CSC458

Peer-to-Peer Systems

This Time

- P2P systems
- Focus
 - How do P2P systems work?
- Topics
 - Unstructured P2Ps
 - Structured P2Ps: Distributed Hash Tables
 - Comparison of two distributed systems: WWW and a DHT (Chord)

Napster (circa 2002)

Unstructured P2Ps



Napster Architecture

- Dedicated, well-provisioned servers
- Peers connect to one server
- Servers index the content of all peers
- To query:
 - Send query to server, wait for reply
- To download:
 - Over HTTP from one peer (uploader) to another (downloader)
 - No server involvement
- Bottom line: Napster was order to shut-down in 2002

Gnutella (circa 2002)



Gnutella

Gnutella ver. 0.4 Architecture

- No servers, only peers
- Upon joining, peer needs to know at least one other peer
 Through out-of-band mechanisms
- Once joining one peer, send a Ping broadcast
 - Discover additional peers to connect to them <-- flooding!!
- To query:
 - Send a Query broadcast <-- flooding!!</p>
- To download:
 - Over HTTP from one peer (uploader) to another (downloader)
- Bottom line: Unscalable protocol (based on flooding)

Kazaa (today)



Kazaa Architecture

- Some peers act as servers (aka ultrapeers/supernodes)
- Upon joining, peer needs to know at least one supernode
 Through out-of-band mechanisms
- Supernodes index the content of all their peers (like Napster)
- To query:
 - Send query to supernode, wait for reply
- To download:
 - Over HTTP from one peer (uploader) to another (downloader)
 - No supernode involvement
- Bottom line: more scalable architecture, servers are not dedicated!!

Structured P2Ps or Distributed Hash Tables (DHTs)

DHT: basic idea



DHT: basic idea



Neighboring nodes are "connected" at the application-level

DHT: basic idea



Operation: take key as input; route messages to node holding key

DHT: basic idea



Operation: take key as input; route messages to node holding key

DHT: basic idea



Operation: take key as input; route messages to node holding key

DHT: basic idea



Operation: take key as input; route messages to node holding key

DHT: basic idea



Operation: take key as input; route messages to node holding key

How to design a DHT?

- · State Assignment:
 - what "(key, value) tables" does a node store?
- Network Topology:
 - how does a node select its neighbors?
- Routing Algorithm:
 - which neighbor to pick while routing to a destination?
- · Various DHT algorithms make different choices
 - Chord, CAN, Pastry, Tapestry, Plaxton, Viceroy, Kademlia, SkipNet, Symphony, Koorde, Apocrypha, Land, ORDI ...

State Assignment in Chord DHT



- · Nodes are randomly chosen points on a clock-wise Ring of values
- Each node stores the *id space* (*values*) between itself and its predecessor

Chord Topology and Route Selection



- · Neighbor selection: ith neighbor at 2ⁱ distance
- Route selection: pick neighbor closest to destination

Chord Architecture

- N: number of peers in the system
- Each node has O(log n) neighbours
- Each query traverses O(log n) hops to the destination
- Each routing table has O(log n) size
- Very scalable!

What Chord style systems do:



Lessons

- Unstructured peer-to-peer systems:
 - Napster: lookup centralized, file Xfer P2P
 - Gnutella: lookup + file Xfer P2P
 - Kazaa: lookup peer to supernode, file Xfer P2P
- DHTs:
 - Chord: lookup based on a distributed hash table
 - Very scalable architecture
- DHTs collapse naming and addressing
- DHTs collapse routing and lookup