

CSC 2515: Structured Prediction

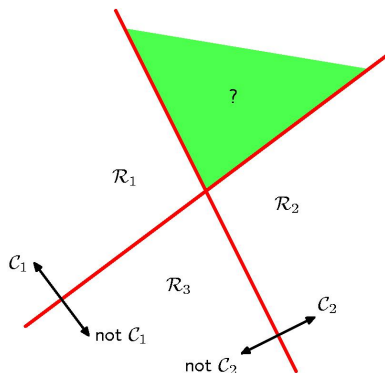
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University of Toronto

March 16, 2015

Discriminant Functions for $K > 2$ classes

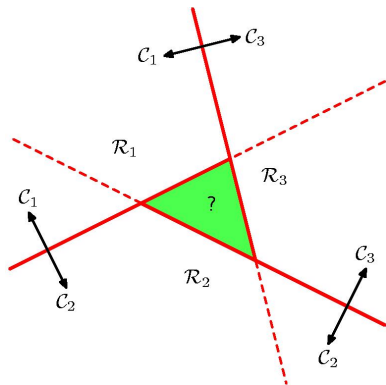
- Use $K - 1$ classifiers, each solving a two class problem of separating point in a class C_k from points not in the class.
- Known as **1 vs all** or **1 vs the rest** classifier



- **PROBLEM:** More than one good answer!

Discriminant Functions for $K > 2$ classes

- Introduce $K(K - 1)/2$ two-way classifiers, one for each possible pair of classes
- Each point is classified according to majority vote amongst the disc. func.
- Known as the **1 vs 1 classifier**



- **PROBLEM:** Two-way preferences need not be transitive

K-Class Discriminant

- We can avoid these problems by considering a single K-class discriminant comprising K functions of the form

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- In this lecture we will look at generalizations of this idea.

- Introduction to Structured prediction
- Inference
- Learning

What is structured prediction?

- In "typical" machine learning

$$f : \mathcal{X} \rightarrow \mathbb{R}$$

the input \mathcal{X} can be anything, and the output is a real number (e.g., classification, regression)

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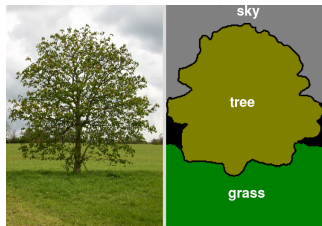
the input \mathcal{X} can be anything, and the output is a **complex** object (e.g., image segmentation, parse tree)

- In this lecture \mathcal{Y} is a discrete space, ask me later if you are interested in continuous variables.

Structured Prediction and its Applications

We want to predict multiple random variables which are related

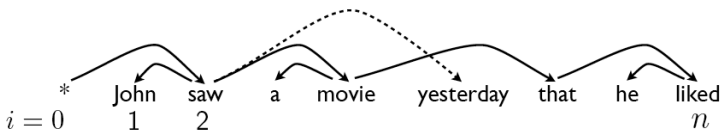
- Computer Vision:
 - Semantic Segmentation (output: pixel-wise labeling)
 - Object detection (output: 2D or 3D bounding boxes)
 - Stereo Reconstruction (output: 3D map)
 - Scene Understanding (output: 3D bounding box reprinting the layout)



Structured Prediction and its Applications

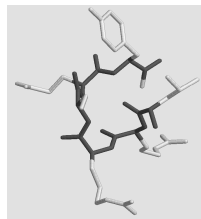
We want to predict multiple random variables which are related

- Natural Language processing
 - Machine Translation (output: sentence in another language)
 - Parsing (output: parse tree)



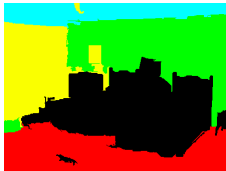
- Computational Biology
 - Protein Folding (output: 3D protein)

```
MRLILALLGICSLTAYIVEGVGSEVSDKR  
TCVSLTTQRLPVSRIKTYTITEGSLRAVIF  
ITKRGLKVCADPQATWVRDVVRSMDRKSNT  
RNNMIQTKPTGTQQSTNTAVTLTG
```



Why structured?

- Independent prediction is good but...

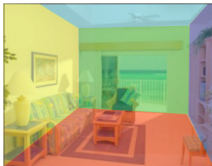


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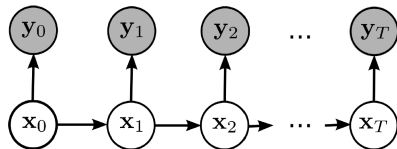
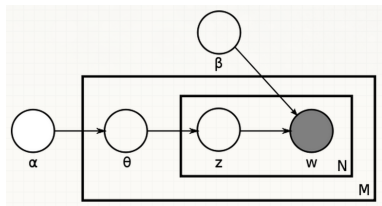
- Neighboring pixels should have same labels (if they look similar).



Graphical Model

A graphical model defines

- A family of probability distributions over a set of random variables
- This is expressed via a graph, which encodes the conditional independences of the distribution



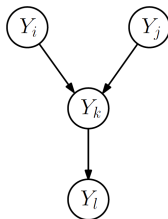
- Two types of graphical models: Directed and undirected

Bayesian Networks (Directed Graphical Models)

- The graph $G = (V, \mathcal{E})$ is acyclic and directed
- Factorization over distributions by conditioning on parent nodes

$$p(\mathbf{y}) = \prod_{i \in V} p(y_i | y_{pa(i)})$$

- Example



$$p(\mathbf{y}) = p(y_l | y_k) p(y_k | y_i, y_j) p(y_i) p(y_j)$$

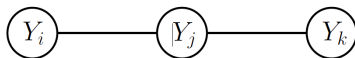
Undirected Graphical Model

- Also called Markov Random Field, or Markov Network
- Graph $G = (V, \mathcal{E})$ is undirected and has no self-edges
- Factorization over cliques

$$p(\mathbf{y}) = \frac{1}{Z} \prod_{r \in \mathcal{R}} \psi_r(\mathbf{y}_r)$$

with $Z = \sum_{\mathbf{y} \in \mathcal{Y}} \prod_{r \in \mathcal{R}} \psi_r(\mathbf{y}_r)$ the partition function

