

Lambda Calculus

$(\lambda x. M)$: define a function for x that returns M .

$$f(x) = M \Rightarrow (\lambda x. M)(x) = M.$$

β reduction:

(in contrast to

α -conversion)

$$((\lambda x. M) N) \rightarrow (M[x := N])$$

e.g. $\lambda x. f(x)$

$\lambda F. \forall y. F(y)$

$$(\lambda F. \forall y. F(y))(\lambda x. f(x)) = ?$$

$$(\lambda x. M) \quad (N)$$
$$(\lambda F. \forall y. F(y))(\lambda x. f(x))$$

$$\Leftrightarrow_{\beta} \forall y. (\lambda x. f(x))(y).$$

$$\Leftrightarrow_{\beta} \forall y. f(y).$$

Q2: Macros: $@\text{lambda}(X, F) \quad \lambda X. F$

$@\text{forall}(X, F, G) \quad \forall X. F \Rightarrow G$

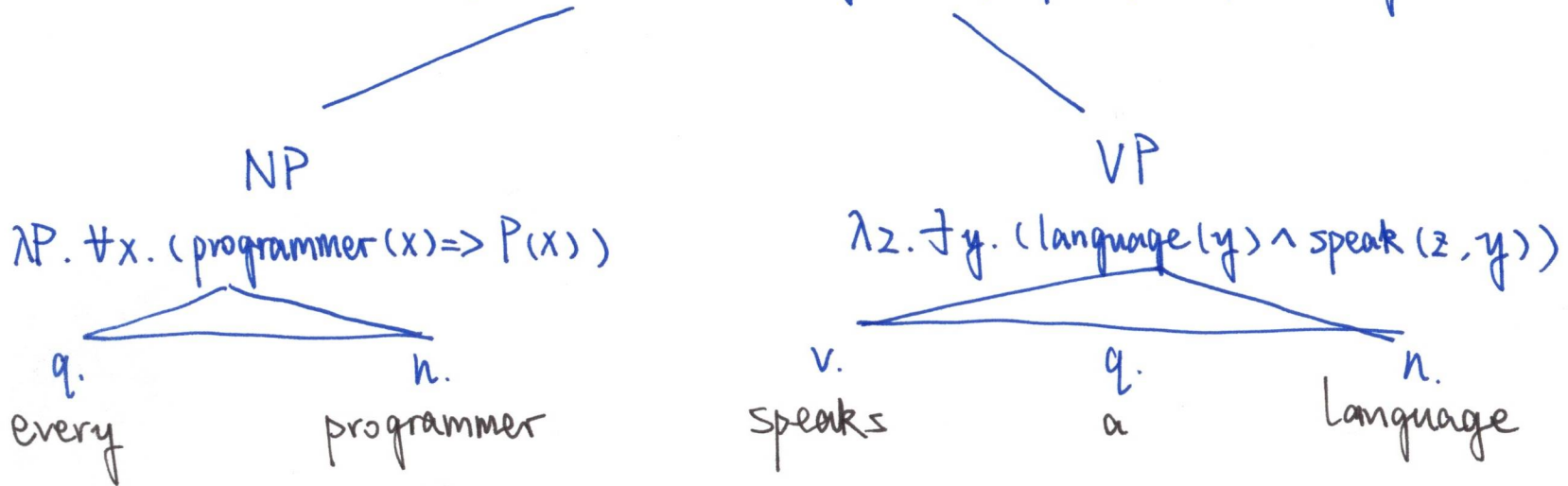
$@\text{exists}(X, F, G) \quad \exists X. F \wedge G$

$@\text{apply}(F, [X, Y, Z]) \quad F(X, Y, Z).$

beta-normalize: $\text{beta-normalize}(@\text{apply}(F, [G]), B).$

Example: Every programmer speaks a language.

$$S \quad \forall x. (\text{programmer}(x) \Rightarrow \exists y. (\text{language}(y) \wedge \text{speaks}(x, y)))$$



$$NP. \forall x. (\text{programmer}(x) \Rightarrow P(x))$$

$$VP. \lambda z. \exists y. (\text{language}(y) \wedge \text{speaks}(z, y))$$

$$\lambda x. \underbrace{(\text{programmer}(x) \Rightarrow P(x))}_{M_1} \underbrace{(\lambda z. \exists y. (\text{language}(y) \wedge \text{speaks}(z, y)))}_{N_1}$$

$$\Leftrightarrow \beta \quad \forall x. (\text{programmer}(x) \Rightarrow \underbrace{\lambda z. \exists y. (\text{language}(y) \wedge \text{speaks}(z, y))}_{M_2} \underbrace{(x)}_{N_2})$$

$$\Leftrightarrow \beta \quad \forall x. (\text{programmer}(x) \Rightarrow \exists y. (\text{language}(y) \wedge \text{speaks}(x, y)))$$

