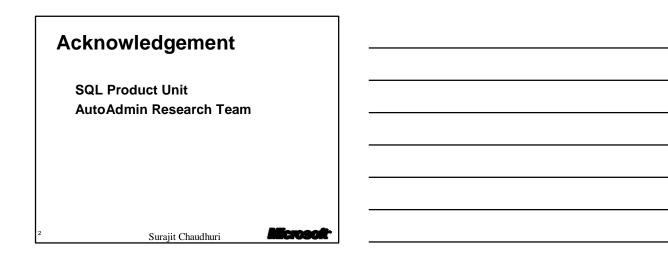
Self-Managing DBMS Technology at Microsoft

Surajit Chaudhuri Microsoft Research

Microsoft[.]



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Easy Solutions	
Throw more hardware	
Use this with caution	
Where do you throw hardware?	
Rules of Thumb approach	
Finding them is harder than you think	
May simply not exist – oversimplified wrong solutions are not helpful	
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Microsoft's Early Focus on Self-Managing Technology

1998: SQL Server 7.0 launch towards a self-tuning database system:

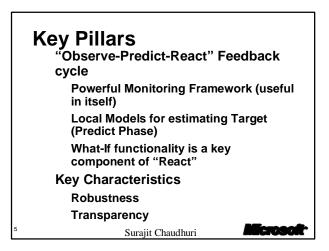
Eliminate outright many knobs and provide adaptive control

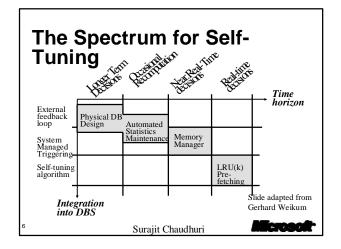
Dynamic Memory Management Auto Stats, Auto Parallelism and Space Management Index Tuning Wizard

1996: AutoAdmin Project at Microsoft Research – exclusive focus on self tuning

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Microsoft







Monitoring SQL Server Activities

Monitoring Tools

7

Microsoft Operations Manager

Track Connectivity, Free space, Long Running Jobs, PERFMON Reporting

Best Practices Analyzer

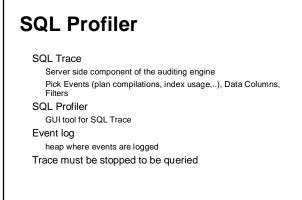
Detect common oversights in managing a SQL Server installation

Simple UI, Rules metadata (70+), Reporting File Compression, File Placement, Index frag

Dedicated Admin connection in SS 2005

Connect even to a "hung" server (uses reserved scheduler, port & resources) Microso

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Need for More
pansparencyMajority of case time is spent diagnosing the
problem (allocation errors, perf degradation)
60% in data collection, 40% in data analysisDependence on ReprosDifficult to ID some performance issues
Unacceptable to many customers
End User experienceHelp requested for cases which don't resolve
within 30 minsFull dump requested on ~40%

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FILL BOOK

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Dynamic Management Views in SQL Server 2005

Simple queries now solve many scenarios (Live in memory stats)

low level system (server-wide) info such as memory, locking & scheduling

Transactions & isolation

Input/Output on network and disks Databases and database objects

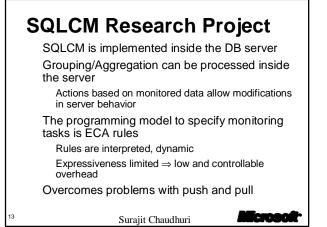
Populate a Data Warehouse

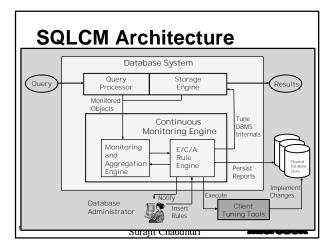
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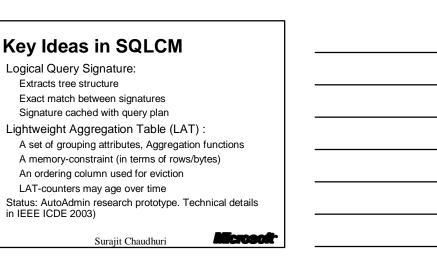
Example: Dynamic Management Views Sys.dm_exec_requests – currently running requests Sys.dm_exec_query_stats