- Example



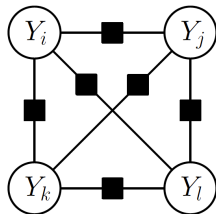
$$p(\mathbf{y}) = \frac{1}{Z} \psi(y_i, y_j) \psi(y_j, y_k) \psi(y_i) \psi(y_j) \psi(y_k)$$

- **Difficulty:** Exponentially many configurations
- Undirected models will be the focus of this lecture

Factor Graph Representation

- Graph $G = (V, \mathcal{F}, \mathcal{E})$, with variable nodes \mathcal{V} , factor nodes \mathcal{F} and edges \mathcal{E}
- **Scope** of a factor $N(F) = \{i \in V : (i, F) \in \mathcal{E}\}$
- Factorization over factors

$$p(\mathbf{y}) = \frac{1}{Z} \prod_{F \in \mathcal{F}} \psi_F(\mathbf{y}_{N(F)})$$



Factor Graph vs Graphical Model

- Factor graphs are explicit about the factorization

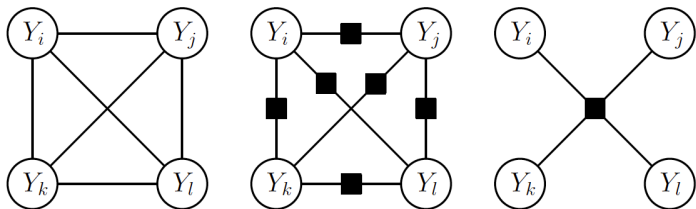


Figure : from [Nowozin et al]

- They define the family of distributions and thus the *capacity*

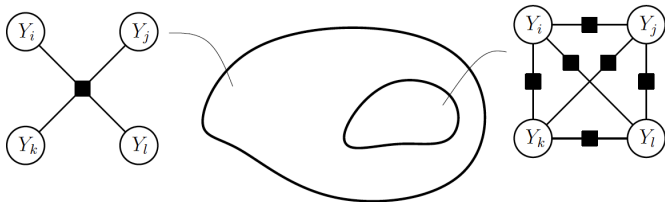


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Markov Random Fields vs Conditional Random Fields

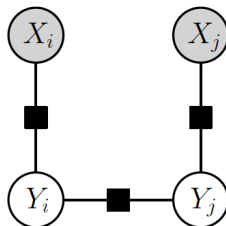
- Markov Random Fields (MRFs) define

$$p(\mathbf{y}) = \frac{1}{Z} \prod_{F \in \mathcal{F}} \psi_F(\mathbf{y}_{N(F)})$$

- Conditional Random Fields (CRFs) define

$$p(\mathbf{y}|\mathbf{x}) = \frac{1}{Z(\mathbf{x})} \prod_{F \in \mathcal{F}} \psi_F(\mathbf{y}_{N(F)}; \mathbf{x})$$

- \mathbf{x} is not a random variable (i.e., not part of the probability distribution)



- The probability is completely determined by the energy

$$\begin{aligned} p(\mathbf{y}) &= \frac{1}{Z} \prod_{F \in \mathcal{F}} \psi_F(\mathbf{y}_{N(F)}) \\ &= \frac{1}{Z} \exp(\log(\psi_F(\mathbf{y}_{N(F)}))) \\ &= \frac{1}{Z} \exp\left(-\sum_{F \in \mathcal{F}} E_F(y_F)\right) \end{aligned}$$

where $E_F(y_F) = -\log(\psi_F(\mathbf{y}_{N(F)}))$

Parameterization: log linear model

- Factor graphs define a family of distributions
- We are interested in identifying individual members by parameters

$$E_F(\mathbf{y}_F) = -\mathbf{w}^T \phi_F(\mathbf{y}_F)$$

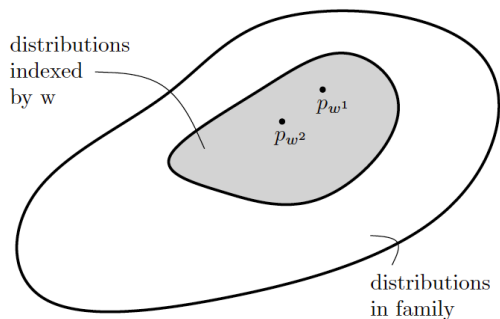


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- Estimation of the parameters \mathbf{w}

$$E_F(\mathbf{y}_F) = -\mathbf{w}^T \phi_F(\mathbf{y}_F)$$

- Learn the structure of the model
- Learn with hidden variables

Inference Tasks

Given an input $x \in \mathcal{X}$ we want to compute

- **MAP estimate** or minimum energy configuration

$$\begin{aligned}\operatorname{argmax}_{y \in \mathcal{Y}} p(\mathbf{y}|\mathbf{x}) &= \operatorname{argmax}_{y \in \mathcal{Y}} \frac{1}{Z} \prod_{F \in \mathcal{F}} \psi_F(\mathbf{y}_{N(F)}; \mathbf{x}, \mathbf{w}) \\ &= \operatorname{argmax}_{y \in \mathcal{Y}} \exp\left(-\sum_{F \in \mathcal{F}} E_F(\mathbf{y}_F, \mathbf{x}, \mathbf{w})\right) \\ &= \operatorname{argmin}_{y \in \mathcal{Y}} \sum_{F \in \mathcal{F}} E_F(\mathbf{y}_F, \mathbf{x}, \mathbf{w})\end{aligned}$$

Inference Tasks

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- **Marginals** $p(y_i)$ or max marginals $\max_{y_i \in \mathcal{Y}_i} p(y_i)$, which requires computing the partition function Z , i.e.,

$$\begin{aligned}\log(Z(\mathbf{x}, \mathbf{w})) &= \log \sum_{\mathbf{y} \in \mathcal{Y}} \exp(-E(\mathbf{y}; \mathbf{x}, \mathbf{w})) \\ \mu_F(\mathbf{y}_F) &= p(\mathbf{y}_F | \mathbf{x}, \mathbf{w})\end{aligned}$$

Inference in Markov Random Fields

Compute the MAP estimate is typically NP-hard

$$\max_{y \in \mathcal{Y}} p(\mathbf{y}|x) = \max_{y \in \mathcal{Y}} \sum_{r \in \mathcal{R}} \mathbf{w}^T \phi_r(\mathbf{y}_r)$$

MAP Inference

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Notable exceptions are:

