CSC2412: Exponential Mechanism & Private PAC Learning

Sasho Nikolov

Classification Basics

Problem: develop an algorithm that classifies avocados into ripe and unripe.

We have a big data set of avocado data. For each avocado, we have:

- colour, firmness, size, shape, skin texture, ...
- ripe or not label

From this data, we want to classify unseen avocados.



The learning problem, formally

Model: all possible setting to the features

• Known data universe $\widehat{\mathcal{X}}$ and an unknown probability distribution \underline{D} on \mathcal{X}

Fraction of the Population lateled incorrectly by c' The *error* of a concept $c' \in C$ is $\mathcal{L}_{\mathsf{D},c}(\mathcal{L}') = \mathbb{P}_{x \sim D}(c'(x) \neq c(x)).$ Loss of c' (w.r.t. D, c) We want an algorithm \mathcal{M} that outputs some $c' \in C$ and satisfies $\mathbb{P}(L_{D,c}(\mathcal{M}(X)) \leq \alpha) \geq 1 - \beta$. We want an algorithm \mathcal{M} that outputs some $c' \in C$ and satisfies $\mathbb{P}(L_{D,c}(\mathcal{M}(X)) \leq \alpha) \geq 1 - \beta$. to taken over sandomness in choosing X and any randomness of lt Probably Approximately Correct learning (PAC) [Valiant]

Empirical risk minimization

Issue: We want to find
$$\arg\min_{c'\in C} L_{D,c}(c')$$
, but we do not know D, c .
by approximate minimizers are also ok
Solution: Instead we solve $\arg\min_{c'\in C} L_X(c')$, where
 $L_X(c) = 0$
 $L_X(c') = \frac{|\{i: c'(x_i) \neq c(x_i)\}|}{n}$
is the empirical error.
Theorem (Uniform convergence) \neg Pop and emp. loss are dose for $\forall c' \in C$
Suppose that $n \ge \frac{\ln(|C|/\beta)}{2\alpha^2}$. Then, with probability $\ge 1 - \beta$,
Hoeffding's inquality
 $\max_{c'\in C} L_{D,c}(c') - L_X(c') \le \alpha$. $\forall c' \in C$
 $L_{D,c}(c') \in L_X(c')$ there is the end of the end o

In private PAC learning, we require that

 an approximately
 when X is a sample of iid labeled data points, we learn the correct concept, as in stadard PAC learning;
 M that on imput X outputs cie C

• the learning algorithm is ε -differentially private for any labeled data set $X \in (\mathcal{X} \times C)^n$.

$$\forall X, X'$$
 weighbouring $\forall S \leq C$ not
 $P(\mathcal{U}(X) \in S) \leq e^{\epsilon} \cdot P(\mathcal{U}(X') \in S)$

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Want to do ERM wy E-DP i.e. (approximately) minimize $L_{X}(c') = \frac{15i \cdot c(x_{i}) \neq c'(x_{i})51}{N} \quad \text{over } c' \in C$ How can we use Laplace noise mechanism $L_X(c)$ is a counting query Exercise: analyze We could release answers to all counting queries $C = \frac{1}{2} C_{1,\cdots}, c_k \frac{1}{2}$ $\frac{1}{2} L_X(c_1), L_X(c_2), \dots, L_X(c_k) \frac{1}{2}$

Exponential mechanism

We want to solve $\arg \min_{c' \in C} L_X(c')$.

How do we minimize with differential privacy?

Sample concepts with less error with higher probability

$$\mathbb{P}(\mathcal{M}(X) = c') \boxtimes \exp\left(-\frac{\varepsilon n}{2}L_X(c')\right)$$

Exponential Mechanism

$$\begin{array}{c} \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score function } u: \mathscr{U}^{n} \times \mathscr{Y} \to \mathbb{R} \\ \text{General set-up: score$$

The mechanism $\mathcal{M}_{ ext{exp}}(X)$ which outputs a random Y so that

$$\mathbb{P}(Y = y) = \frac{e^{\varepsilon u(X,y)/2\Delta u}}{\sum_{z \in \mathcal{Y}} e^{\varepsilon u(X,z)/2\Delta u}} \Rightarrow \text{ normalizing factor}$$

is ε -differentially private

Privacy analysis

 $\mathbb{P}(Y = y) = \frac{e^{\varepsilon u(X,y)/2\Delta u}}{\sum_{z \in \mathcal{Y}} e^{\varepsilon u(X,z)/2\Delta u}}$

$$\frac{\operatorname{Enough to show}}{\operatorname{P}(\mathcal{M}(X)=Y)} \stackrel{\times}{\to} \exp\left(\frac{\operatorname{E}(u(X,y)-u(X',y))}{2\Delta u}\right) \stackrel{\times}{\to} \exp\left(\frac{\operatorname{E}(u(X,y)-u(X',y))}{2\Delta u}\right) \stackrel{\times}{\to} \frac{\operatorname{E}^{\operatorname{E}(u(X,z)/2\Delta u}}{\sum_{z \in Y} e^{\operatorname{E}(u(X,z)/2\Delta u}}\right)$$