One row per query plan currently in the cache Min, max, avg, last; – Physical reads, logical reads, physical writes; Execution count; First and last execution times "Performance Statistics" Trace event

Log "query_stats" for plans which are removed from the cache







Workload Analysis

Variety of tasks leverage workload DBA (ad-hoc analysis) Physical design tuning tools Approximate query processing

Workload typically gathered by logging events on server

Workloads can be very large Few DBAs can eyeball 1GB workload file! Few tools can scale

Need infrastructure for summarizing and analyzing workloads

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Approaches to Workload Analysis

Populate a schematized database

Model as multi-dimensional analysis problem Good for ad-hoc analysis using SQL and OLAP Insufficient support for summarization

Summarizing Workload:

Random sampling Application specific workload clustering (SIGMOD 2002)

Plug-in "distance" function, adapt K-Mediod clustering Novel declarative primitives (VLDB 2003)

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Estimating Progress of SQL Query Execution

Decision support systems need to support long running SQL queries

Today's DBMS provides little feedback to DBA <u>during query execution</u>

Goal: Provide <u>reliable progress estimator</u> during query execution

Accuracy, Fine Granularity, Low Overhead, Monotonicity, Leverage feedback from execution Status: AutoAdmin Research Project and prototype: technical details in SIGMOD 2004

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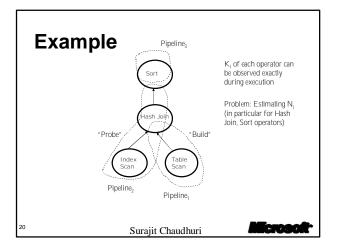
Microso

$\begin{array}{l} \textbf{Modeling Total Work} \\ \text{Want a simpler model than query optimizer's cost} \\ \text{estimate} \\ \text{Query execution engines use iterator model} \\ \text{Total work = Total number of GetNext() calls} \\ \text{Let N}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let K}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let K}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let K}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be total number of GetNext() calls for Op}_i \\ \text{Let M}_i \text{ be$

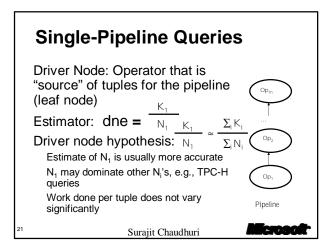
where c_i is relative weight of Op Problem: Estimating $N_i\, \text{during query execution}$

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F-TO-BOOK









Other Key Considerations Leverages execution information Observed cardinalities (K,'s) Algebraic properties of operators Internal state of the operator Spills due to insufficient memory Model as a new (runtime) pipeline Trade-off between guaranteeing monotonicity and accuracy Non-uniform weights of operators

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Recap of Monitoring Highlights Transparency of current server state crucial for easing DBA tasks, supported by DMVs Online aggregation of server state can support a monitoring framework (SQLCM)

Logging of workloads as well as server events using SQL Profiler is crucial for offline analysis

Tool to estimate progress of queries

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Self-Tuning Memory Management

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Dynamic Self Tuning Memory Manager

SQL 7.0 pioneered idea of dynamic selftuning memory

Sufficient memory set aside so that Windows and other applications can run without hiccups Amount depends on system load

Observe:

Query Windows for the amount of free physical memory periodically

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Considers page life expectancy for the buffer pool

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Self-Tuning Memory Manager

Predict: Available memory compared to required threshold of Target Pages (PERFMON values consulted)

No explicit model-based prediction

Takes physical memory size into account

React:

Keep a given number of free pages (for new allocation requests) at all times

Grab if low page life expectancy

If memory pressure from OS, free up buffers

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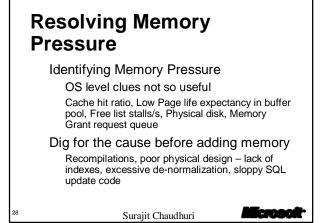
Memory Management by Query Execution Engine Among competing queries

