CSC2209 Computer Networks

#### **DCCP = Unreliable TCP**

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### **Treat Your Advisor Like Your Grandparents!**

## **Treat Your Advisor Like Your Grandparents!**

- Advisors have very short attention spans
- They always forget what they said and to whom
- Every time you see them, spend time:
  - Introducing yourself
  - Telling them what you work on
  - Ignoring them telling the same story over and over again
    - Act surprised!
- Also, make sure you visit them at least once in awhile

### **Project Proposals**

- Next deadline:
  - Friday, November 19th: project report
- Websites up by the end of the week
- Take this seriously!

- Analyzing and Enhancing Content Availability in BitTorrent-like P2P. Nadeem, Jin, Di.
- *Large-Scale Measurement of Bluetooth Activity*. Ivan, Alireza.
- Non-Interfering VMs. Tom.
- *Proxy Support for Application Replay.* Vladan.
- Cooperative Browser Caches. Robert.
- Something related to Spoofing & Phishing. Shvet.

• Waqas, Andrew, Fareha?!

## **Take-away Messages?**

Congestion Management

• DCCP

## **Take-away Messages?**

- Congestion Management
  - Application-level API for congestion control

- DCCP
  - How to build an unreliable, congestion control protocol

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  Streaming media, VoIP, games, video-conferencing
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- What does YouTube use?
  - Streaming media works just fine with TCP, unless it's live
- How much live content is there on the Internet?
  - Not much: is video-conferencing the only app?



- Pure Seq #, not bytes
- Each packet carries a Seq #
- Seq # increase per packet
- Pure Acks also consume Seq #







### **Two congestion control schemes**

- CCID 2: TCP-like
  - Slow-start, timeouts like TCP
  - Additionally, congestion control for ACKs
- CCID 3: equation-based
  - Receiver measures loss rate and feeds back to sender
  - Sender uses feedback messages to estimate RTT
  - Loss rate and RTT are fed into equation
  - Sender adjusts actual rate to match expected rate
- When should one use CCID 2 vs. CCID 3?

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  - Loss rate and RTT are fed into equation
  - Sender adjusts actual rate to match expected rate
- When should one use CCID 2 vs. CCID 3?
  - Whenever reacting to congestion must be abrupt
    - How about CBR vs. online games?
  - Many unknown, poorly-understood issues are left

# **Adding Reliability Back**

• How would you add reliability to DCCP?

# **Adding Reliability Back**

- How would you add reliability to DCCP?
  - FEC
  - Bloom filters
  - Likely very different than TCP...
- How about in-order delivery?

## **Congestion Manager Original Motivation**

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- 1. Many short Web Xfers
- 2. Web accelerators
- 3. Application heterogeneity: many apps over UDP to avoid TCP's reliability mechanism

• Are any of these motivations valid today?

### **Congestion Manager Current Motivation**

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- Congestion is located at the end-points today
  - CM manages flows to different destinations but sharing same bottleneck
- Should we expose API to applications?