

CSC2231: Security in P2Ps

<http://www.cs.toronto.edu/~stefan/courses/csc2231/05au>

Stefan Saroiu
Department of Computer Science
University of Toronto

Outline

- **Power-law networks and flat overlays**
- **Sybil attack**
- **P2P routing attacks**
- **Discussion**

Outline

- **Power-law networks and flat overlays**
- Sybil attack
- P2P routing attacks
- Discussion

Power-Law Networks are here to Stay

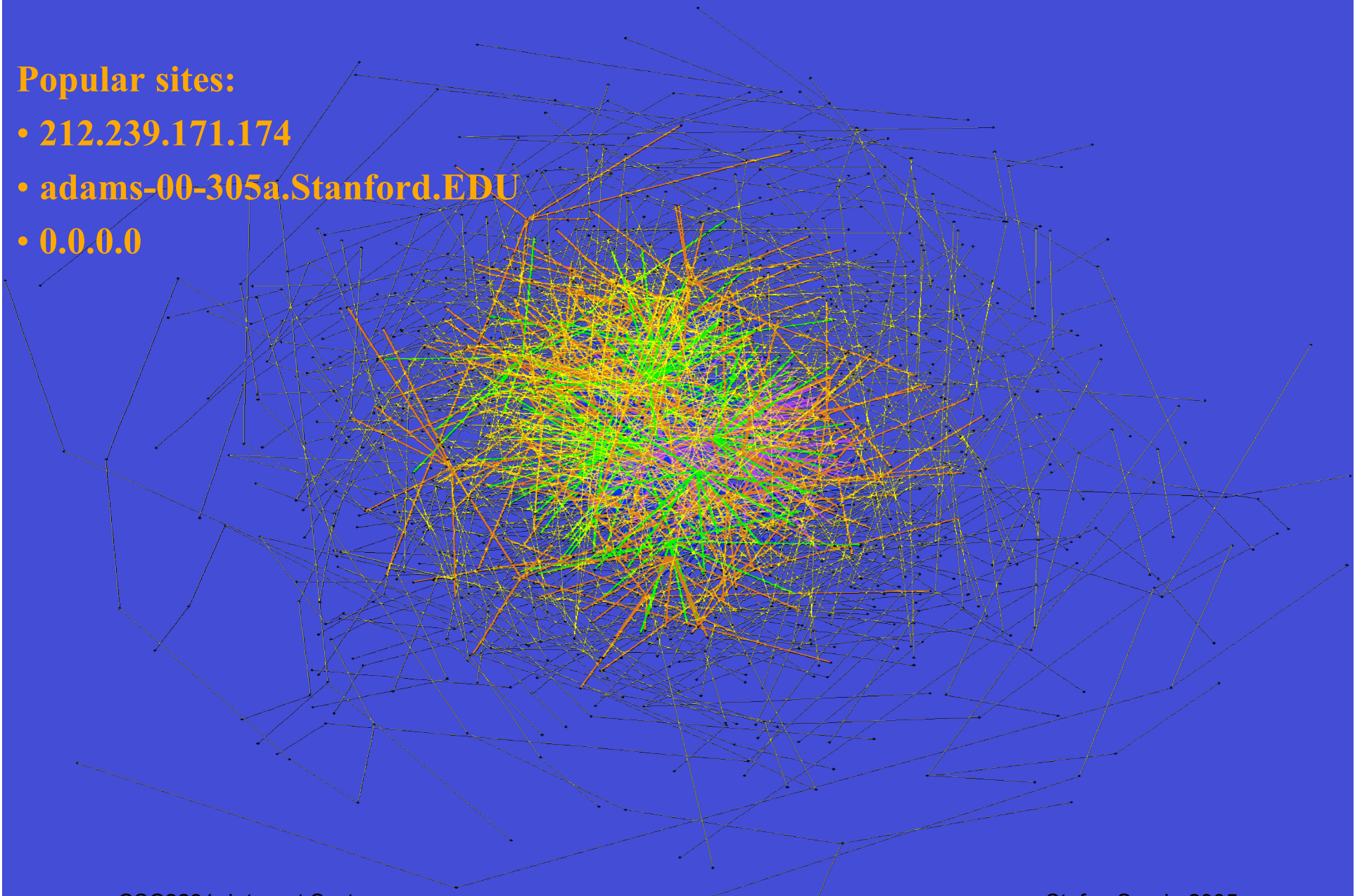
- **Barabasi and Albert showed that networks which...**
 - grow by continuous addition of new nodes
 - exhibit preferential attachment (likelihood of connecting to a node depends on the node's degree)
- **...power-law distribution of vertex degree**
- **Internet, WWW, Gnutella**

