(x)=\infty$ \rightarrow y, $D_z(x)=6$ | | No change | \rightarrow w, $D_z(x)=\infty$ \rightarrow y, $D_z(x)=11$ | |
| W | \rightarrow y, $D_w(x)=\infty$ \rightarrow z, $D_w(x)=5$ | | \rightarrow y, $D_w(x)=\infty$ \rightarrow z, $D_w(x)=10$ | | No change |
| Y | \rightarrow w, $D_y(x)=4$ \rightarrow z, $D_y(x)=4$ | \rightarrow w, $D_y(x)=9$ \rightarrow z, $D_y(x)=\infty$ | | No change | \rightarrow w, $D_y(x)=14$ \rightarrow z, $D_y(x)=\infty$ |

We see that w, y, z form a loop in their computation of the costs to router x. If we continue the iterations shown in the above table, then we will see that, at t_{27} , z detects that its least cost to x is 50, via its direct link with x. At t_{29} , w learns its least cost to x is 51 via z. At t_{30} , y updates its least cost to x to be 52 (via w). Finally, at time t_{31} , no updating, and the routing is stabilized.

| time | t_{27} | t_{28} | t_{29} | t_{30} | t_{31} |
|------|--|--|--|--|---|
| Z | \rightarrow w, $D_z(x)=50$ \rightarrow y, $D_z(x)=50$ | | | | via w, ∞ via y, 55 via z, 50 |
| W | | \rightarrow y, $D_w(x)=\infty$ \rightarrow z, $D_w(x)=50$ | \rightarrow y, $D_w(x)=51$ \rightarrow z, $D_w(x)=\infty$ | | via w, ∞ via y, ∞ via z, 51 |
| Y | | \rightarrow w, $D_y(x)=53$ \rightarrow z, $D_y(x)=\infty$ | | \rightarrow w, $D_y(x)=\infty$ \rightarrow z, $D_y(x)=52$ | via w, 52 via y, 60 via z, 53 |

c) cut the link between y and z.

Problem 2 Solution:

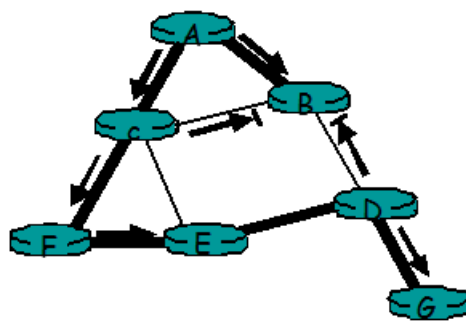
- a) eBGP
- b) iBGP
- c) eBGP
- d) iBGP

Problem 3 Solution:

- a) I1 because this interface begins the least cost path from 1d towards the gateway router 1c.
- b) I2. Both routes have equal AS-PATH length but I2 begins the path that has the closest NEXT-HOP router.
- c) I1. I1 begins the path that has the shortest AS-PATH.

Problem 4 Solution:

The minimal spanning tree has z connected to y via x at a cost of $14(=8+6)$.
 z connected to v via x at a cost of $11(=8+3)$;
 z connected to u via x and v, at a cost of $14(=8+3+3)$;
 z connected to w via x, v, and u, at a cost of $17(=8+3+3+3)$.

Problem 5 Solution:

The thicker shaded lines represent
 The shortest path tree from A to all
 destination. Other solutions are
 possible, but in these solutions, B
 can not route to either C or D from A.