Within a query Among parallel threads Nodes of a plan Phases within an operator Give each query, once admitted to execution, adequate memory

Waiting memory, Waiting operators Preempt on demand

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Alicrose



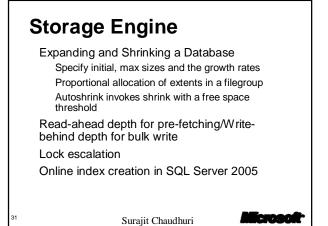


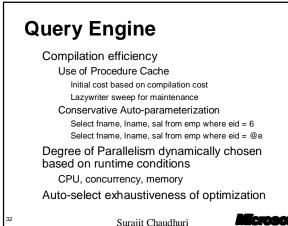
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Automatic Checkpointing Uniform time interval is not ideal

Based on number of records in the log Specified recovery interval – max time SQL Server should take for restart Log manager estimates if it is time for checkpointing For simple recovery model Log 70% full Restart may take more than recovery interval





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Why Statistics Management?

Having "right" statistics is crucial for good quality plans. When to build statistics? Which columns to build statistics on?

How to build statistics on any column efficiently?

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Auto Statistics in SQL Server

Created dynamically at query compilation time

On single table columns for which optimizer needs to estimate distribution Uses sampling of data to create statistics

Statistics auto-maintained

Novel feature supported since SQL Server 7.0

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Uniform vs. Block-Level Sampling

Uniform random sampling is too expensive.

Block-level sampling:

Pick a few blocks at random and retain all tuples in those

Block level sampling is efficient but tuples may be placed in blocks arbitrarily Reduced quality of the resulting estimate

uced quality of the resulting esti

AutoUpdate of Statistics

Triggered by Query Optimization Involves only a subset of the columns in the query

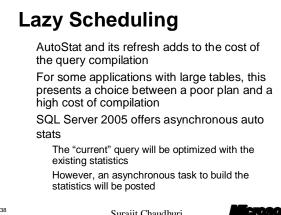
Refreshed when a certain fraction (roughly) of the rows have been modified

Uses rowmodctr information to check if threshold has been reached

Statistics that are auto-created are aged and retired if appropriate.

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Frontiers for Further Thinking

Determining the appropriate Block Level Sampling Identifying the interesting subset of statistics for a query

Statistics on views and query expressions Leveraging execution feedback

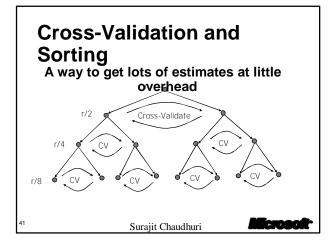
Remaining slides in this part are on some research ideas being pursued at Microsoft

Adaptive 2-phase approach for Block Level Sampling

Get initial sample While sorting get error estimate for r/2, r/4, r/8 ... etc. Find the best-fit curve of the form c/sqrt(r) through these points Read off the required sample size Experimentally found to almost always reach the error target or very close. AutoAdmin research prototype, SIGMOD 2004

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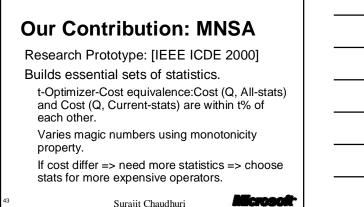
Recommending Base-Table Statistics Find subset as good as having all

statistics ("essential" set) Depends on workload, data distribution, optimizer... Determining an essential set is non-

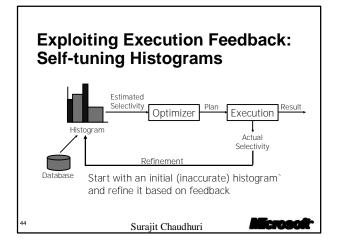
trivial.