- Belief propagation for tree-structure models

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- Graph cuts for binary energies with sub modular potentials

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Difficulties

- Deal with the exponentially many states in \mathbf{y}

Belief Propagation

- Compact notation

$$\theta_r(\mathbf{y}_r) = \mathbf{w}^T \phi_r(\mathbf{y}_r)$$

- Inference can be written as

$$\max_{\mathbf{y} \in \mathcal{Y}} \sum_{r \in \mathcal{R}} \theta_r(\mathbf{y}_r)$$



- For the example

$$\max_{y_i, y_j, y_k, y_l} \{ \theta_F(y_i, y_j) + \theta_G(y_j, y_k) + \theta_G(y_k, y_l) \}$$

Belief Propagation



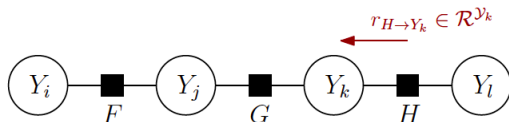
$$\theta^*(\mathbf{y}) = \max_{y_i, y_j, y_k, y_l} \{\theta_F(y_i, y_j) + \theta_G(y_j, y_k) + \theta_H(y_k, y_l)\}$$
$$=$$

Belief Propagation



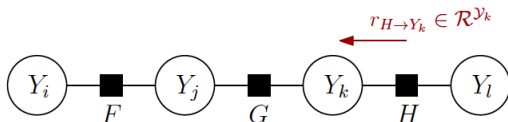
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Belief Propagation



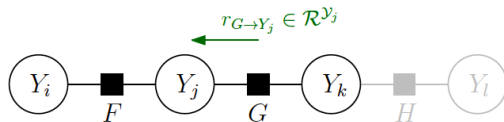
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Belief Propagation



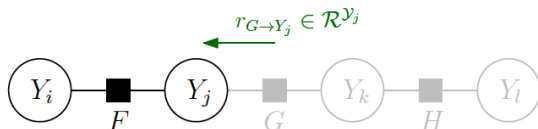
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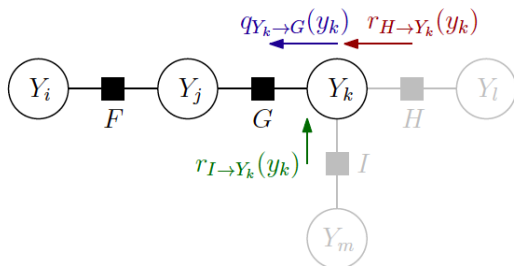
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Belief Propagation



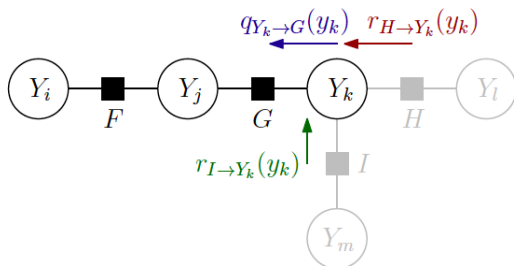
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Tree Generalization



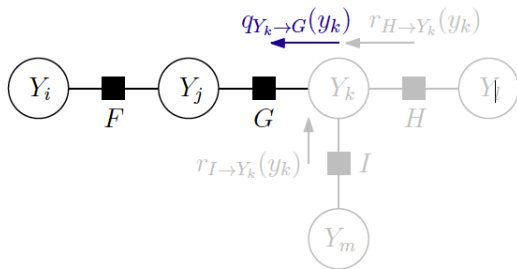
$$\theta^*(\mathbf{y}) = \max_{y_i, y_k, y_l, y_m} \theta_F(y_i, y_j) + \theta_G(y_j, y_k) + \theta_I(y_m, y_k) + \theta_H(y_l, y_k)$$

Tree Generalization



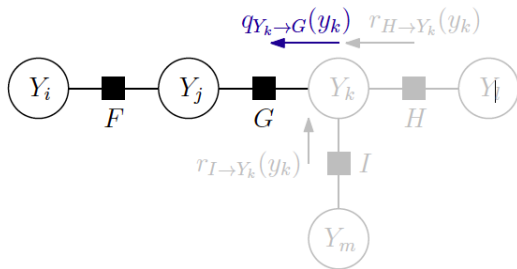
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 &= \max_{Y_i, Y_j} \theta_F(y_i, y_j) + \max_{y_k} \theta_G(y_j, y_k) + q_{y_k \rightarrow G}(y_k)
 \end{aligned}$$

Factor Graph Max Product

Iteratively updates and passes messages:

- $r_{F \rightarrow Y_i} \in \mathbb{R}^{\mathcal{Y}_i}$: factor to variable message
- $q_{Y_i \rightarrow F} \in \mathbb{R}^{\mathcal{Y}_i}$: variable to factor message

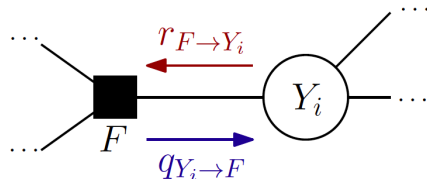


Figure : from [Nowozin et al]

Variable to factor

- Let $M(i)$ be the factors adjacent to variable i , $M(i) = \{F \in \mathcal{F} : (i, F) \in \mathcal{E}\}$
- Variable-to-factor message

$$q_{y_i \rightarrow F}(y_i) = \sum_{F' \in M(i) \setminus \{F\}} r_{F' \rightarrow y_i}(y_i)$$

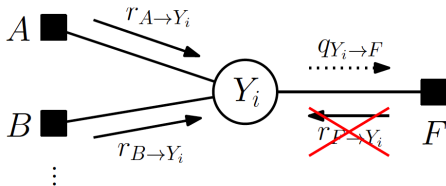


Figure : from [Nowozin et al]

Factor to variable

- Factor-to-variable message

$$r_{F \rightarrow y_i}(y_i) = \max_{y'_F \in \mathcal{Y}_F, y'_i = y_i} \left(\theta(y'_F) + \sum_{j \in \mathcal{N}(F) \setminus \{i\}} q_{y_j \rightarrow F}(y'_j) \right)$$