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Accuracy of the exponential mechanism

$$\sum_{x \in \mathcal{Y}} |u(X, y)| = \max_{y \in \mathcal{Y}} |u(X, y)|$$
Then, for the output $Y = \mathcal{M}_{exp}(X)$,

$$\mathbb{P}(u(X, Y) \leq OPT(X) - t) \leq \frac{e^{\sum_{x \in \mathcal{Y}} e^{\sum_{x \in \mathcal{Y}} u(X, x)/2 \Delta u}}{\sum_{x \in \mathcal{Y}} e^{\sum_{x \in \mathcal{Y}} u(X, x)/2 \Delta u}} \cdot |\zeta_{y}: u(X, y) \leq OPT - t)$$

$$\leq \frac{e^{\sum_{x \in \mathcal{Y}} e^{\sum_{x \in \mathcal{Y}} u(X, x)/2 \Delta u}}{e^{\sum_{x \in \mathcal{Y}} u(X, x)/2 \Delta u}} \cdot |\zeta_{y}: u(X, y) = e^{\sum_{x \in \mathcal{Y}} u(X, y)/2 \Delta u}$$

$$\leq \frac{e^{\sum_{x \in \mathcal{Y}} (OPT - t)/2 \Delta u}}{e^{\sum_{x \in \mathcal{Y}} (OPT/2 \Delta u}} \cdot |(|||_{1}) = e^{-\sum_{x \in \mathcal{Y}} t/2 \Delta u}$$

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Private Learning

Unknown distribution D on known & Unknown c in a known concept class C Data set $X = f(x_1, c(x_1)), \dots, (x_n, c(x_n))f$ where $x_{1,\dots}, x_n \cap id D$ $L_{D,c}(c') = P(c(x) \neq c'(x)) \int L_{x}(c') = \frac{Hi:c(x_i) \neq c'(x_i)f}{n}$ $If n = \frac{ln(lcl/B)}{2d^2}$ then $w_j \text{ prob } = l-f^3,$ $L_{D,c}(c') \in L_{x}(c') + d$ Exporential mechanism: sample y el w/ prob. proportional to exp(Eu(X,y)/2Dy) $Au = \max \max \left[u(X, y) - u(X', y) \right]$

A concept class C can be learned by an $\mathscr{O}s$ -differentially private mechanism when the sample size is

$$n \geq \max\left\{rac{4\ln(2|\mathcal{C}|/eta)}{arepsilon lpha}, rac{2\ln(2|\mathcal{C}|/eta)}{lpha^2}
ight\}$$

Use exp. mechanism with
$$g = C$$

with $u(X, c') = -L_X(c')$
with $prob = 1 - B_2$, c' output by $Mup(X)$ has
 $L_X(c') \leq \frac{d}{2}$. By unif. convergence, w/ prob $\geq 1 - B_2$, $L_{D,c}(c') - L_X(c') \leq \frac{d}{2}$.

Putting things together

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$$\begin{split} \mathcal{U}(X, c') &= -\mathcal{L}_{X}(c'); \quad \Delta \mathcal{U} = \frac{1}{n}; \quad \text{sample } c' \in \mathcal{C} \quad \text{wprob} \\ \text{prop. to } exp[-en \mathcal{L}_{X}(c')] \\ \text{opt}(X) &= \max -\mathcal{L}_{X}(c') \\ c' \in \mathcal{C} \\ &= -\min \mathcal{L}_{X}(c') = \mathcal{O} \\ c' \in \mathcal{C} \\ \end{split}$$

$$\begin{split} &= -\min \mathcal{L}_{X}(c') = \mathcal{O} \\ c' \in \mathcal{C} \\ &= \mathcal{O} \\ \text{prop. to } exp[-en \mathcal{L}_{X}(c')] \\ &= \mathcal{O} \\ c' \in \mathcal{C} \\ \end{split}$$

$$\begin{split} &= -\min \mathcal{L}_{X}(c') = \mathcal{O} \\ c' \in \mathcal{C} \\ &= \mathcal{O} \\ e' \in \mathcal{C} \\ \end{split}$$

$$\begin{split} &= \mathcal{O} \\ &= \mathcal{O} \\ &= \mathcal{O} \\ &= \mathcal{O} \\ e' \in \mathcal{O} \\ \end{aligned}$$

$$\begin{split} &= \mathcal{O} \\ &= \mathcal{O} \\$$