"Chicken-and-egg" problem: cannot tell if additional statistics are necessary until we actually build them! Need a test for equivalence *without* having to build any new statistics

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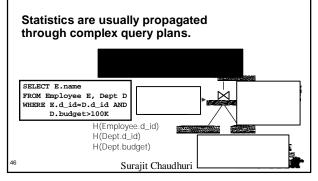
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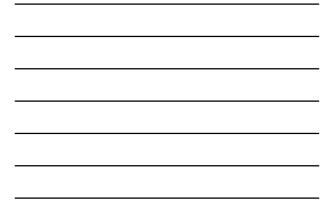


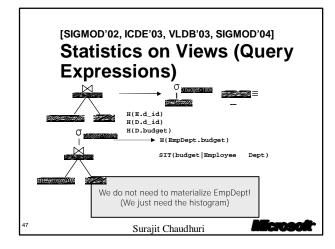
Self Tuning Histograms: STGrid and STHoles Assume uniformity and independence until execution feedback shows otherwise (no data set examination) Exploit workload to allocate buckets. Query feedback captures uniformly dense regions Differences: Bucket structure and refining STGrid: Multidimensional Grid [SIGMOD'99].

STHoles: Bucket nesting [SIGMOD'01].

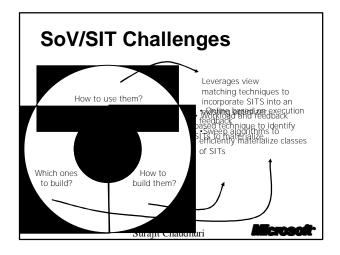
Are base-table statistics sufficient?













Self-Tuning Physical Database Design

Microsoft SQL Server Milestones

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SQL Server 7.0: Ships index tuning wizard (1998): Industry's first

SQL Server 2000: Integrated recommendations for indexes and materialized (indexed) Views: Industry's first

SQL Server 2005: Integrated recommendations for indexes, materialized views, and partitioning, offering time bound tuning, Industry's first

Results of collaboration between AutoAdmin Research and the SQL Server teams

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Key Insights Robustness was a design priority Every system is different – track wor

Every system is different – track workloads (VLDB 1997)

"What-If" API for DBMS (SIGMOD 1998) is key to driving selection of physical design Efficient search for physical design (VLDB 1997, 2000, 2004)

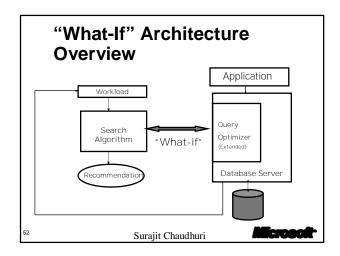
Significant thinking on system usability (VLDB 2004)

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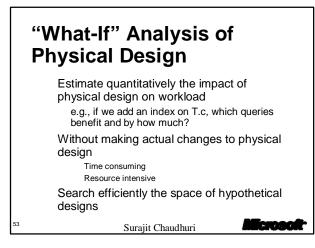
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Crosoft







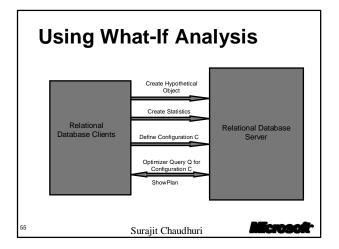
Realizing "What-If" Indexes

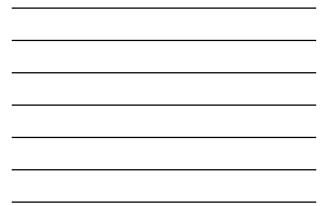
Query Optimizer decides which plan to choose given a physical design

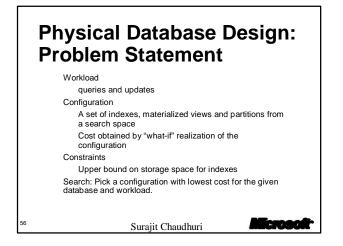
Query optimizer does not require physical design to be materialized Relies on statistics to choose right plan Sampling based techniques for building