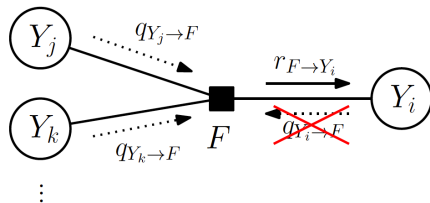


Figure : from [Nowozin et al]

Message Scheduling

- 1 Select one variable as tree root
- 2 Compute leaf-to-root messages
- 3 Compute root-to-leaf messages

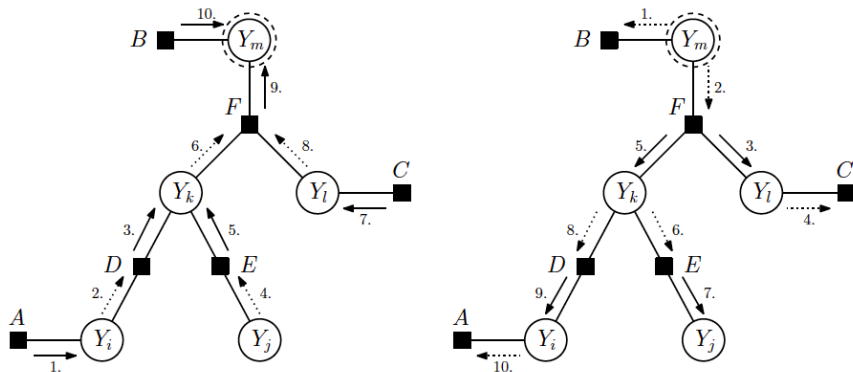


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Max Product v Sum Product

Max sum version of max-product

- 1 Compute leaf-to-root messages

$$q_{y_i \rightarrow F}(y_i) = \sum_{F' \in M(i) \setminus \{F\}} r_{F' \rightarrow y_i}(y_i)$$

- 2 Compute root-to-leaf messages

$$r_{F \rightarrow y_i}(y_i) = \max_{y'_F \in \mathcal{Y}_F, y'_i = y_i} \left(\theta(y'_F) + \sum_{j \in N(F) \setminus \{i\}} q_{y_j \rightarrow F}(y'_j) \right)$$

Max Product v Sum Product

Max sum version of max-product

- 1 Compute leaf-to-root messages

$$q_{y_i \rightarrow F}(y_i) = \sum_{F' \in M(i) \setminus \{F\}} r_{F' \rightarrow y_i}(y_i)$$

- 2 Compute root-to-leaf messages

$$r_{F \rightarrow y_i}(y_i) = \max_{y'_F \in \mathcal{Y}_F, y'_i = y_i} \left(\theta(y'_F) + \sum_{j \in N(F) \setminus \{i\}} q_{y'_j \rightarrow F}(y'_j) \right)$$

Sum-product

- 1 Compute leaf-to-root messages

$$q_{y_i \rightarrow F}(y_i) = \sum_{F' \in M(i) \setminus \{F\}} r_{F' \rightarrow y_i}(y_i)$$

- 2 Compute root-to-leaf messages

$$r_{F \rightarrow y_i}(y_i) = \log \sum_{y'_F \in \mathcal{Y}_F, y'_i = y_i} \exp \left(\theta(y'_F) + \sum_{j \in N(F) \setminus \{i\}} q_{y'_j \rightarrow F}(y'_j) \right)$$

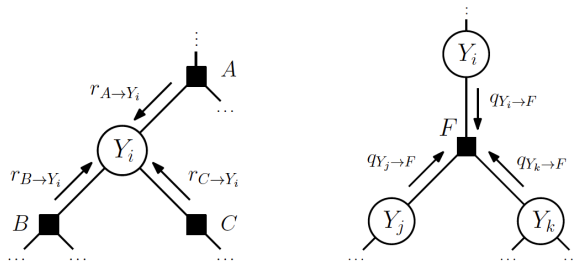
Computing marginals

- Partition function can be evaluated at the root

$$\log Z = \log \sum_{y_r} \exp \left(\sum_{F \in M(r)} r_{F \rightarrow y_r}(y_r) \right)$$

- Marginal distributions, for each factor

$$\mu_F(y_F) = p(y_F) = \frac{1}{Z} \exp \left(\theta_F(y_F) + \sum_{i \in N(F)} q_{y_i \rightarrow F}(y_i) \right)$$



Computing marginals

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- Marginals at every node

$$\mu_{y_i}(y_i) = p(y_i) = \frac{1}{Z} \exp \left(\sum_{F \in M(i)} r_{F \rightarrow y_i}(y_i) \right)$$

Generalizations to loops

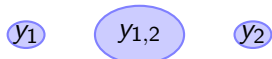
- It is call **loopy belief propagation** (Perl, 1988)
- No schedule that removes dependencies
- Different messaging schedules (synchronous/asynchronous, static/dynamic)
- Slight changes in the algorithm

MAP LP Relaxation Task

Integer Linear Program (LP) equivalence [Werner 2007]:

- Inference task:

$$\hat{\mathbf{y}} = \arg \max_{\mathbf{y}} \sum_r \theta_r(\mathbf{y}_r)$$



- Variables $b_r(\mathbf{y}_r)$:

$$\max_{b_1, b_2, b_{12}} \begin{bmatrix} b_1(0) \\ b_1(1) \\ b_2(0) \\ b_2(1) \\ b_{12}(0,0) \\ b_{12}(1,0) \\ b_{12}(0,1) \\ b_{12}(1,1) \end{bmatrix}^T \begin{bmatrix} \theta_1(0) \\ \theta_1(1) \\ \theta_2(0) \\ \theta_2(1) \\ \theta_{12}(0,0) \\ \theta_{12}(1,0) \\ \theta_{12}(0,1) \\ \theta_{12}(1,1) \end{bmatrix}$$

MAP LP Relaxation Task

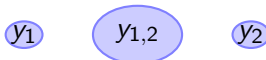
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MAP LP Relaxation Task

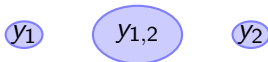
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MAP LP Relaxation Task

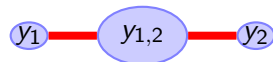
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$$\text{s.t.} \quad \begin{aligned} b_r(\mathbf{y}_r) &\in \{0, 1\} \\ \sum_{\mathbf{y}_r} b_r(\mathbf{y}_r) &= 1 \\ \sum_{\mathbf{y}_p \setminus \mathbf{y}_r} b_p(\mathbf{y}_p) &= b_r(\mathbf{y}_r) \end{aligned}$$

