

Lossless Join Decomposition - 3NF Solutions

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1. Are these schemas in 3NF?
2. Decompose the relations, as necessary, into collections of relations that are in 3NF.

(a) $R = \{city, street, zip\}$
 $\mathcal{F} = \{city, street \rightarrow zip, zip \rightarrow city\}$

$\{city, street\}$ and $\{zip, street\}$ are the only keys. The left hand side of $city, street \rightarrow zip$ is a key, so it's ok. The right hand side of $zip \rightarrow city$ is part of a key, so it's ok. Therefore the schema is in 3NF.

(b) $R = ABC$
 $\mathcal{F} = \{A \rightarrow B, B \rightarrow C\}$

The only key is A . The lhs of $B \rightarrow C$ is not a superkey, the rhs is not part of a key. Therefore the schema is not in 3NF.

We decompose R into $R_1 = AB$ and $R_2 = BC$.

(c) $R = ABCD$
 $\mathcal{F} = \{AB \rightarrow C, C \rightarrow D, D \rightarrow A\}$

The keys are AB , BC and BD . The lhs of $AB \rightarrow C$ is a key, the rhs of $C \rightarrow D$ is part of a key, the rhs of $D \rightarrow A$ is part of a key. Therefore the schema is in 3NF.

(d) $R = ABCD$
 $\mathcal{F} = \{B \rightarrow C, B \rightarrow D\}$

The only key is AB . The lhs of $B \rightarrow C$ is not a superkey, the rhs is not part of a key. Therefore the schema is not in 3NF.

We decompose R into $R_1 = BC$ and $R_2 = BD$. The key AB isn't contained in any schema so we add $R_3 = AB$. We can replace R_1 and R_2 by $R_4 = BCD$. The decomposition is then $\rho = \{BCD, AB\}$.

- (e) $R = ABCD$
 $\mathcal{F} = \{AB \rightarrow C, BC \rightarrow D, CD \rightarrow A, AD \rightarrow B\}$

The only keys are $\{AB, BC, CD, AD\}$. All the functional dependencies have their lhs as keys. Therefore the schema is in 3NF.