Quantifier Storage

$$\text{Storage: } \lambda h. \exists y. (\text{Language}(y) \wedge h(y))$$

$$\Downarrow$$
$$\lambda F. F(z), \text{QSTORE} \langle z; \lambda h. \exists y. (\text{Language}(y) \wedge h(y)) \rangle$$

$$\text{Retrieval: } \exists x. \text{programmer}(x) \Rightarrow \text{speaks}(x, z), \text{QSTORE} \langle z; \lambda h. \exists y. (\text{Language}(y) \wedge h(y)) \rangle$$

$$\Downarrow$$
$$\lambda z. \forall x. \text{programmer}(x) \Rightarrow \text{speaks}(x, z), \text{QSTORE} \langle \dots \rangle$$

$$\Downarrow$$
$$\underbrace{\lambda x}_{M_1} \underbrace{(\lambda h. \exists y. (\text{Language}(y) \wedge h(y)))}_{N_1} \underbrace{(\lambda z. \forall x. \text{programmer}(x) \Rightarrow \text{speaks}(x, z))}_{N_2}$$

$$\Leftrightarrow \beta \exists y. (\text{Language}(y) \wedge \underbrace{\lambda z. \forall x. \text{programmer}(x) \Rightarrow \text{speaks}(x, z)}_{M_2} \underbrace{(y)}_{N_2})$$

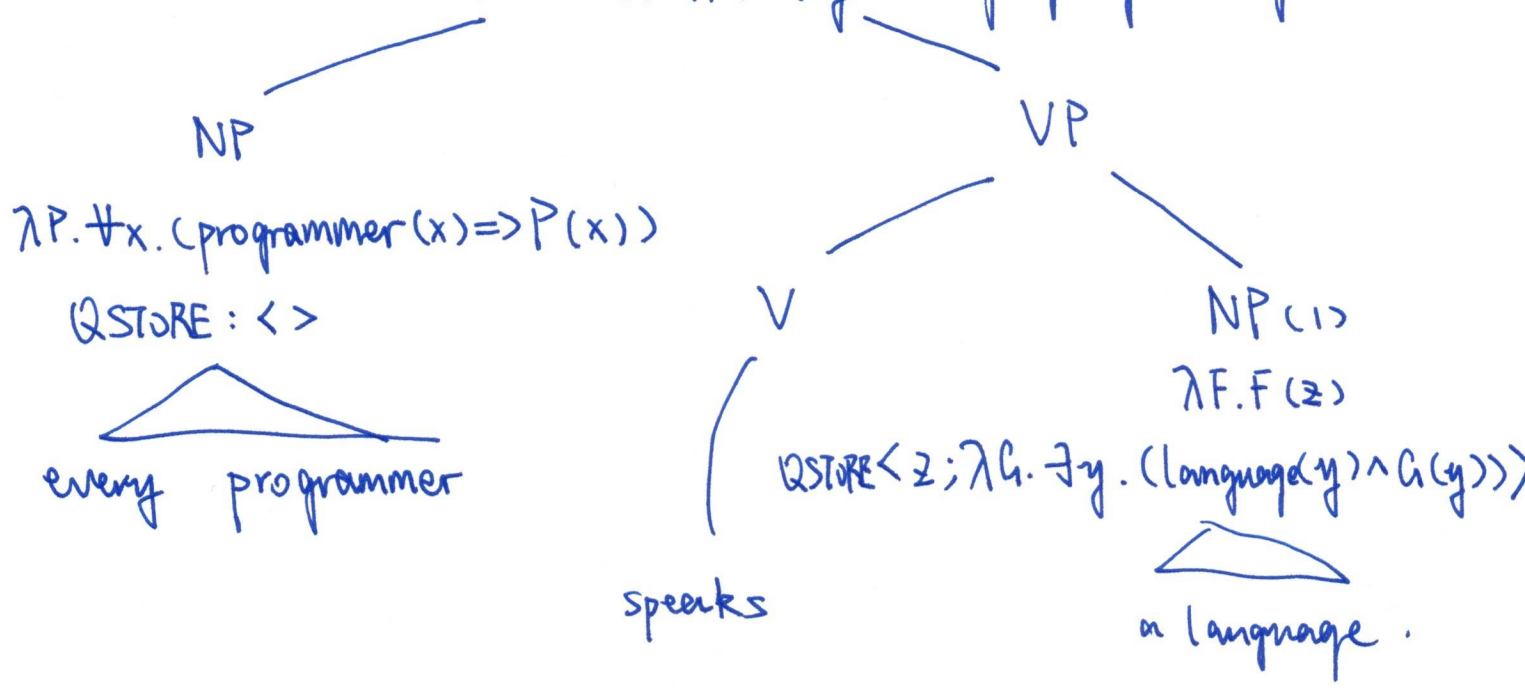
$$\Leftrightarrow \beta \exists y. (\text{Language}(y) \wedge \forall x. \text{programmer}(x) \Rightarrow \text{speaks}(x, y))$$

Quantifier Storage

S (2)

$\forall x. \text{programmer}(x) \Rightarrow \text{speaks}(x, z)$

QSTORE: $\langle z; \lambda G. \exists y. (\text{language}(y) \wedge G(y)) \rangle$



STORE at (1) , Retrieve at (2)

action (Logic , QStore , NewLogic , NewQStore) :

↳ start with Logic and QSTORE,
comes out with NewLogic and NewQStore

retrieve (QStore , Logic , NewQStore , NewLogic)

↳ start with QStore and Logic
comes out as NewQStore , what's out is added To Logic
→ NewLogic