MAP LP Relaxation Task

$$\max_{b_1, b_2, b_{12}} \begin{bmatrix} b_1(1) \\ b_1(2) \\ b_2(1) \\ b_2(2) \\ b_{12}(1, 1) \\ b_{12}(2, 1) \\ b_{12}(1, 2) \\ b_{12}(2, 2) \end{bmatrix}^T \begin{bmatrix} \theta_1(1) \\ \theta_1(2) \\ \theta_2(1) \\ \theta_2(2) \\ \theta_{12}(1, 1) \\ \theta_{12}(2, 1) \\ \theta_{12}(1, 2) \\ \theta_{12}(2, 2) \end{bmatrix}$$

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MAP LP Relaxation Task

$$\begin{aligned} \max_{b_r} \quad & \sum_{r, \mathbf{y}_r} b_r(\mathbf{y}_r) \theta_r(\mathbf{y}_r) \\ \text{s.t.} \quad & b_r(\mathbf{y}_r) \in \{0, 1\} \\ & \sum_{\mathbf{y}_r} b_r(\mathbf{y}_r) = 1 \\ & \sum_{\mathbf{y}_p \setminus \mathbf{y}_r} b_p(\mathbf{y}_p) = b_r(\mathbf{y}_r) \end{aligned}$$

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MAP LP Relaxation Task

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MAP LP Relaxation Task

LP relaxation:

\max_{b_r}

$$\sum_{r, \mathbf{y}_r} b_r(\mathbf{y}_r) \theta_r(\mathbf{y}_r)$$

s.t.

~~$$b_r(\mathbf{y}_r) \in \{0, 1\}$$~~

Local probability b_r

Marginalization

MAP LP Relaxation Task

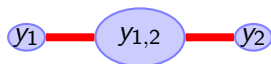
LP relaxation:

$$\begin{aligned} \max_{b_r} \quad & \sum_{r, \mathbf{y}_r} b_r(\mathbf{y}_r) \theta_r(\mathbf{y}_r) \\ \text{s.t.} \quad & \cancel{b_r(\mathbf{y}_r) \in \{0, 1\}} \\ & \text{Local probability } b_r \\ & \text{Marginalization} \end{aligned}$$

Can be solved by any standard LP solver but **slow** because of typically many variables and constraints. Can we do better?

MAP LP Relaxation Task

Observation: Graph structure in marginalization constraints.



Use dual to take advantage of structure in constraint set

- Set of parents of region r : $P(r)$
- Set of children of region r : $C(r)$

$$\forall r, \mathbf{y}_r, p \in P(r) \quad \sum_{\mathbf{y}_p \setminus \mathbf{y}_r} b_p(\mathbf{y}_p) = b_r(\mathbf{y}_r)$$

- Lagrange multipliers for every constraint:

$$\forall r, \mathbf{y}_r, p \in P(r) \quad \lambda_{r \rightarrow p}(\mathbf{y}_r)$$

MAP LP Relaxation Task

Re-parameterization of score $\theta_r(\mathbf{y}_r)$:

$$\hat{\theta}_r(\mathbf{y}_r) = \theta_r(\mathbf{y}_r) + \sum_{p \in P(r)} \lambda_{r \rightarrow p}(\mathbf{y}_r) - \sum_{c \in C(r)} \lambda_{c \rightarrow r}(\mathbf{y}_c)$$

Properties of dual program:

$$\min_{\lambda} q(\lambda) = \min_{\lambda} \sum_r \max_{\mathbf{y}_r} \hat{\theta}_r(\mathbf{y}_r)$$

- **Dual upper-bounds primal** $\forall \lambda$
- Convex problem
- Unconstrained task
- Doing block coordinate descent in the dual results on message passing (Lagrange multipliers are your messages)

Block-coordinate descent solvers iterate the following steps:

- Take a block of Lagrange multipliers
- Optimize sub-problem of dual function w.r.t. this block while keeping all other variables fixed

Advantage: fast due to analytically computable sub-problems

Same type of algorithms also exist to compute approximate marginals

Theorem [Kolmogorov and Zabih, 2004]: If the energy function is a function of binary variables containing only unary and pairwise factors, the discrete energy minimization problem

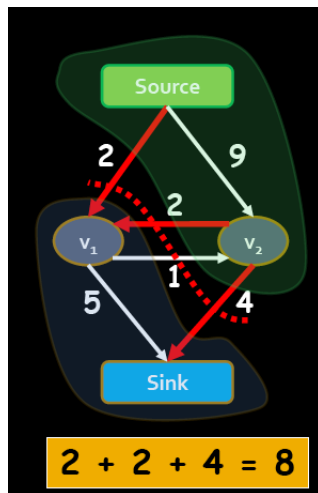
$$\min_{\mathbf{y}} \sum_{r \in \mathcal{R}} E(\mathbf{y}_r, x)$$

can be formulated as a graph cut problem if and only if all pairwise energies are sub modular

$$E_{i,j}(0,0) + E_{i,j}(1,1) \leq E_{i,j}(0,1) + E_{i,j}(1,0)$$

The ST-mincut problem

- The st-mincut is the st-cut with the minimum cost

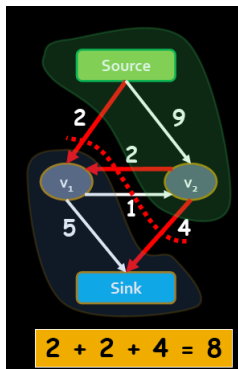


[Source: P. Kohli]

Back to our energy minimization

Construct a graph such that

- 1 Any st-cut corresponds to an assignment of x
- 2 The cost of the cut is equal to the energy of x : $E(x)$



[Source: P. Kohli]

St-mincut and Energy Minimization

$$E(\mathbf{x}) = \sum_i \theta_i(x_i) + \sum_{i,j} \theta_{ij}(x_i, x_j)$$

For all ij $\theta_{ij}(0,1) + \theta_{ij}(1,0) \geq \theta_{ij}(0,0) + \theta_{ij}(1,1)$



Equivalent (transformable)

$$E(\mathbf{x}) = \sum_i c_i x_i + \sum_{i,j} c_{ij} x_i(1-x_j) \quad c_{ij} \geq 0$$

[Source: P. Kohli]

How are they equivalent?

