CSC2412: Definition of Differential Privacy

Sasho Nikolov

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 - In particular, believes Sasho has four fingers on each hand.

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Another example:

- Adversary believes there is no link between smoking and cancer.
 - Also knows that Sasho smokes
- Study reveals link between smoking and cancer.

earning Hout the world also means learning about we In the examples, the adversary learns statistical information that pertains to Sasho.

• If science works, it better reveal something about me.

What information is *statistical* and what information is *personal*?

Test: Could the adversary have learned this information if my data were not analyzed?

The algorithm doing the analysis should do almost the same in all the following cases:

- my data is **included** in the data set
- my data is not included in the data set
- my data is **changed** in the data set

Data set: (multi-)set X of n data points $X = \{x_1, \ldots, x_n\}$.

• each data point (or row) x_i is the data of one person χ =

• each data point comes from a *universe*
$$\lambda$$

E.g. d binary attributes
 $x_i \in \mathcal{X} = \{0, 1\}$



A data analysis algorithm (a *mechanism*) is a **randomized** algorithm \mathcal{M} that takes a data set X and produces the results of the data analysis as output.

The output of
$$M(X)$$
 is random
for any X

single individual X=fx,...,xn} We call two data sets X and X' (neighbouring) if 1. (*variable* \mathfrak{D} we can get X' from X by adding or removing an element $X' = \{x_{i-1}, x_{i+1}, x_{i+1},$ 2. (fixed n) we can get X' from X by replacing an element with another $X = \{x_{1}, \dots, x_{n}\}$ X'=hx, x:, x:, X:, X:+...., x, S

Definition

An mechanism \mathcal{M} is differentially private if, for any two neighbouring datasets X, X'

M(X) (X) (X) and
$$\mathcal{U}(X')$$
 are
"similar" as random variables

Total Variation Distance Differential Privacy

 $d_{tv}(\mathcal{M}(\mathcal{K}), \mathcal{M}(\mathcal{K}')) = \max | \mathbb{P}(\mathcal{M}(\mathcal{K}) \in S) - \mathbb{P}(\mathcal{M}(\mathcal{K}') \in S) |$ Definition An mechanism \mathcal{M} is δ -tv differentially private if, for any two neighbouring datasets X, X', and any set of outputs S $X = \{ \bigotimes_{j=1}^{n}, \bigotimes_{j=1}^{n} \}$ $X' = \{ x_{1}, \dots, x_{n} \}$ $|\mathbb{P}(\mathcal{M}(X) \in S) - \mathbb{P}(\mathcal{M}(X') \in S)| \le \delta.$ What should δ be? Not nec. veighbouring What should δ be? For any X, X', there are $k \leq n$ data sets $\delta < \frac{1}{2n}$? reighbouring $TO X^{(1)} X^{(1)} X^{(2)}$, $f^{(k-1)} X X^{(k)} = X'$ max $\left[P(\mathcal{M}(X) \in S) - P(\mathcal{M}(X^{1}) \in S) \right] \leq \delta n < \frac{1}{2}$ all datasets • $\delta \ge \frac{1}{2n}$? "Name and shame" mechanism: For all i, output x_i of prob. of $\forall X \sim X', X = \{x_1, \dots, x_n\}$ of $\forall prob \cdot \delta \times i$ is not published and M(K) = M(K')neighbouring $K = \{x_1, \dots, x_n\}$) not intuitively private is some data pt published of const prob

Finally, Differential Privacy

Definition

An mechanism \mathcal{M} is ε -differentially private if, for any two neighbouring datasets X, X', and any set of outputs S

$$\mathbb{P}(\mathcal{M}(X) \in S) \leq \mathbb{E}^{\mathbb{P}}(\mathcal{M}(X') \in S).$$