Resilience to Failures

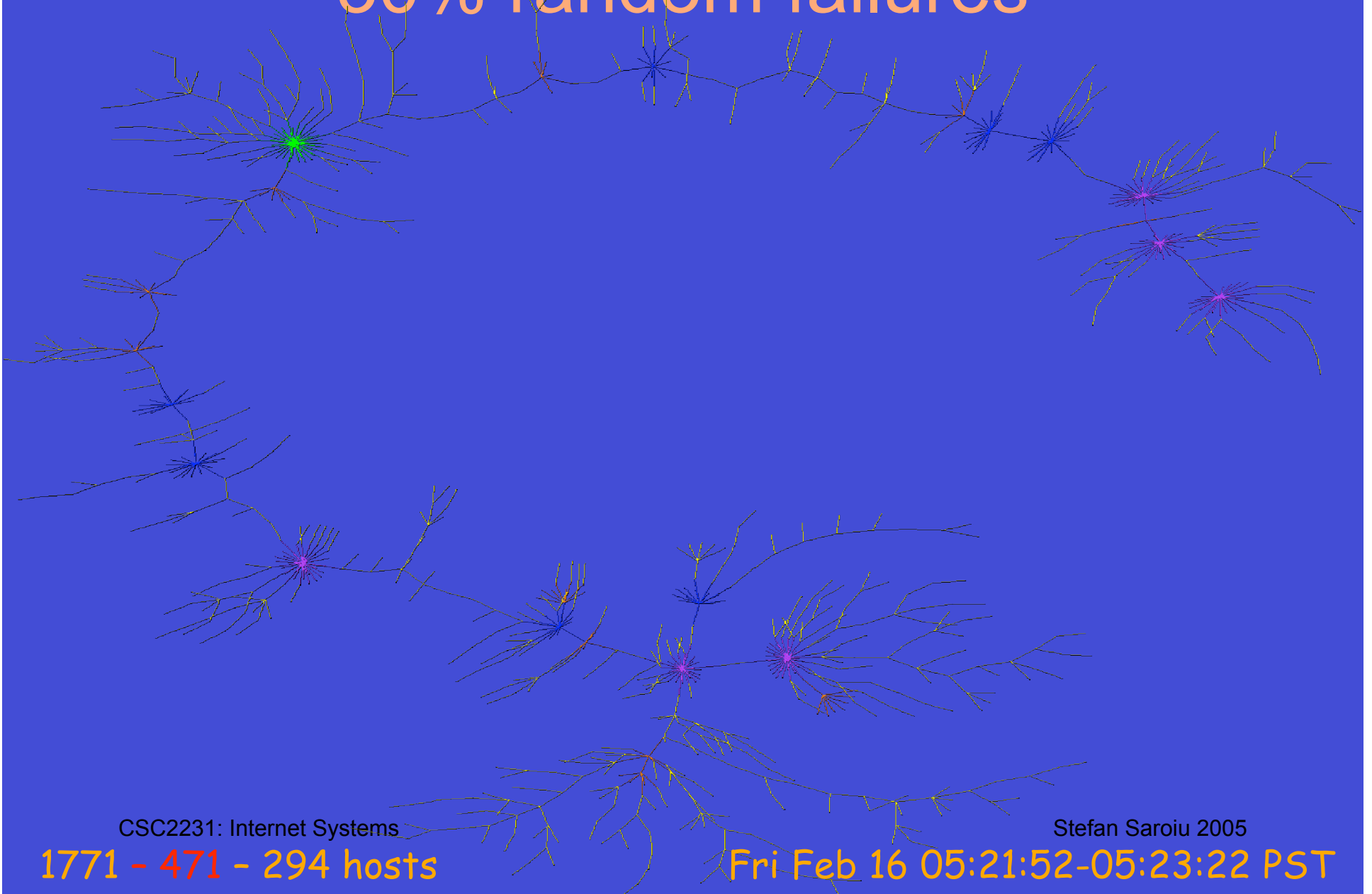
- **Power-law networks (Cohen et al.):**
 - very resilient in face of random node failures
 - a giant spanning cluster still exists
 - fairly resilient in face of cascading failures
 - very vulnerable in face of orchestrated attacks (towards high-degree nodes)