$A = \theta_{ij}(0,0)$

$B = \theta_{ij}(0,1)$

$C = \theta_{ij}(1,0)$

$D = \theta_{ij}(1,1)$

		x_j		
		0	1	
x_i	0	A	B	= A +
	1	C	D	

		x_j		
		0	1	
x_i	0	0	0	+ A +
	1	C-A	C-A	

if $x_1=1$ add C-
A

		x_j		
		0	1	
x_i	0	0	D-C	+ A +
	1	0	D-C	

if $x_2 = 1$ add
D-C

		x_j		
		0	1	
x_i	0	0	B +C- A-D	+ A +
	1	0	0	

$$\theta_{ij}(x_i, x_j) = \theta_{ij}(0,0)$$

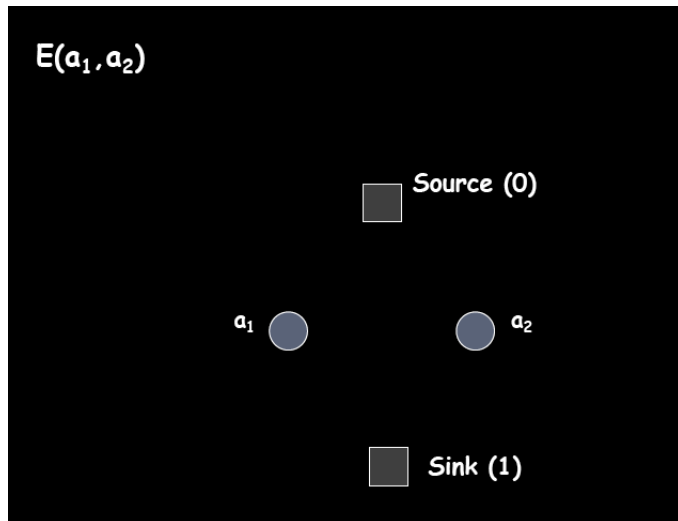
$$+ (\theta_{ij}(1,0) - \theta_{ij}(0,0)) x_i + (\theta_{ij}(0,1) - \theta_{ij}(0,0)) x_j$$

$$+ (\theta_{ij}(1,0) + \theta_{ij}(0,1) - \theta_{ij}(0,0) - \theta_{ij}(1,1)) (1-x_i) x_j$$

$B+C-A-D \geq 0$ is true from the submodularity of θ_{ij}

[Source: P. Kohli]

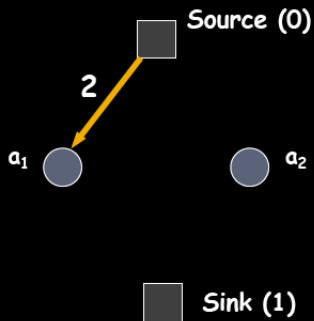
Graph Construction



[Source: P. Kohli]

Graph Construction

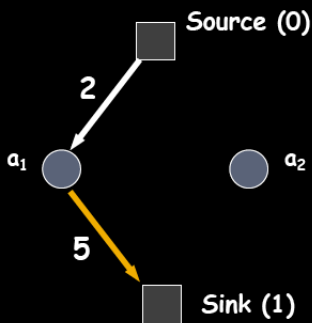
$$E(\mathbf{a}_1, \mathbf{a}_2) = 2\mathbf{a}_1$$



[Source: P. Kohli]

Graph Construction

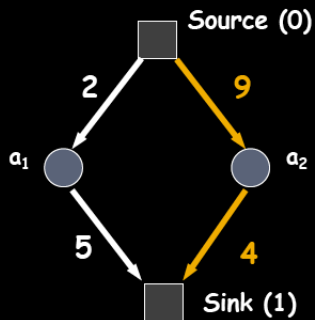
$$E(a_1, a_2) = 2a_1 + 5\bar{a}_1$$



[Source: P. Kohli]

Graph Construction

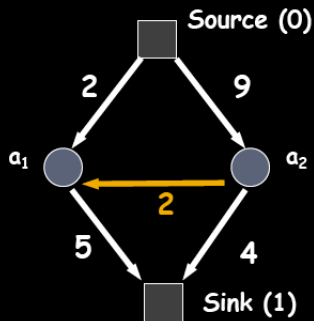
$$E(a_1, a_2) = 2a_1 + 5\bar{a}_1 + 9a_2 + 4\bar{a}_2$$



[Source: P. Kohli]

Graph Construction

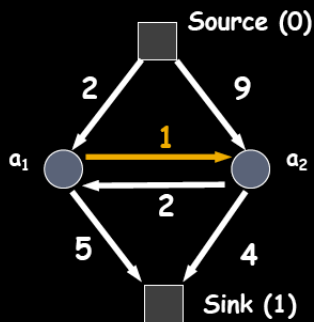
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[Source: P. Kohli]

Graph Construction

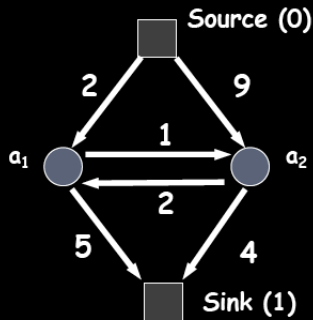
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[Source: P. Kohli]

Graph Construction

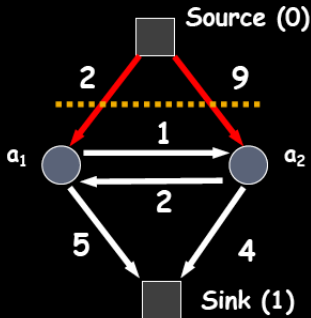
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[Source: P. Kohli]

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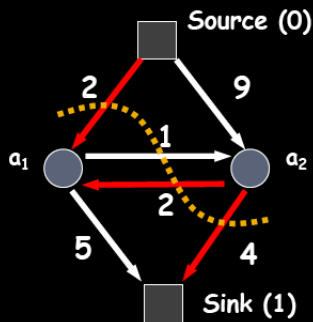
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[Source: P. Kohli]

Graph Construction

$$E(a_1, a_2) = 2a_1 + 5\bar{a}_1 + 9a_2 + 4\bar{a}_2 + 2a_1\bar{a}_2 + \bar{a}_1a_2$$



st-minicut cost = 8

$$a_1 = 1 \quad a_2 = 0$$

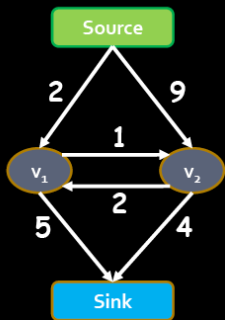
$$E(1, 0) = 8$$

[Source: P. Kohli]

How to compute the St-mincut?