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$$\mathbb{P}(\mathbb{P}(X') \in S)$$

$$\mathbb{P}(\mathbb{P}$$

A Hypothesis Testing Viewpoint

Suppose
$$X = \{X_1, \dots, X_n\}$$
 are drawn (IID) from some distribution.
The adversary A wants to use $\mathcal{M}(X)$ to test which hypothesis holds:
 $H_0: X_i = y_0$
• E.g., "Sasho does not smoke"
 $H_1: X_i = y_1$
• E.g., "Sasho smokes"
Then for any A (that soes $\mathcal{M}(K)$ and outputs "Ho", "H,")
 $\mathbb{P}(\mathcal{A}(\mathcal{M}(X)) = "H_1" \mid H_1) \leq e^{\varepsilon} \mathbb{P}(\mathcal{A}(\mathcal{M}(X)) = "H_1" \mid H_0)$
True Positive rate False fositive rate
I- Type I error Type I error

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Warner

Given

• dataset
$$X = \{x_1, \ldots, x_n\} \subseteq \mathcal{X}$$
,
• query $q : \mathcal{X} \to \{0, 1\}$ E.g. $q(x) = \begin{cases} i & \text{if } x \text{ is a surder} \\ 0 & 0 \end{cases}$

output $\mathcal{M}(X) = (Y_1(x_1), \dots, Y_n(x_n))$, where, independently

$$Y_i(x_i) = \begin{cases} q(x_i) & \text{w/ prob. } \frac{e^{\varepsilon}}{1+e^{\varepsilon}} > \frac{1}{2} \\ 1-q(x_i) & \text{w/ prob. } \frac{1}{1+e^{\varepsilon}} < \frac{1}{2} \end{cases}$$

Privacy Analysis

$$\text{ETS for any } y \in \{0,1\}^n, \text{ and any neighbouring } X, X' \qquad Y_i(x_i) = \begin{cases} q(x_i) & \text{w/ prob. } \frac{e^e}{1+e^e} \\ 1-q(x_i) & \text{w/ prob. } \frac{1}{1+e^e} \end{cases}$$

$$\text{P}(\mathcal{M}(X) = y) \bigoplus e^e \mathbb{P}(\mathcal{M}(X') = y).$$

$$P(\mathcal{M}(X) = S) = \sum_{\substack{y \in S}} \mathbb{P}(\mathcal{M}(X) = y) \stackrel{e}{=} \sum_{\substack{y \in S}} \mathbb{P}(\mathcal{M}(X') = y).$$

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$$P(\mathcal{M}(X) = Y) = \mathbb{P}(\mathcal{M}(X) = Y). = P(\mathcal{M}(X_1) = \mathcal{M}_1, \mathcal{M}_2, \mathcal{M}_1) = \mathcal{M}_1, \mathcal{M}_2, \mathcal{M}_1) = \mathcal{M}_1(\mathcal{M}_1, \mathcal{M}_1) = \mathcal{M}_1, P(\mathcal{M}_1, \mathcal{M}_1) = \mathcal{M}_1).$$

$$P(\mathcal{M}(X') = Y) = \mathbb{P}(\mathcal{M}_1(\mathcal{M}_1) = \mathcal{M}_1, \mathcal{M}_2(\mathcal{M}_1) = \mathcal{M}_2). - \mathbb{P}(\mathcal{M}_1(\mathcal{M}_1) = \mathcal{M}_1).$$

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Accuracy Analysis

$$Y_{i}(x_{i}) = \begin{cases} q(x_{i}) & \text{w/ prob. } \frac{e^{e}}{1+e^{e}} > \frac{1}{2} \\ 1-q(x_{i}) & \text{w/ prob. } \frac{1}{1+e^{e}} > \frac{1}{2} \end{cases}$$

$$F(x_{i}) = \begin{cases} q(x_{i}) & \text{w/ prob. } \frac{1}{1+e^{e}} > \frac{1}{2} \\ 1-q(x_{i}) & \text{w/ prob. } \frac{1}{1+e^{e}} > \frac{1}{2} \end{cases}$$

$$F(z_{i}) = q(x_{i})$$

$$F(z_{i}) = q(x_{i}) \cdot \frac{e^{2}}{e^{2}+1} + \frac{1}{e^{2}+1}$$
Want to approximate $q(X) = \frac{1}{n} \sum_{i=1}^{n} q(x_{i})$. Claim: $\frac{1}{n} \sum_{i=1}^{n} \sum_{i=1}^{n} \sum_{i=1}^{n} \sum_{i=1}^{n} \frac{1}{e^{e}-1} \approx q(X)$

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