Popular sites:

- 212.239.171.174
- adams-00-305a.Stanford.EDU
- 0.0.0.0



30% random failures



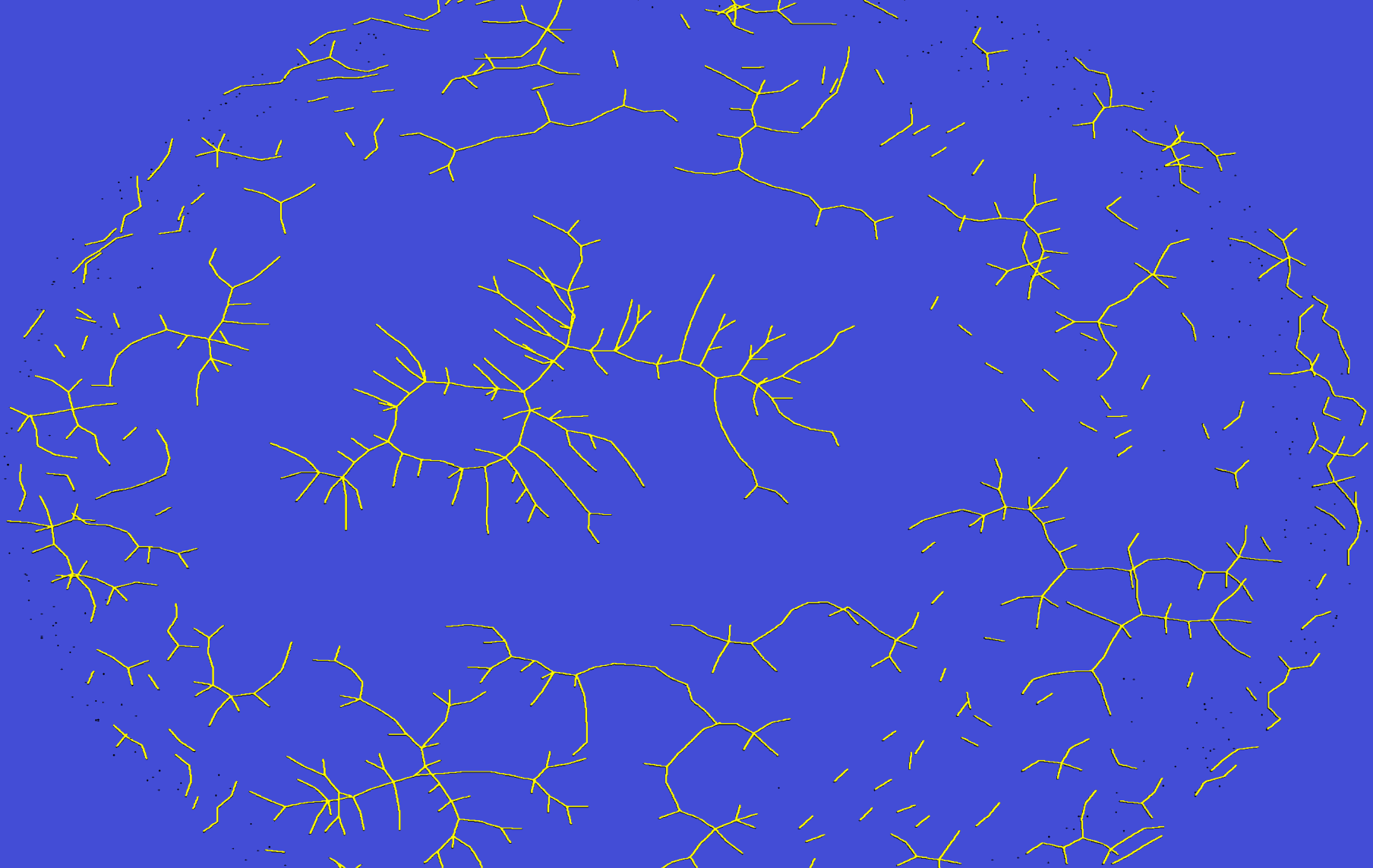
CSC2231: Internet Systems

1771 - 471 - 294 hosts

Stefan Saroiu 2005

Fri Feb 16 05:21:52-05:23:22 PST

4% orchestrated failures



CSC2231: Internet Systems

1771 - 63 hosts

Stefan Saroiu 2005

Fri Feb 16 05:21:52-05:23:22 PST

Outline

- Power-law networks and flat overlays
- **Sybil attack**
- P2P routing attacks
- Discussion