Solve the dual **maximum flow** problem

Compute the maximum flow between Source and Sink s.t.



Edges: Flow < Capacity

Nodes: Flow in = Flow out

Min-cut \ Max-flow Theorem

In every network, the maximum flow equals the cost of the st-mincut

Assuming non-negative capacity

[Source: P. Kohli]

How does the code look like

```
Graph *g;
```

For all pixels p

```
/* Add a node to the graph */
```

```
nodeID(p) = g->add_node();
```

```
/* Set cost of terminal edges */
```

```
set_weights(nodeID(p), fgCost(p), bgCost(p));
```

```
end
```

for all adjacent pixels p,q

```
add_weights(nodeID(p), nodeID(q), cost(p,q));
```

```
end
```

```
g->compute_maxflow();
```

```
label_p = g->is_connected_to_source(nodeID(p));
```

```
// is the label of pixel p (0 or 1)
```

 **Source (0)**

 **Sink (1)**

[Source: P. Kohli]

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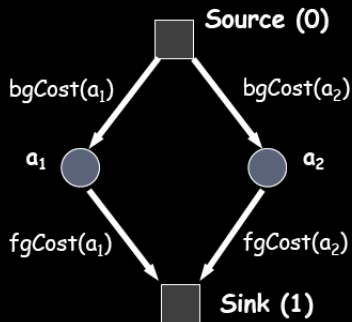
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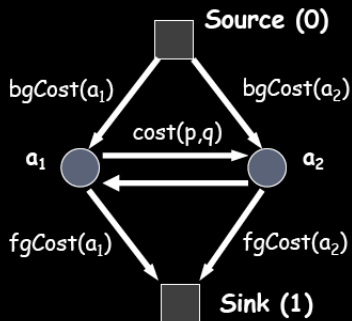
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[Source: P. Kohli]

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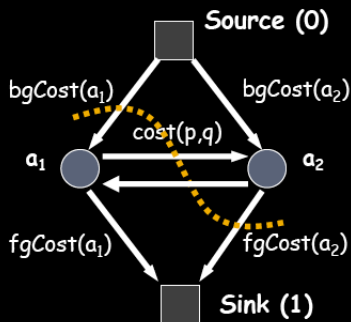
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```
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```



$a_1 = bg$ $a_2 = fg$

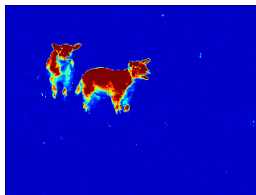
[Source: P. Kohli]

Example: Figure-Ground Segmentation

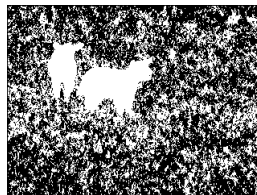
Binary labeling problem



(Original)



(Color model)



(Indep. Prediction)

Figure : from [Nowozin et al]

Example: Figure-Ground Segmentation

- Markov Random Field

$$E(\mathbf{y}, \mathbf{x}, \mathbf{w}) = \sum_i \log p(y_i | x_i) + w \sum_{(i,j) \in \mathcal{E}} C(x_i, x_j) I(y_i \neq y_j)$$

with $C(x_i, x_j) = \exp(\gamma \|x_i - x_j\|^2)$, and $w \geq 0$.

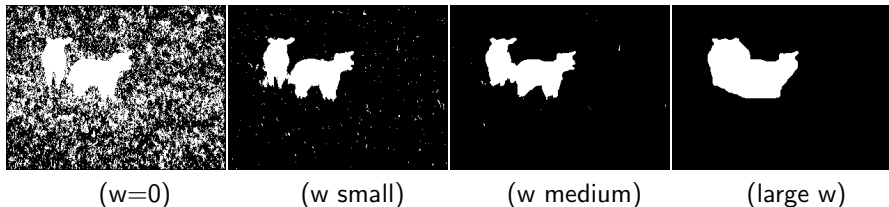


Figure : from [Nowozin et al]

- Why do we need the condition $w \geq 0$?

Generalization to Multi-label Problems

- Optimal solution is not possible anymore
- Solve to optimality subproblems that include current iterate
- This guarantees decrease in the objective

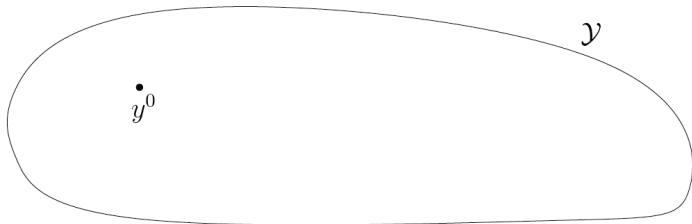


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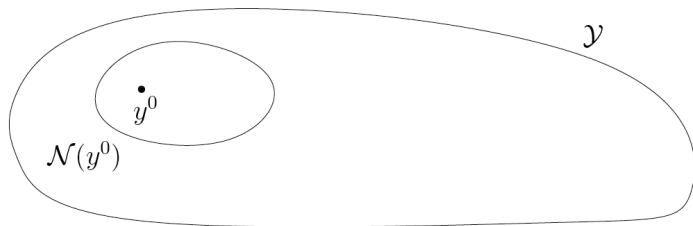


Figure : from [Nowozin et al]

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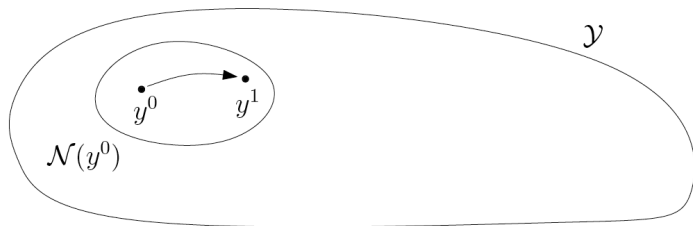


