## Question 1 (segmentation)



Figure 1
a) $L_{f}=1$ MByte
$R=1 \mathrm{Mbps}$
Time for the file to move one hop $=L_{f} f R=1 \times 8 \times 10^{6} / 10^{6}=8 \mathrm{sec}$
Total time $=3 \times 8=24$ seconds
b) Number of packets $=4000$
$L_{p}=2 \mathrm{~Kb}$
$R=1 \mathrm{Mbps}$
Let $\mathrm{T}(1)$ denote the time for 1 packet to arrive at destination, i.e. to move 3 hops. $\mathrm{T}(1)=3 \times\left(2 \times 10^{3} / 10^{6}\right)=3 \times 0.002=0.006 \mathrm{sec}$

Packet $n$ arrives at destination 0.002 sec after packet $n-1$ :

$$
\mathrm{T}(n)=\mathrm{T}(n-1)+0.002
$$

By recurrence, $\quad \mathrm{T}(n)=\mathrm{T}(1)+(n-1) \times 0.002$

$$
\begin{aligned}
& \mathrm{T}(n)=0.006+0.002 \times n-0.002 \\
& \mathrm{~T}(n)=0.002 \times n+0.004
\end{aligned}
$$

Hence, $T(4000)=0.002 \times 4000+0.004=8.004 \mathrm{sec}$

## Question 2 (non persistent HTTP)

a) $4 \times \mathrm{RTT}_{0}=4 \times 1.5=6 \mathrm{sec}$ time for name resolution
$\mathrm{RTT}_{1}+\mathrm{RTT}_{1}=2+2=4 \mathrm{sec}$ time for one TCP connection + receiving the HTML file $7 \times\left(\mathrm{RTT}_{1}+\mathrm{RTT}_{1}\right)=7 \times 4=28 \mathrm{sec}$ time for receiving 7 objects, each requires a TCP connection

A total of 38 seconds.
b) $4 \times \mathrm{RTT}_{0}=4 \times 1.5=6 \mathrm{sec}$ time for name resolution
$\mathrm{RTT}_{1}+\mathrm{RTT}_{1}=2+2=4 \mathrm{sec}$ time for one TCP connection + receiving the HTML file $\mathrm{RTT}_{1}+\mathrm{RTT}_{1}=2+2=4 \mathrm{sec} \quad$ time for one TCP connection + receiving 4 objects in parallel $\mathrm{RTT}_{1}+\mathrm{RTT}_{1}=2+2=4 \mathrm{sec} \quad$ time for one TCP connection + receiving 3 objects in parallel

A total of 18 seconds.

Question 3 (steady state) .
Assumptions: Assume that the system is in state $S_{0}, S_{1}, S_{2}, S_{3}$ with probability $P_{0}, P_{1}, P_{2}, P_{3}$ respectively. Also, $\frac{P_{a}}{P_{d}}=\frac{1}{3}$.


Figure 2
Based on Figure 2 and since $\frac{P_{a}}{P_{d}}=\frac{1}{3}$ :

$$
\begin{equation*}
\mathrm{P}_{0}=3 \mathrm{P}_{1} \quad \mathrm{P}_{1}=3 \mathrm{P}_{2} \quad \mathrm{P}_{2}=3 \mathrm{P}_{3} \tag{1}
\end{equation*}
$$

Also, $P_{3}+P_{2}+P_{1}+P_{0}=1$ :
(1) and (2) $==>\quad P_{3}+3 P_{3}+9 P_{3}+27 P_{3}=1$

$$
40 P_{3}=1 \Rightarrow P_{3}=1 / 40 \quad P_{2}=3 / 40 \quad P_{1}=9 / 40 \quad P_{0}=27 / 40
$$

In other words: $P_{3}=0.025 \quad P_{2}=3 * 0.025=0.075 \quad P_{1}=3 * 0.075=0.225 \quad P_{0}=3 * 0.225=0.675$

## Part 1: true

- DNS, the domain name system, is a transport layer protocol.
- A circuit-switched network can guarantee a certain amount of end-to-end bandwidth for the duration of a call.
- HTTP and DNS have in-band control messages.
- Transmission delay is normally shorter than propagation delay for small packets.
- It takes 1 msec for a packet of length 1 Kb to transmit over a link of distance 5 Km , propagation speed of $10^{6} \mathrm{mps}$, and transmission rate of $10^{6} \mathrm{bps}$, neglecting any other delays.
- In BitTorrent, while a peer does not have any chunk it cannot become a top-four uploader for any other peer.
- HTTP response messages can have an empty message body.


## Part 1: false

- If there are 50 users in a system each of which is active only $10 \%$ of the time, the probability of having at least 11 users active at the same time is significantly (e.g. more than 10 times) higher than the probability of having exactly 11 users active.
- The time between any two packet arrivals is called inter-arrival time.
- TLS, the transport layer security, is a transport layer protocol.

