

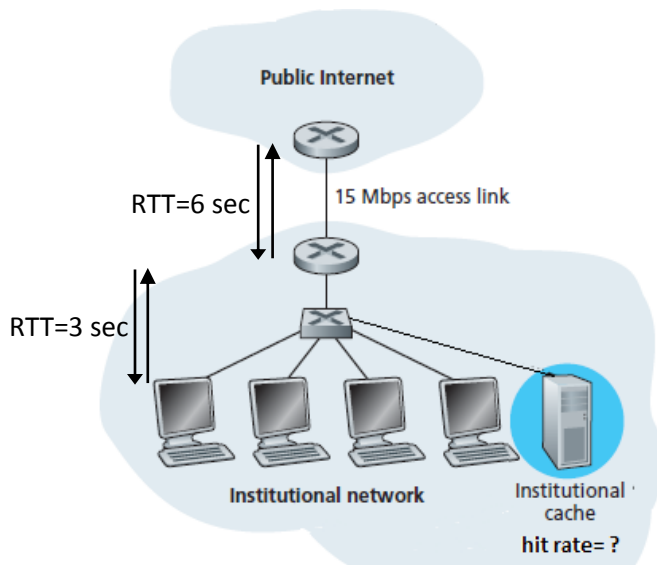
**Question 1 (packet switching):**

a 4 Mbps link  
 each user requires 200 Kbps only 50 percent of the time  
 assume  $2^{10} \cong 1000$

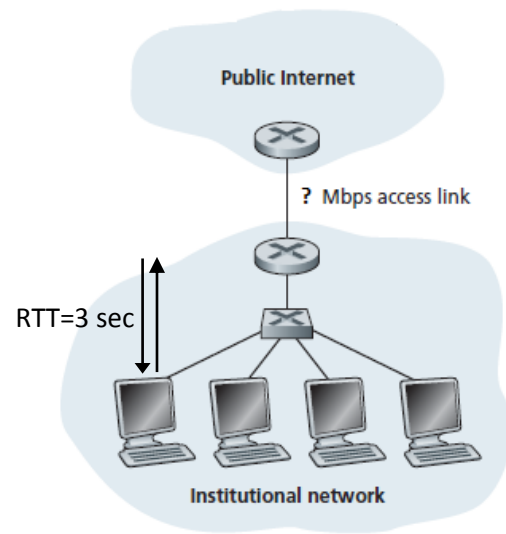
- a) Yes, 50 users can simultaneously use the packet switching network. It just needs congestion control.
- b) Probability of having minimum 8 users (out of 10) active at the same time.

$$\binom{10}{8} \left(\frac{1}{2}\right)^{10} + \binom{10}{9} \left(\frac{1}{2}\right)^{10} + \binom{10}{10} \left(\frac{1}{2}\right)^{10} = \left(\frac{10 \times 9}{2} + \frac{10}{1} + 1\right) \left(\frac{1}{2}\right)^{10} = 56 \left(\frac{1}{2}\right)^{10} \cong 0.056$$

**Question 2 (cache):**



**Figure 1a**



**Figure 1b**

- a) Let  $x$  denote the cache hit rate,  $RTT_0$  denote the local round trip time, and  $RTT_1$  the round trip time between the routers.  $RTT_0=3$ ,  $RTT_1=6$ .

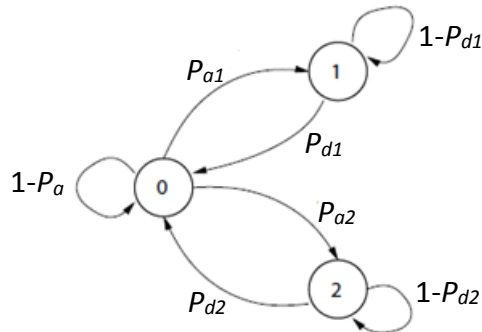
$$\begin{aligned} x \cdot RTT_0 + (1-x) (RTT_0+RTT_1) &< 6 \\ 3x + 9(1-x) &< 6 \\ \Rightarrow x &> 0.5 \end{aligned}$$

- b) Objective:  $RTT_0 + RTT_1 = 6$  sec.  
 Assumptions:  $RTT_0=3$  sec, and  $RTT_1 = 6$  sec when  $R= 15$ Mbps.  
 Hence, in order to have  $RTT_1=3$  sec,  $R$  should be 30 Mbps.

**Question 3 (steady state):**

Assumptions: Assume that the system is in state  $S_0, S_1, S_2$  with probability  $P_0, P_1, P_2$  respectively. Also,

$$\frac{P_{a1}}{P_{d1}} = \frac{1}{2} \text{ and } \frac{P_{a2}}{P_{d2}} = \frac{1}{3} \text{ and } P_a = P_{a1} + P_{a2}.$$



**Figure 2**

Based on Figure 2 and since  $\frac{P_{a1}}{P_{d1}} = \frac{1}{2}$  and  $\frac{P_{a2}}{P_{d2}} = \frac{1}{3}$ :

$$P_0 = 2P_1 \quad P_0 = 3P_2 \quad (1)$$

$$\text{Also, } P_0 + P_1 + P_2 = 1 \quad (2)$$

$$(1) \text{ and } (2) \implies P_0 + P_0/2 + P_0/3 = 1$$

$$\implies 11P_0/6 = 1 \implies P_0 = 6/11 \quad P_1 = 3/11 \quad P_2 = 2/11$$

$$\text{In other words: } P_0 \cong 0.55 \quad P_1 \cong 0.27 \quad P_2 \cong 0.18$$

**Part 1: true:**

- ⇒ HTTP is stateless.
- ⇒ DNS, the domain name system, is an example of a client-server architecture.
- ⇒ DHT, the distributed hash table, can be designed so that the number of messages per query is  $O(\log N)$ , where  $N$  is the number of peers.
- ⇒ After a packet arrives to a switch, its first bits can be propagated while remaining bits are being processed.

**Part 1: false:**

- ⇒  $\mu$ TP is a transport layer protocol.
- ⇒ In order to join a BitTorrent torrent, a peer must have at least one chunk.
- ⇒ TCP provides minimum data transmission rate between processes.
- ⇒ It takes 1 msec for a packet of length 1 Kbits to propagate over a link of distance 5 Km, propagation speed of  $10^6$  mps, and transmission rate of  $10^6$  bps, neglecting any other delays.
- ⇒ Suppose that a file size is 1000 Kbits and the path from Host A to Host B has two links of rates  $R_1 = 1$  Mbps and  $R_2 = 500$  Kbps. The throughput for the file transfer is 2 seconds.
- ⇒ A circuit-switched network is well suited for applications in which the transmission rate is unknown and bursty.