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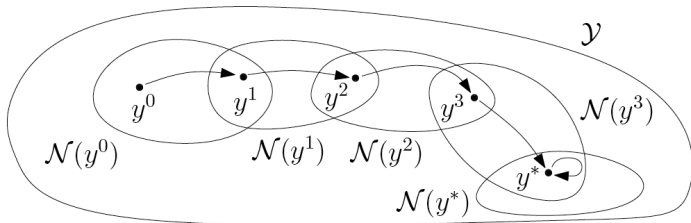


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Two general classes of pairwise interactions

- **Metric** if it satisfies for any set of labels α, β, γ

$$V(\alpha, \beta) = 0 \quad \Leftrightarrow \quad \alpha = \beta$$

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Examples for 1D label set

- Truncated quadratic is a semi-metric

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- Potts model is a metric

$$V(\alpha, \beta) = K \cdot T(\alpha \neq \beta)$$

with $T(\cdot) = 1$ if the argument is true and 0 otherwise.

Move Making Algorithms

- **Alpha Expansion:** Checks if current nodes want to switch to label α
- **Alpha - Beta Swaps:** Checks if a node with class α wants to switch to β .
- Binary problems that can be solve exactly for certain type of potentials

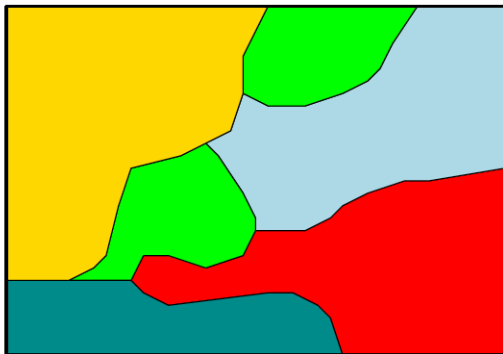


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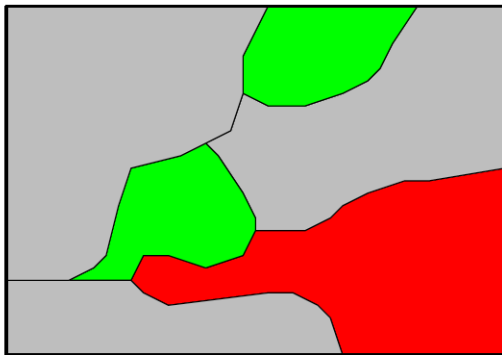


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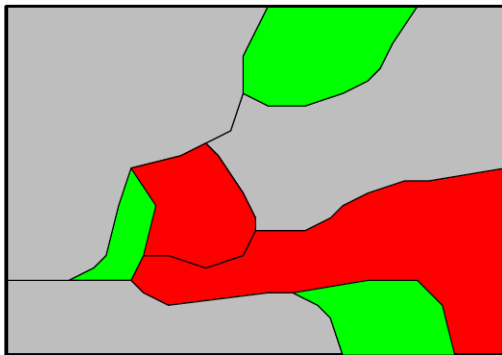


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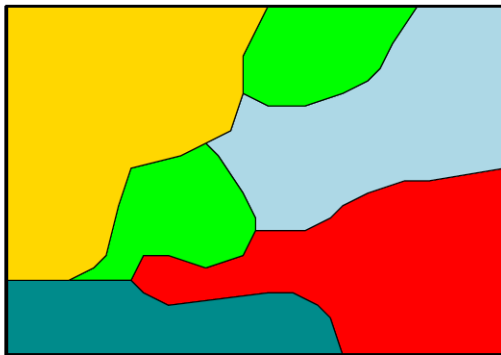


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Binary Moves

- $\alpha - \beta$ moves works for semi-metrics
- α expansion works for V being a metric

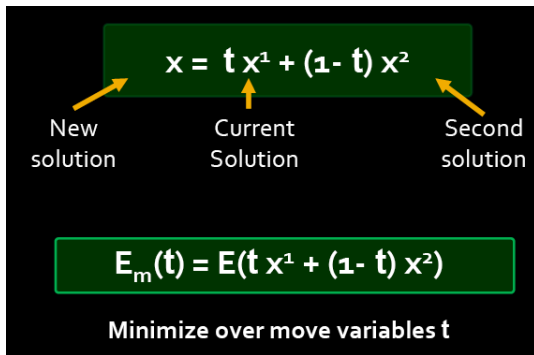
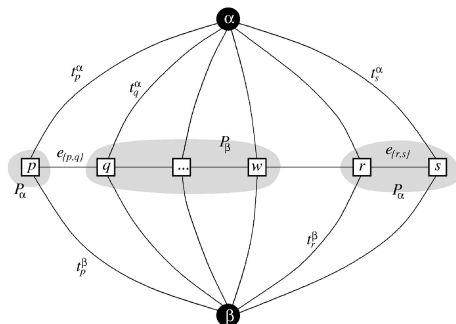


Figure : from P. Kohli tutorial on graph-cuts

- For certain x^1 and x^2 , the move energy is sub-modular

Graph Construction

- The set of vertices includes the two terminals α and β , as well as image pixels p in the sets \mathcal{P}_α and \mathcal{P}_β (i.e., $f_p \in \{\alpha, \beta\}$).
- Each pixel $p \in \mathcal{P}_{\alpha\beta}$ is connected to the terminals α and β , called t -links.
- Each set of pixels $p, q \in \mathcal{P}_{\alpha\beta}$ which are neighbors is connected by an edge $e_{p,q}$



edge	weight	for
t_p^α	$D_p(\alpha) + \sum_{\substack{q \in \mathcal{N}_p \\ q \in \mathcal{P}_{\alpha\beta}}} V(\alpha, f_q)$	$p \in \mathcal{P}_{\alpha\beta}$
t_p^β	$D_p(\beta) + \sum_{\substack{q \in \mathcal{N}_p \\ q \in \mathcal{P}_{\alpha\beta}}} V(\beta, f_q)$	$p \in \mathcal{P}_{\alpha\beta}$
$e_{\{p,q\}}$	$V(\alpha, \beta)$	$\{p,q\} \in \mathcal{N}$ $p, q \in \mathcal{P}_{\alpha\beta}$