Fundamental problems for P2P

- **One can have, some claim, as many electronic personas as one has the time and energy to create. – *Judith S. Donath.***
- **“Sybil attack....”**
 - Correct assumption: attacker has access to infinite resources
 - It registers *many times* using different identities
 - Cannot distinguish whether these are real or not. Real clients get poor service
 - Impossible to get around in a fully decentralized manner



Identities and Entities

- **Ideally want one to one mapping**
 - Virtual servers idea was a many-to-one
- **How to convince that two different identities correspond to two different entities**
 - Perform a task that no single entity can perform
 - Have others vouch that these are two different entities

Observations on performing tasks

- **The # of identities one entity can create is proportional to the ratio of an entity's resources to the resources of the weakest peer**
 - Can use computational resources (i.e., solve a puzzle)
 - Can use communication resources (i.e., use large packets)
 - Can use storage resources (i.e., challenge large data)
- **If the check on entities is not done simultaneously, one entity can create an infinite # of identities**
 - Can't be used for computational, communication resources
 - Could be used for storage, however it's extremely wasteful

Observations on Vouching

- **If one accepts an identity vouched by q accepted identities, a malicious set F of nodes can create an arbitrary number of identities if $|F| > q$**
- **If the accepting identities is not coordinated in time across the system, one can create an infinite # of identities**

Outline

- Power-law networks and flat overlays
- Sybil attack
- **P2P routing attacks**
- **Discussion**

P2P Routing

- **A key identifier space**
- **A node identifier space**
- **Rules for associating keys to nodes**
- **Per-node routing tables**
- **Rules for forwarding packets to neighbors**
- **Rules for updating routing entries when nodes join/leave**

- **Routing attacks:**
 - Attacks on the ID space (manipulating IDs)
 - Attacks on the routing tables (manipulating latencies/proximity)
 - Attacks on forwarding protocol (ignore messages)

Node ID Assignment

- **Example of possible attacks:**
 - Surround a target node
 - Partition the overlay network
 - Become the root of a certain document
- **How are IDs generated?**
 - Randomly in FreeNet
 - Hashing something in Chord/Pastry
 - Both are easy to manipulate
- **Solution: centralized certification authority for Ids**
 - Single point-of-failure for the system (DoS attacks)
 - Small overlay networks are problematic
 - Dynamic IDs DHTs problematic

Attacks on Routing Tables

- **Fake being the closest node**
- **If f faulty nodes, probability of receiving bad update**
 - $f + (1-f)f$ quite a bit
- **Solution: use a backup routing table**
 - trade complexity for security
 - Sounds like a really bad idea
- **Potential issues:**
 - Bootstrapping nodes
 - Bugs
 - Probability of receiving bad updates is still $f + (1-f)f$

Attacks on forwarding

- **Attacks:**
 - Ignore forwarding requests
 - Be the root of a key and ignore lookups
 - $(1-f)^{\log N}$
 - $f = 10\%$, routing fails 35% of the time
- **Solution: failure test**
 - Have several roots for a document
 - The attack only works when faulty nodes are becoming the roots for the document
 - This means that the root set ID density is higher than the node ID's density in the overlay

Additional Attacks

- **Content pollution**
 - Fair-exchange is hard to enforce in a fully-decentralized manner
 - A central authority is a great solution (i.e., escrow)
- **Rapid joins and leaves**
- **Lots of key inserts**
- **Lots of negative lookups**
- **Inconsistent behavior**
 - “good citizen” to half of the network, “bad citizen” to the other half

High-level Lesson from P2P Security Paper

- **Simple idea:**
 - Build a P2P around a verifiable invariant
 - Check for invariant to enforce security
- **Observation:**
 - Very hard to do in practice
 - A centralized check for the invariant would simplify the problem tremendously

Discussion

- **How decentralized should a P2P be?**

Discussion

- **How decentralized should a P2P be?**
 - CDNs are semi-decentralized:
 - Central authority, management, etc..
 - Can be used to bootstrap security/trust

Discussion

- **How decentralized should a P2P be?**
 - CDNs are semi-decentralized:
 - Central authority, management, etc..
 - Can be used to bootstrap security/trust
- **Is the Internet decentralized?**
 - If yes, how can it avoid these problems?
 - How come BGP works?