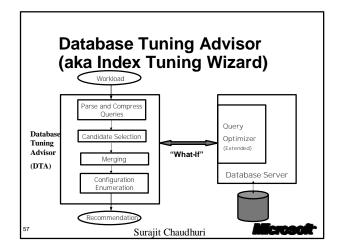
Sufficient to fake existence of physical design

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Some Key Ideas Prefiltering of search space Adapt cost-based frequent itemset idea from data mining (VLDB 2000) Quantitative analysis at per query level to isolate candidates Watch out for over-fitting View Merging Search Efficiency crucial Server bears the cost of "searching" as we ping the optimizer, Robustness - unaffected by most optimizer changes

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DTA for Microsoft SQL Server 2005 Time bound tuning Complete tuning in batch window Range partitioning recommendations Integrated Recommendation with Indexes and $\ensuremath{\mathsf{MVs}}$ Manageability: Can recommend "Aligned" partitioning User-specified configuration (USC) Exposes "What-if" analysis Manageability: Allows specifying partial configuration for tuning Input/Output via XML Public schema: http://schemas.microsoft.com/sqlserver/2004/07/dta/dtaschema.xsd More scriptable Easy for ISVs to build value added tools on top Microsoft ri

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DTA: Microsoft SQL Server 2005
Production/Test Server Tuning
Exploit test server to reduce tuning load on production server
Recommendation same as if tuning done on production server
Servers need not be H/W identical
Improved Performance and Scalability
Workload compression
Reduced statistics creation
Exploit multiple processors on server
Scaling to large schema
Multi-database tuning
Recommends online indexes
Drop-only mode
Clean up unused indexes, MVs
More details in VLDB 2004 paper
¹⁰ Surajit Chaudhuri Alicrosoft

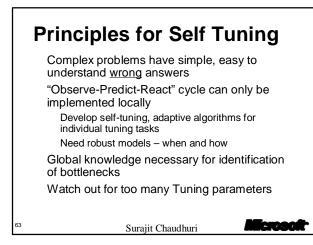
Lessons for Self-Tuning and Rethinking System Design

61

The Spectrum for Self-Tuning Heren le Les Ces Time horizon External feedback loop Physical DB Design Automated Statistics Maintenanc System Managed Triggering Memory Manager Self-tuning algorithm LRU(k) chi 1 Slide adapted from Gerhard Weikum Integration into DBS Microsoft Surajit Chaudhuri

Microsoft





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Rethinking Systems: Wishful Thinking?

VLDB 2000 Vision paper (Chaudhuri and Weikum 2000)

Enforce Layered approach and Strong limits on interaction (narrow APIs)

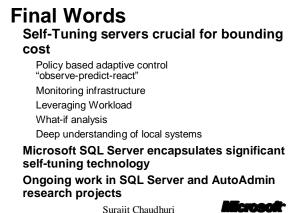
Package as components of modest complexity Encapsulation must be equipped with self-tuning

Featurism can be a curse

Don't abuse extensibility - Eliminate 2nd order optimization

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Microsoft SQL Server Self Tuning Technology Talks

Vivek Narasayya "Database Tuning Advisor for Microsoft SQL Server 2005" (Industrial Session 4, Thu)

David Campbell "Production Database Systems: Making Them Easy is Hard Work" (industrial Session 6, Thu)

Self-Tuning Overview **Papers**

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Self-Tuning Physical Design

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Agrawal S., Chaudhuri S., Kollar L., Marathe A., Narasayya V. and Syamala M. Database Tuning Advisor for Microsoft SQL Server 2005. VLDB 2004 Surajit Chaudhuri . 1 -

Statistics Management

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Monitoring and Workload

Analysis and Management Chaudhuri S., König, A., and Narasayya V. SQLCM: A Continuous Monitoring Framework for Relational Database Engines. *ICDE 2004*. Chaudhuri S., Narasayya V., and Ramamurthy, R. Estimating Progress of Execution for SQL Queries. SIGMOD 2004.

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Chaudhuri S., Ganesan P., and Narasayya V. Primitives for Workload Summarization and Implications for SQL. *VLDB 2003*

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Microsoft[.]